

**CS 181:**  
**NATURAL LANGUAGE**  
**PROCESSING**  
*Lecture 14: Semantic Representations*

K I M B R U C E  
P O M O N A C O L L E G E  
S P R I N G 2 0 0 8

*Disclaimer: Slide contents borrowed from many sources on web!*

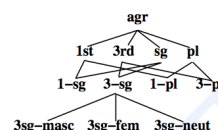
# TYPES & INHERITANCE

## CONSTRAINING FEATURES

- Each feature structure is labeled by a type.
- Types have **appropriateness conditions** expressing which features are appropriate for it.
- Types organized into **hierarchy** -- more specific types inherit properties of more abstract ones.
- The unification operation is modified to unify the types of feature structures in addition to unifying the attributes and values.

## TYPES

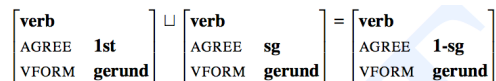
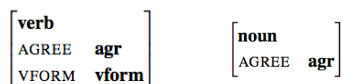
- Simple types:
  - Atomic features -- sg, pl, 3rd, ...
  - Arranged into multiple hierarchies:



- Note unification moves down hierarchy!
- Is this the best approach?

## COMPLEX TYPES

- Specify set of appropriate features
- Type restrictions on values of features
- Equality constraints between values.



## OTHER FEATURES

- Instead of writing:
  - S → NP VP
  - <NP Head Agreement> = <VP Head Agreement>
  - <S Head> = <VP Head>
- Could push categories of non-terminals into features (next slide)

## UNIFICATION PARSING

- ✱ Write:
  - ✱  $X \rightarrow Y Z$ 
    - $\langle X \text{ Cat} \rangle = S$
    - $\langle Y \text{ Cat} \rangle = NP$
    - $\langle Z \text{ Cat} \rangle = VP$
    - $\langle Y \text{ Head Agreement} \rangle = \langle Z \text{ Head Agreement} \rangle$
    - $\langle X \text{ Head} \rangle = \langle Z \text{ Head} \rangle$
- ✱ Why bother?

## POLYMORPHISM

- ✱ 'and' and 'or' polymorphism
  - ✱  $X \rightarrow Y \text{'and'} Z$ 
    - $\langle Y \text{ Cat} \rangle = \langle Z \text{ Cat} \rangle$
    - $\langle X \text{ Cat} \rangle = \langle Y \text{ Cat} \rangle$
- ✱ Requires revision of parser to look at categories rather than non-terminals.

## SUMMARY

- ✱ Feature structures allow us to capture fine-grained distinctions without multiplying our non-terminals.
- ✱ Constraints express restrictions on sentences generated.
  - ✱ Provide fixed values for lexical entries
  - ✱ Other rules have constraints requiring unification.
- ✱ Unification requires sharing.

## SEMANTIC ANALYSIS

## WHY SEMANTICS

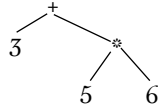
- ✱ Information retrieval
  - ✱ Some queries simple:
    - ✱ What is the current Democratic delegate count?
  - ✱ Others are not:
    - ✱ What is the city east of Claremont?
- ✱ Dialog systems, e.g., as travel agents

## SEMANTIC ANALYSIS

- ✱ How do we represent meanings of sentences?
  - ✱ First-order logic (or extensions)
  - ✱ Semantic network (graph)
  - ✱ Conceptual Dependency Graph
  - ✱ Frame-based representation (fill slots)
- ✱ Worry later about how we compute them.

## PROGRAMMING LANGUAGES

- What is meaning of  $3+5*6$ ?
- First parse into tree representing  $3 + (5 * 6)$
- Then interpret recursively from meanings of nodes
- What about  $3 + x * 6$ ?
  - Need to find meaning of  $x$  from environment



## REPRESENTATIONS

- Symbols correspond to objects, properties, and relations among objects
- True or false in some “state of affairs” or model.
- Two points of view:
  - Meanings of statements that may be true or false in model
  - Set of models in which statement is true.
- Ignore context of other statements for now

## WHAT DO WE NEED?

- Verifiability
  - Can we answer questions using representation?
- Unambiguous representation
  - Statement may be ambiguous, but better to provide several distinct semantic reps than one ambiguous one.
    - A woman in the U.S. has a baby every 15 minutes.
  - By contrast, vagueness is OK
- Canonical form
  - Identical meanings should have same representation, if possible.

## WHAT DO WE NEED?

- Inference
  - Can we deduce conclusions from semantic representation?
- Expressiveness
  - Can we represent everything necessary?
- *Prefer syntax-driven construction of meaning representations – Principle of Compositionality:*
  - *The meaning of a whole is a function of the meanings of the parts and of the way they are syntactically combined.*

## THEMATIC ROLES

- Verb subcategorization frames enforce linking of arguments with their semantic roles.
  - Ex: NP<sub>1</sub> gave NP<sub>2</sub> NP<sub>3</sub>
  - What are restrictions on each NP?
  - *Giver vs Given vs Receiver*
  - Consider meaning of gave as three-place relation *gave(x,y,z)* w/ restrictions on  $x$ ,  $y$ , and  $z$ .
  - Ex: NP<sub>1</sub> was given NP<sub>2</sub> by NP<sub>3</sub>
- Prepositions similar: NP<sub>1</sub> on NP<sub>2</sub>.

## THEMATIC ROLES

- Must support relations with
  - Semantic labeling of arguments
  - Semantic constraints on allowable arguments

## MEANING

- Meaning will be understood as a relation between
  - Semantic representations (e.g., FOL)
  - State-of-Affairs (SOA) or model providing meaning of representations.

## MODEL

- Model should represent all relevant info about current context
  - Set(s) of objects -- individual elements
  - Names of objects
  - Properties of objects -- unary relations
  - Relations between objects -- n-ary relations

## VOCABULARY OF REPRESENTATION

- Non-logical vocabulary: Open-ended set of names for objects, properties, and relations of model.
- Logical vocabulary: Closed set of symbols, operators, quantifiers, etc.
  - Allow us to compose more complex expressions from simpler ones.

## MODEL

- Model consists of a mapping of all non-logical symbols to objects, sets and sets of tuples of objects of the model.
- Logical symbols are used to build up meaning of more complex representations in the model.

## FIRST-ORDER LOGIC

- Non-logical symbols,  $L$ , consist of:
  - a finite set of constant names:  $c_1, \dots, c_n$
  - a finite set of function symbols:  $f_1^{(n_1)}, \dots, f_m^{(n_m)}$  where  $f^{(j)}$  is a function symbol of arity  $j$ .
  - a finite set of relation symbols:  $R_1^{(k_1)}, \dots, R_m^{(k_m)}$  where  $R^{(j)}$  is a  $j$ -ary relation symbol.
- Logical symbols:
  - potentially infinite set of variables:  $v_1, \dots, v_j, \dots$
  - logical connectives:  $\neg, \wedge, \vee, \Rightarrow$
  - Quantifiers:  $\forall, \exists$

## TERMS

- A term  $t$  is one of:
  - a constant,  $t = c_i$ .
  - a variable,  $t = v_j$ .
  - a function application,  $t = f_j^{(k)}(t_1, \dots, t_k)$  where  $f_j^{(k)}$  is a  $k$ -ary function symbol and  $t_1, \dots, t_k$  are terms.
- Example:  $f^{(3)}(g^{(2)}(c, x), y, z)$

## FORMULAS

- \* A formula  $\phi$  is one of:
  - \* An atomic formula,  $\phi = R_i^{(j)}(t_1, \dots, t_j)$ , where  $R_i^{(j)}$  is a  $j$ -ary relation symbol and  $t_1, \dots, t_j$  are terms.
  - \*  $\phi = \neg\psi$ ,  $\phi = \psi \vee \theta$ ,  $\phi = \psi \wedge \theta$ , or  $\phi = \psi \Rightarrow \theta$ , where  $\psi$  and  $\theta$  are formulas
  - \*  $\phi = \forall x.\theta$  or  $\phi = \exists x.\theta$ , where  $x$  is a variable and  $\theta$  is a formula.
  - \* Quantifiers have tighter binding than connectives
  - \* Can define scope of variable, bound and free variables

## FOL MODEL

- \* A *model*  $M$  for language  $L$  is an ordered pair  $\langle D, F \rangle$  consisting of a non-empty domain,  $D$ , and an interpretation function  $F$  mapping the non-logical symbols of the language  $L$  to elements of the model,  $n$ -ary functions, or  $m$ -ary relations to match the arity of the symbols.
- \* There are also typed models with a collection of domains  $D_i$ .

## ASSIGNMENTS

- \* Cannot determine meaning of term or assignment unless know meaning of its free variables.
- \*  $g$  is an *assignment* in  $M$  iff  $g$  is a partial function from the set of variables to  $D$ .
- \* If  $g$  is an assignment in  $M$ , then if  $d$  in  $D$ , define  $g[d/x]$  to be function  $g'$  s.t. for all  $y \neq x$ ,  $g'(y) = g(y)$ , and  $g'(x) = d$ .

## MEANING OF TERMS

- \* Let  $\tau$  be a term,  $M = \langle D, F \rangle$  be a model, and  $g$  be an assignment in  $M$  whose domain includes all variables in  $\tau$ . Define the meaning of  $\tau$  by induction as follows:
  - \*  $M^g(c) = F(c)$ .
  - \*  $M^g(x) = g(x)$ .
  - \*  $M^g(f(t_1, \dots, t_n)) = F(f)(M^g(t_1), \dots, M^g(t_n))$ .

## MEANING OF FORMULAS

- \* Let  $\phi$  be a formula,  $M = \langle D, F \rangle$  be a model, and  $g$  be an assignment in  $M$  whose domain includes all variables in  $\phi$ . Define the satisfaction of  $\phi$  by  $(M, g)$  by induction:
  - \*  $(M, g) \models (R(t_1, \dots, t_n))$  iff  $(M^g(t_1), \dots, M^g(t_n)) \in F(R)$ .
  - \*  $(M, g) \models \psi \wedge \theta$  iff  $(M, g) \models \psi$  and  $(M, g) \models \theta$
  - \*  $(M, g) \models \psi \vee \theta$  iff  $(M, g) \models \psi$  or  $(M, g) \models \theta$
  - \*  $(M, g) \models \neg\psi$  iff not  $(M, g) \models \psi$

## MEANING OF FORMULAS

- \* Let  $\phi$  be a term,  $M = \langle D, F \rangle$  be a model, and  $g$  be an assignment in  $M$  whose domain includes all variables in  $\phi$ . Define the satisfaction of  $\phi$  by  $(M, g)$  by induction:
  - \*  $(M, g) \models \exists x.\psi$  iff there is an  $a \in D$ , s.t.  $(M, g[a/x]) \models \psi$ .
  - \*  $(M, g) \models \forall x.\psi$  iff for all  $a \in D$ ,  $(M, g[a/x]) \models \psi$ .

## EXAMPLE

- \* Model w/  $D$  = set of students in room and the chairs, w/interpretations for
  - \* names,
  - \* awake:  $D \rightarrow \text{boolean}$
  - \*  $\text{ifo} \subseteq D \times D$
  - \*  $\text{isi} \subseteq D \times D$
- \* Write sentences and determine which true.

**ANY QUESTIONS?**