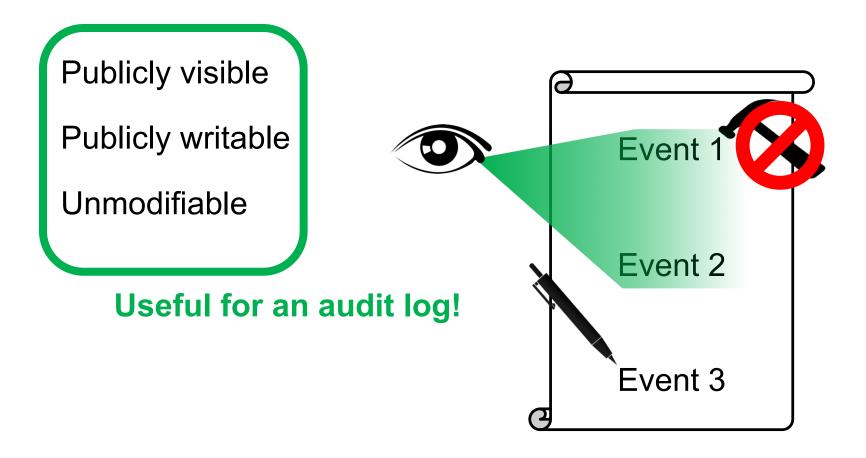
Blockchains and Audit

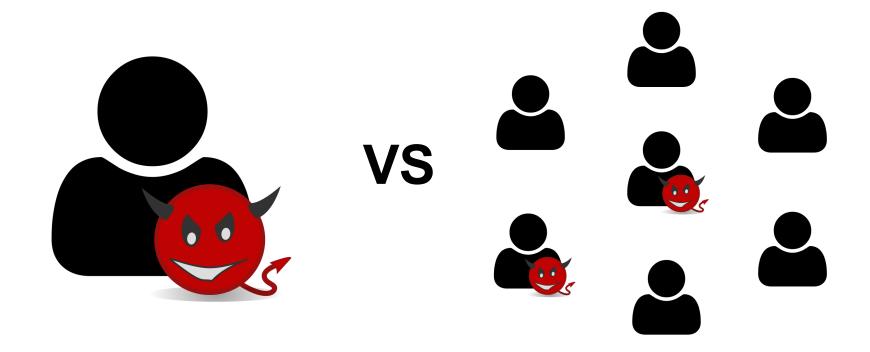
CS 181S

November 28, 2018

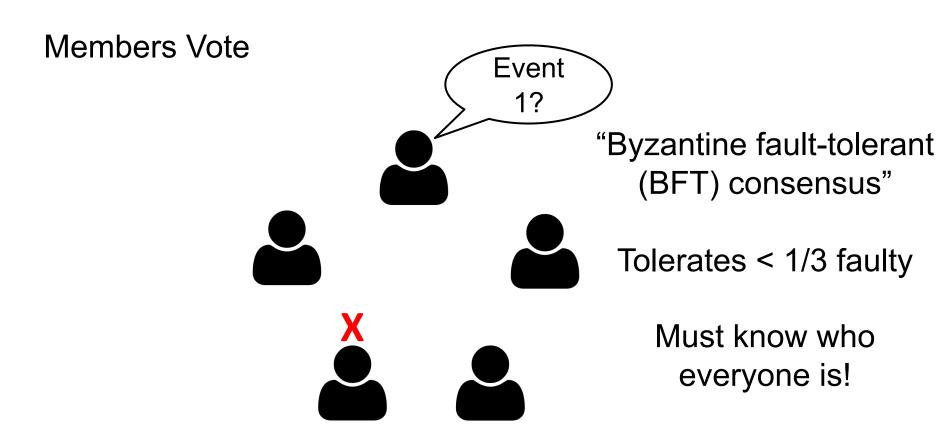
Blockchain: A public tamper-proof log



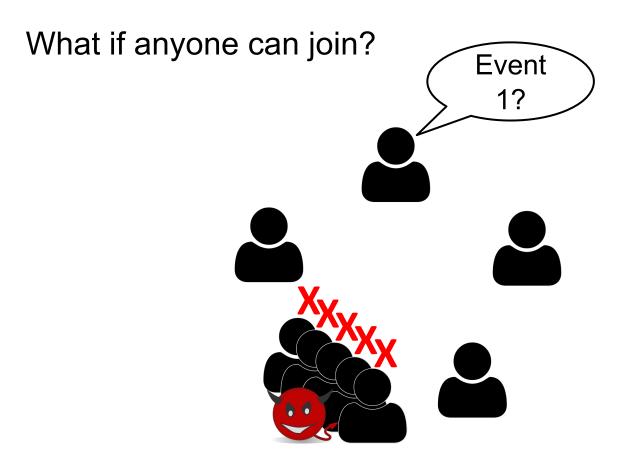
Preventing Tampering



Traditional Consensus



Sybil Attacks



Defending against Sybil

Need a scarce resource

- BFT consensus uses identity you only get one
- What else can we use?
 - Money (Proof of Stake)
 - Computational power (Proof of Work)

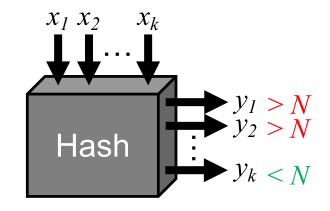
COMPUTATION AS A SCARCE RESOURCE: PROOF OF WORK

Proof of Work: The basics

Find x such that Hash(x) < NThis could take a while...

What about replays?

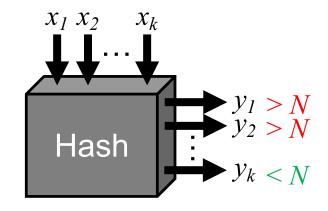
Add a nonce rLook for $Hash(r \parallel x) \leq N$



Proof of Work: Building a log

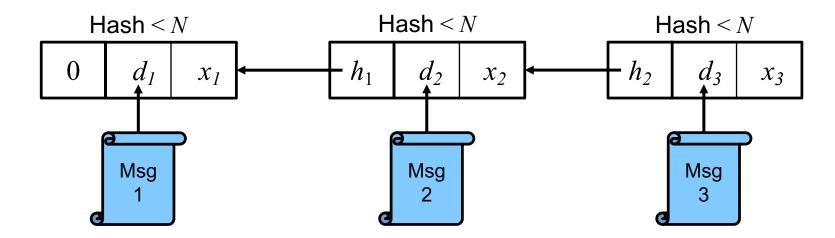
Make the nonce useful Use a message digest!

> $d = \mathsf{Digest}(m)$ Find x such that $\mathsf{Hash}(d \parallel x) < N$

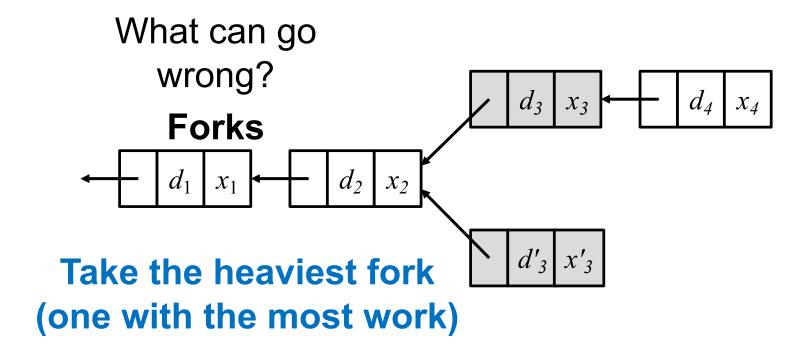


Proof of Work: Building a log

- 1. To add a message, generate a proof of work with that message
- 2. Connect each message to previous



Proof of Work: Coming to consensus



Nakamoto Consensus



If majority of computation is honest, honest parties will agree (eventually)

- Log is tamper-proof
 - It would require redoing all of the work to tamper

Blockchains for Audit

- Individual accountability
 - Everything is visible. Everyone is accountable.
- Event reconstruction
 - All of the events are there. Easy to reconstruct.
- Real-time intelligence
 - Miners can verify everything it goes on the log.

before!

Not just a log!

Authoritative record

- Instead of logging events elsewhere, the blockchain can record the definition of events (e.g. transactions)
- Online validation can prevent illegal events from ever happening!

What restrictions make sense?

Transaction Processing System

- Each block has a limited number of transactions (1 MB)
- Transactions cannot create money
 - Except coinbase transaction to reward miner
- Coins can only be spent once (spending creates new unspent coins)
- To spend a coin conditions must be met (e.g., owner authorizes)

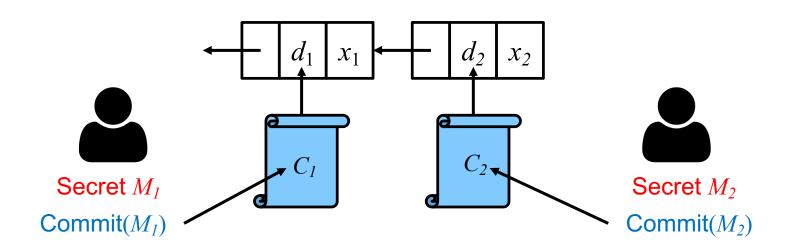
Bitcoin



BLOCKCHAINS AND CONFIDENTIALITY

What do we do with private data?

Cannot put it on the blockchain – everything is public Only publish commitments



Commitment Schemes

- A commitment scheme Com is a two-phase, two-party protocol such that:
 - Secrecy: receiver does not learn anything about x from Com(x)
 - Binding: sender cannot produce alternative x' such that Com(x) = x'

- Example Protocol
 - 1. B->A: r
 - A: choose random bit b. If b=0, Com(x) = Hash(x) else Com(x) = Hash(x) xor r
 - 3. A->B: Com(x)
 - . . .
 - 1. A->B: x

What do we do with private data?

Cannot put it on the blockchain – everything is public Only publish commitments

- Still tamper-proof
- No longer able to see actions
 - Cannot reconstruct events X
 - Cannot perform online validation X

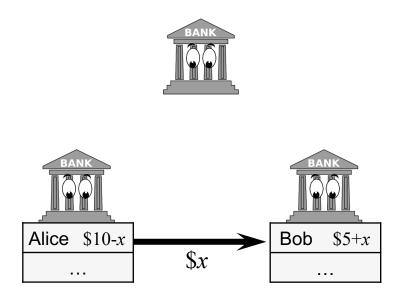


Doing better with private data

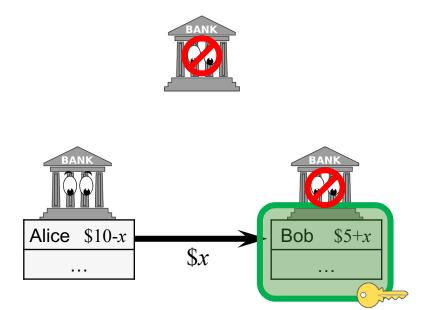
Verify data validity without leaking secrets Ongoing research with two main tools

- 1. Heavy-duty cryptographic constructs
 - Complex zero-knowledge proofs
- 2. Trusted hardware
 - Places trust in hardware instead of crypto or a large group

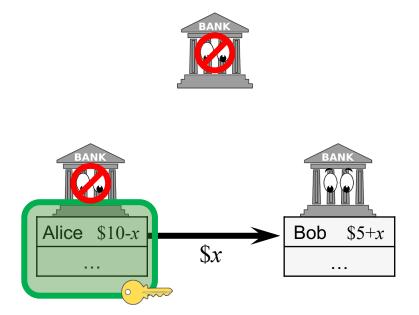
Bank-base confidential transactions



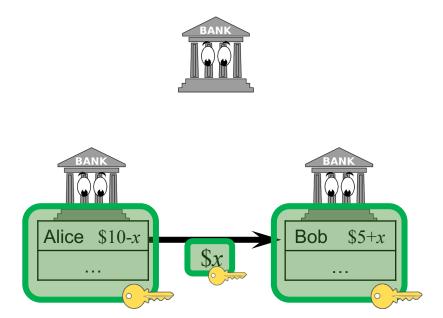
Sending bank can see sender and value



Receiving bank can see recipient and value



Everyone can see banks involved



Strong publicly verifiable integrity guarantees

- Sender authorized transaction
- Sender had money to send
- Transaction value was not negative
- Transaction was processed correctly

Can (provably) furnish transaction details to external

auditor





Trusted Hardware

Special machine instructions

Isolate process from the surrounding system

Can remotely attest that they're running specific code

Uses (literally) hard-wired keys in the CPU

Trustworthy code can operate on secret data and attest to correctness

Examples:

- Intel Software Guard eXtensions (SGX)
- ARM TrustZone