

# Lecture 30: Maps & Hashing Continued

Fall 2016

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## This week

- Lab 10: introduction to C (input/output)
  - Come prepared
- Assignment 10: priority queue in C
  - Implement a heap data structure
  - Review heap algorithms from book
  - Read Bailey's heap implementation before starting

## Naïve Version

```
public class Map<K, V> {  
    protected V[] entries;  
  
    public V get(K key) {  
        int index = key.hashCode() % entries.length;  
        return entries[index];  
    }  
    public void put(K key, V value) {  
        int index = key.hashCode() % entries.length;  
        entries[index] = value;  
    }  
}
```

Warning: This code is simplified!

## Hash Collisions

- `k1.hashCode() == k2.hashCode()` but `k1 != k2`
  - May also be caused by the modulus operation
- This is inevitable (e.g., the birthday paradox)
- A “good” hash function rarely collides

## Avoiding Collisions

Two main strategies:

1. Open addressing (rehashing):
  - Linear probing (use the next open space)
  - Cuckoo hashing
2. Bucketing (wide cells):
  - Separate chaining

## Linear Probing

- If we collide, check next entry until one is empty
- Deletion is complicated
- Can only hold `entries.length` items
- Resizing the table requires rehashing everything

## Separate Chaining

- Turn each entry into a linked list (or array, etc.)
- On collision add to the bucket
- Searching list is fast if lists are small
- Deletion is simple
- Can hold more than `entries.length` items easily
- In practice slower than linear probing

## Cuckoo Hashing

- Create two tables `table1` and `table2`
- Create an alternate hash function

```
int altHash(Object o)
```

  - Can't use `hashCode...`
- On collision kick out an item and take its spot
- Homeless item tries to get into the other table
  - Repeat as necessary

## Cuckoo Hashing

- Lookup is worst-case  $O(1)$  (only two places)
- Deletion is easy
- Can cause an infinite loop (!)
  - Unlikely if hash functions are random
- Fast in practice

## Load Factor

- Performance depends on *load factor*
- Load factor is  $\alpha = \frac{n}{N}$  where  $n$  = items in table and  $N$  = size of table
- Higher load factor  $\rightarrow$  more collisions  $\rightarrow$  slow
- Usually want to ensure  $\alpha < 0.75$
- Generally  $\alpha > 0.75$  means resize the table (& rehash everything)

We'll analyze hash table performance in lab next week

## Hashing and Equality

```
public class Point {
    public int x, y;

    public boolean equals(Object other) {
        if (other instanceof Point) {
            return (this.x == other.x
                && this.y == other.y);
        }
        return false;
    }

    public int hashCode() {return addr(this);}
}
```

## Respecting Equality

```
public class Point {
    public int x, y;

    public boolean equals(Object other) {
        ...
    }

    public int hashCode() {
        return (31 * this.x) ^ this.y;
    }
}
```

See: [http://www.javamex.com/tutorials/collections/hash\\_function\\_guidelines.shtml](http://www.javamex.com/tutorials/collections/hash_function_guidelines.shtml)