

Homework 7

Due midnight, Thursday, 3/26/2015

Please submit your homework solutions online at <http://www.dci.pomona.edu/tools-bin/cs081upload.php>. If you have more than one file to be turned in, please put it in a folder and zip it up before turning it in.

Problems from the texts are given in the form c.n where c is the chapter and n is the problem number. Thus problem 2.7 is problem 7 from Chapter 2.

1. (5 points) **Predicate Logic**

Problem 2.1.3 from H & R page 157.

2. (10 points) **Free & Bound Variables**

- (a) In the formula below, or a copy of it, underline the occurrences of free variables and circle the occurrences of bound variables. (Consider the variable right after a quantifier to be bound.)

$$\forall z \left(P(x) \wedge \forall x \exists y (Q(y, f(z)) \rightarrow R(g(g(w, x), f(y)))) \right)$$

Using the definitions included in the LaTeX file, the command `\encircle x` can be used to draw a letter with a circle around it: $\text{\encircle{x}}$ `\underline{x}` will generate \underline{x} .

- (b) Let ϕ be the formula in part a. Write out the formulas below, renaming bound variables when necessary to avoid “capture.”
- i. $\phi[f(y)/x]$
 - ii. $\phi[f(y)/z]$
 - iii. $\phi[f(y)/w]$

3. (10 points) **Proofs**

Exactly one of the inferences below is correct. Give a (constructive) proof of the correct member and a short intuitive explanation of why the other is incorrect.

- (a) $\forall x (\phi \vee \psi) \vdash \forall x \phi \vee \forall x \psi$ and
 (b) $\forall x \phi \vee \forall x \psi \vdash \forall x (\phi \vee \psi)$

4. (10 points) **Bad Proof**

Recall the axioms for an equivalence relation \sim :

- A0. $\forall x (x \sim x)$
 A1. $\forall x \forall y (x \sim y \rightarrow y \sim x)$
 A2. $\forall x \forall y \forall z (x \sim y \rightarrow (y \sim z \rightarrow x \sim z))$

Consider the following proof sketch that purports to prove that $A1, A2 \vdash A0$. Formalize the sketch and use your formulation to identify the error in it.

Let x be arbitrary and suppose that $x \sim y$. Then by A1, $y \sim x$. Using these two facts, we may use A2 to conclude that $x \sim x$. Because x was arbitrary, we have $\forall x (x \sim x)$.

5. (10 points) **Proofs**

Problem 2.3.11ac from H & R page 162.

6. (10 points) **Semantics**

Problem 2.4.5ab from H & R page 163.

7. (10 points) **Consistency**

Problem 2.4.11ad from H & R page 164. Show each set of formulas is consistent by finding a model for all formulas in the set.

8. (5 points) **Validity**

Problem 2.5.1d from H & R page 165.

9. (5 points) **Compactness Theorem**

Consider a predicate logic with a two-place predicate symbol $<$, as well as the usual arithmetic operators $+$ and $*$. Let \mathbb{N} be the model of this language that consists of the natural numbers with their natural ordering and where $+$ and $*$ have their ordinary meanings. Let $\text{Th}(\mathbb{N})$ be the set of all sentences of this language that are true in \mathbb{N} . Use the compactness theorem to show that there is another model \mathfrak{M} of this language that satisfies all of the sentences in $\text{Th}(\mathbb{N})$, but also has an infinite descending chain, i.e., an infinite set of elements of the model $\{a_i | i > 0\}$ such that for all $i > 0$, $a_{i+1} <^{\mathfrak{M}} a_i$.