Homework 3

Due 27 September

Handout 7
CSC 131: Fall, 2006
20 September

Reading

1. Read Mitchell, skim 4.4–4.5, 5.3-5.5, skim 6.1.
2. Read Ullman, as necessary for programming assignments.
3. *(Optional)* J. Backus, Can programming be liberated from the von Neuman style?, Comm. ACM 21, 8 (1978) 613-641. You can find this on the cs131 web site.

Problems

1. (*10 points*) ..... Translation into Lambda Calculus

Mitchell, Problem 4.6

2. (*20 points*) ..... Lazy Evaluation and Parallelism

Mitchell, Problem 4.11

The function g should be defined as follows (there is a typo in the book):

```
fun g(x, y) = if x = 0
              then 1
              else if x + y = 0
                   then 2
                   else 3;
```

3. (*5 points*) ..... Algol 60 Procedure Types

Mitchell, Problem 5.1

4. (*10 points*) ..... ML Types

Mitchell, Problem 6.1

5. (*50 points*) ..... ML Programming

For this problem, use the ml interpreter on the Unix machines in the computer lab. To run the program in the file "example.sml", either

- type

  ```
  -> sml < example.sml
  ```

  at the command line.

- open the file in emacs and then type C-c C-b to evaluate the buffer.

As with Lisp, the ML compiler will process the program in the file and print the result. For example, if "example.sml" contains

```
(* double an integer *)
fun double x = x + x;

(* return the length of a list *)
fun listLength nil = 0
  | listLength (l::ls) = 1 + listLength ls
  ;

double (10);
listLength (1::[2,3,4]);
```

tlaoding the program will produce the following:

```
val double = fn : int -> int
val listLength = fn : 'a list -> int
val it = 100 : int
val it = 4 : int
```

You can also run "sml" and enter in declarations and expressions to evaluate at the prompt.

Start early on this part so you can see the TA or me if you have problems understanding the language. Looking at the examples in Mitchell and Ullman books and in your notes will help a great deal in understanding how to use ML. Use pattern matching where possible.

Comments in ML appear inside `(*` and `*)` characters. Also, put the following line at the top of your file to ensure that large data types are fully printed:

```
Control.Print.printDepth:= 100;
```

(a) **Basic Functions**

Define a function `sumSquares` that, given a nonnegative integer n, returns the sum of the squares of the numbers from 1 to n:

```
- sumSquares 4;
val it = 30 : int
- sumSquares 5;
val it = 55 : int
```

Define a function `listDup` that takes a pair of an element, e, of any type, and a non-negative number, n, and returns a list with n copies of e:

```
- listDup("moo", 4);
val it = ["moo","moo","moo","moo"] : string list
- listDup(1, 2);
val it = [1,1] : int list
- listDup(listDup("cow", 2), 2);
val it = [["cow","cow"],["cow","cow"]] : string list list
```

Your function will have a type like `'a * int -> 'a list`. What does this type mean? Why is it the appropriate type for your function.

(b) **Zipping and Unzipping**

Write the function `zip` to compute the pairwise interleaving of two lists of arbitrary length. You should use pattern matching to define this function:

```
- zip [1,3,5,7] ["a","b","c","de"];
val it = [(1,"a"),(3,"b"),(5,"c"),(7,"de")]: (int * string) list
```

Note: If the lists don't have the same length, you may decide how you would like the function to behave. If you don't specify any behavior at all you will get a warning from the compiler that you have not taken care of all possible patterns— this is fine.

Write the inverse function, `unzip`, which behaves as follows:

```
- unzip [(1,"a"),(3,"b"),(5,"c"),(7,"de")];
val it = ([1,3,5,7], ["a","b","c","de"]): int list * string list
```

Write `zip3`, to zip three lists.

```
- zip3 [1,3,5,7] ["a","b","c","de"] [1,2,3,4];
val it = [(1,"a",1),(3,"b",2),(5,"c",3),(7,"de",4)]: (int * string * int) list
```

Why can't you write a function `zip_any` that takes a list of any number of lists and zips them into tuples? From the first part of this question it should be pretty clear that for any fixed n, one can write a function `zipn`. The difficulty here is to write a single function that works for all n! I.e., can we write a single function `zip_any` such that `zip_any [list1,list2,...,listk]` returns a list of k-tuples no matter what k is?

(c) **find**

Write a function find with type `''a * ''a list -> int` that takes a pair of an element and a list and returns the location of the first occurrence of the element in the list. For example:

```
- find(3, [1, 2, 3, 4, 5]);
val it = 2 : int
- find("cow", ["cow", "dog"]);
val it = 0 : int
- find("rabbit", ["cow", "dog"]);
val it = ~1 : int
```

First write a definition for find where the element is guaranteed to be in the list. Then, modify your definition so that it returns `~1` if the element is not in the list.

(d) **Trees**

Here is the datatype definition for a binary tree storing integers at the leaves:

```
datatype IntTree = LEAF of int | NODE of (IntTree * IntTree);
```

Write a function `sum:IntTree -> int` that adds up the values in the leaves of a tree:

```
- sum(LEAF 3);
val it = 3 : int
- sum(NODE(LEAF 2, LEAF 3));
val it = 5 : int
- sum(NODE(LEAF 2, NODE(LEAF 1, LEAF 1)));
val it = 4 : int
```

Write a function `height : IntTree -> int` that returns the height of a tree:

```
- height(LEAF 3);
val it = 1 : int
- height(NODE(LEAF 2, LEAF 3));
val it = 2 : int
- height(NODE(LEAF 2, NODE(LEAF 1, LEAF 1)));
val it = 3 : int
```

Write a function `balanced: IntTree -> bool` that returns true if a tree is balanced (ie, both subtrees are balanced and differ in height by at most one). You may use your height function above.

```
- balanced(LEAF 3);
val it = true : bool
- balanced(NODE(LEAF 2, LEAF 3));
val it = true : bool
- balanced(NODE(LEAF 2, NODE(LEAF 3, NODE(LEAF 1, LEAF 1))));
val it = false : bool
```

Is your implementation as efficient as possible? What is wrong with using the height? function in the definition of `balanced`? How would you write `balanced` to be more efficient? (You need not write code, but describe how you would do this.)

(e) **Stack Operations**

Certain programming languages (and HP calculators) evaluate expressions using a stack. As some of you may know, PostScript is a programming language of this ilk for describing images when sending them to a printer. We are going to implement a simple evaluator for such a language. Computation is expressed as a sequence of operations, which are drawn from the following data type:

```
datatype OpCode =
    PUSH of real
  | ADD
  | MULT
  | SUB
  | DIV
  | SWAP
  ;
```

The operations have the following effect on the operand stack. (The top of the stack is shown on the left.)

OpCode	Initial Stack	Resulting Stack
`PUSH(r)`	...	r ...
`ADD`	a b ...	(b + a) ...
`MULT`	a b ...	(b * a) ...
`SUB`	a b ...	(b - a) ...
`DIV`	a b ...	(b / a) ...
`SWAP`	a b ...	b a ...

The stack may be represented using a list for this example, although we could also define a stack data type for it.

```
type Stack = real list;
```

Write a recursive evaluation function with the signature

```
eval : OpCode list * Stack -> real
```

It takes a list of operations and a stack. The function should perform each operation in order and return what is left in the top of the stack when no operations are left. For example,

```
eval([PUSH(2.0),PUSH(1.0),SUB],[])
```

returns 1.0. The `eval` function will have the following basic form:

```
fun eval (nil,a::st) = (* ... *)
  | eval (PUSH(n)::ops,st) = (* ... *)
  | (* ... *)
  | eval (_,_) = 0.0
  ;
```

You need to fill in the blanks and add cases for the other opcodes.

The last rule handles illegal cases by matching any operation list and stack not handled by the cases you write. These illegal cases include ending with an empty stack, performing addition when fewer than two elements are on the stack, and so on. You may ignore divide-by-zero errors for now (or look at exception handling in Ullman– we will cover that topic in a few weeks).