# Problem Session 3: Dynamic Memory 

Wednesday, April 15, 2020

1. Consider an allocator that uses an implicit free list. Assume that all blocks (both allocated and free) contain at 32-bit header and a 32-bit footer. Assume that each block has a total size (including header and footer) that is a multiple of 8 bytes, thus only the 29 higher order bits in the header and footer are needed to record block size. The usage of the remaining 3 lower order bits is as follows:

- bit 0 indicates the use of the current block: 1 for allocated, 0 for free.
- bit 1 indicates the use of the previous adjacent block: 1 for allocated, 0 for free.
- bit 2 is unused and is always set to be 0 .

Given the contents of the heap shown on the left, show the new contents of the heap (in the right table) after a call to free ( $0 \times 400 \mathrm{~b} 010$ ) is executed. Your answers should be given as hex values. Note that the address grows from bottom up. Assume that the allocator uses immediate coalescing, that is, adjacent free blocks are merged immediately each time a block is freed.

| Address |  |
| :---: | :---: |
| 0x400b028 | 0x00000012 |
| 0x400b024 | 0x400b611c |
| 0x400b020 | 0x400b512c |
| 0x400b01c | 0x00000012 |
| 0x400b018 | 0x00000013 |
| 0x400b014 | 0x400b511c |
| 0x400b010 | 0x400b601c |
| 0x400b00c | 0x00000013 |
| 0x400b008 | 0x00000013 |
| 0x400b004 | 0x400b601c |
| 0x400b000 | 0x400b511c |
| 0x400affc | 0x00000013 |


| Address |  |
| :---: | :---: |
| 0x400b028 |  |
| 0x400b024 | 0x400b611c |
| 0x400b020 | 0x400b512c |
| 0x400b01c |  |
| 0x400b018 |  |
| 0x400b014 | 0x400b511c |
| 0x400b010 | 0x400b601c |
| 0x400b00c |  |
| 0x400b008 |  |
| 0x400b004 | 0x400b601c |
| 0x400b000 | 0x400b511c |
| 0x400affc |  |

2. Assume you are running on the same system as in the previous problem. Five helper routines are defined to facilitate the implementation of free (void $*$ p). The functionality of each routine is explained in the comment above the function definition. Fill in the body of the helper routines the code section label that implement the corresponding functionality correctly.
```
/* given a pointer p to an allocated block, i.e., p is a pointer returned by some
previous malloc() call; returns the pointer to the header of the block */
void * header(void * p){
    void *ptr;
    _----_-;
    return ptr;
}
A. ptr=p-1
B. ptr=(void *)((int *)p-1)
C. ptr=(void *)((int *)p-4)
/* given a pointer to a valid block header, returns the size of the block */
int size(void *hp){
    int result;
    _-_---_;
    return result;
}
A. result=(*hp)&(~7)
B. result=((*(char *)hp)&(~5))<<2
C. result=(*(int *)hp)&(~7)
```

A. ptr=p+size(header (p))-8
B. $p t r=p+s i z e(h e a d e r(p))-4$
C. $\operatorname{ptr}=($ int $*) p+$ size (header $(p))-2$

```
```

/* given a pointer p to an allocated block,i.e. p is a pointer returned by some

```
/* given a pointer p to an allocated block,i.e. p is a pointer returned by some
previous malloc() call; returns the pointer to the footer of the block */
previous malloc() call; returns the pointer to the footer of the block */
void * footer(void * p){
void * footer(void * p){
    void *ptr;
    void *ptr;
    ___-_-_;
    ___-_-_;
    return ptr;
    return ptr;
}
```

}

```
```

/* given a pointer to a valid block header, returns the usage of the currect block,
1 for allocated, O for free */
int allocated(void * hp){
int result;
___-_-_;
return result;
}
A. result=(*(int $*) \mathrm{hp}) \& 1$
B. result $=(*$ (int $* h p) \& 0$
C. result=(*(int $*) h p) \mid 1$
/* given a pointer to a valid block header, returns the pointer to the header of previous block in memory */
void * prev(void *hp)\{
void *ptr;

```
\(\qquad\)
```

return ptr;
\}
A. $\mathrm{ptr}=\mathrm{hp}-\operatorname{size}(\mathrm{hp})$
B. $\mathrm{ptr}=\mathrm{hp}-\operatorname{size}(\mathrm{hp}-4)$
C. $\operatorname{ptr}=\mathrm{hp}-\operatorname{size}(\mathrm{hp}-4)+4$

```
```

