## Lecture 1: Introduction to Computer Systems

CS 105 January 22, 2020

## Abstraction



## Correctness

- Example 1: Is $\mathbf{x}^{2} \geq \mathbf{0}$ ?
- Floats: Yes!
- Ints:

- 40000 * $40000 \rightarrow 1600000000$
- 50000 * $50000 \rightarrow$ ??
- Example 2: Is $(x+y)+z=x+(y+z) ?$
- Ints: Yes!
- Floats:
- $\left(2^{\wedge} 30+-2^{\wedge} 30\right)+3.14 \rightarrow 3.14$
- $2^{\wedge} 30+\left(-2^{\wedge} 30+3.14\right) \rightarrow$ ??


## Performance

```
* void copyij(int src[2048][2048],
* void copyij(int src[2048][2048],
* void copyij(int src[2048][2048],
* void copyij(int src[2048][2048],
* void copyij(int src[2048][2048],
* void copyij(int src[2048][2048],
* void copyij(int src[2048][2048],
* void copyij(int src[2048][2048],
```


## 4.3 ms

## 81.8 ms

- Hierarchical memory organization
- Performance depends on access patterns
- Including how step through multi-dimensional array


## Security

```
void admin_stuff(int authenticated) {
    if(authenticated) {
        // do admin stuff
    }
}
int dontTryThisAtHome(char * user_input, int size) {
    char data[size];
    int ret = memcpy(*user_input, data);
    return ret;
}
```


## Bits

- a bit is a binary digit that can have two possible values
- can be physically represented with a two state device



## Bits



## Storing bits

- Static random access memory (SRAM): stores each bit of data in a flip-flop, a circuit with two stable states
- Dynamic Memory (DRAM): stores each bit of data in a capacitor, which stores energy in an electric field (or not)
- Magnetic Disk: regions of the platter are magnetized with either N-S polarity or S-N polarity
- Optical Disk: stores bits as tiny indentations (pits) or not (lands) that reflect light differently
- Flash Disk: electrons are stored in one of two gates separated by oxide layers


## Bytes and Memory

- A byte is a unit of eight bits
- Memory is an array of bytes
- An index into the array is an address, location, or pointer
- Often expressed in hexadecimal
- We speak of the value in memory at an address
- The value may be a single byte ...
- ... or a multi-byte quantity starting at that address


## Binary Numbers

$$
\begin{gathered}
4211 \\
=4 \cdot 10^{3}+2 \cdot 10^{2}+1 \cdot 10^{1}+1 \cdot 10^{0} \\
=4211
\end{gathered}
$$

## 1011

$$
\begin{gathered}
=1 \cdot 2^{3}+0 \cdot 2^{2}+1 \cdot 2^{1}+1 \cdot 2^{0} \\
=11
\end{gathered}
$$

Binary Numbers


## Hexidecimal Numbers

- Use digits $0,1,2,3,4,5,6,7,8,9, A, B, C, D, E, F$
- Compute numbers base 16 1011

$$
\begin{gathered}
=1 \cdot 2^{3}+0 \cdot 2^{2}+1 \cdot 2^{1}+1 \cdot 2^{0} \\
=11
\end{gathered}
$$

$$
=1 \cdot 10^{3}+0 \cdot 10^{2}+1 \cdot 10^{1}+1 \cdot 10^{0}
$$

$$
=1011
$$

$$
=1 \cdot 16^{3}+0 \cdot 16^{2}+1 \cdot 16^{1}+1 \cdot 16^{0}
$$

$$
=4113
$$

- one byte is two digits in hex


## ASCII characters

| Char | Dec | Oct | Hex | Char | Dec | Oct | Hex | \| Char | Dec | Oct | Hex |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| (sp) | 32 | 0040 | 0x20 | @ | 64 | 0100 | 0x40 | \| | 96 | 0140 | 0x60 |
| ( | 33 | 0041 | $0 \times 21$ | A | 65 | 0101 | 0x41 | a | 97 | 0141 | $0 \times 61$ |
| " | 34 | 0042 | $0 \times 22$ | B | 66 | 0102 | 0x42 | b | 98 | 0142 | $0 \times 62$ |
| \# | 35 | 0043 | $0 \times 23$ | C | 67 | 0103 | $0 \times 43$ | \| c | 99 | 0143 | $0 \times 63$ |
| \$ | 36 | 0044 | $0 \times 24$ | D | 68 | 0104 | 0x44 | d | 100 | 0144 | 0x64 |
| \% | 37 | 0045 | 0x25 | E | 69 | 0105 | 0x45 | \| e | 101 | 0145 | 0x65 |
| \& | 38 | 0046 | $0 \times 26$ | F | 70 | 0106 | $0 \times 46$ | \| f | 102 | 0146 | $0 \times 66$ |
| ' | 39 | 0047 | $0 \times 27$ | G | 71 | 0107 | $0 \times 47$ | g | 103 | 0147 | $0 \times 67$ |
| ( | 40 | 0050 | 0x28 | H | 72 | 0110 | 0x48 | \| h | 104 | 0150 | 0x68 |
| ) | 41 | 0051 | 0x29 | I | 73 | 0111 | 0x49 | i | 105 | 0151 | 0x69 |
| * | 42 | 0052 | 0x2a | J | 74 | 0112 | 0x4a | \| ${ }^{\text {j }}$ | 106 | 0152 | $0 \times 6 \mathrm{a}$ |
| + | 43 | 0053 | 0x2b | K | 75 | 0113 | 0x4b | \| k | 107 | 0153 | 0x6b |
| , | 44 | 0054 | 0x2c | L | 76 | 0114 | 0x4c | I | 108 | 0154 | $0 \times 6 \mathrm{c}$ |
| - | 45 | 0055 | 0x2d | M | 77 | 0115 | 0x4d | m | 109 | 0155 | 0x6d |
|  | 46 | 0056 | 0x2e | N | 78 | 0116 | 0x4e | \| $n$ | 110 | 0156 | $0 \times 6 \mathrm{e}$ |
| 1 | 47 | 0057 | 0x2f | 0 | 79 | 0117 | 0x4f | \| 0 | 111 | 0157 | 0x6f |
| 0 | 48 | 0060 | $0 \times 30$ | P | 80 | 0120 | $0 \times 50$ | \| $p$ | 112 | 0160 | 0x70 |
| 1 | 49 | 0061 | $0 \times 31$ | Q | 81 | 0121 | $0 \times 51$ | \| 9 | 113 | 0161 | 0x71 |
| 2 | 50 | 0062 | 0x32 | R | 82 | 0122 | $0 \times 52$ | r | 114 | 0162 | 0x72 |
| 3 | 51 | 0063 | $0 \times 33$ | S | 83 | 0123 | $0 \times 53$ | \| s | 115 | 0163 | 0x73 |
| 4 | 52 | 0064 | 0x34 | T | 84 | 0124 | $0 \times 54$ | t | 116 | 0164 | 0x74 |
| 5 | 53 | 0065 | 0x35 | U | 85 | 0125 | $0 \times 55$ | \| u | 117 | 0165 | 0x75 |
| 6 | 54 | 0066 | $0 \times 36$ | V | 86 | 0126 | $0 \times 56$ | \| v | 118 | 0166 | 0x76 |
| 7 | 55 | 0067 | $0 \times 37$ | W | 87 | 0127 | $0 \times 57$ | \| w | 119 | 0167 | 0x77 |
| 8 | 56 | 0070 | $0 \times 38$ | X | 88 | 0130 | $0 \times 58$ | \| x | 120 | 0170 | 0x78 |
| 9 | 57 | 0071 | 0x39 | Y | 89 | 0131 | $0 \times 59$ | \| y | 121 | 0171 | 0x79 |
| : | 58 | 0072 | 0x3a | Z | 90 | 0132 | 0x5a | \| | 122 | 0172 | $0 \times 7 \mathrm{a}$ |
| ; | 59 | 0073 | 0x3b | [ | 91 | 0133 | 0x5b | \| | 123 | 0173 | 0x7b |
| < | 60 | 0074 | 0x3c | 1 | 92 | 0134 | 0x5c | \| | | 124 | 0174 | 0x7c |
| $=$ | 61 | 0075 | 0x3d | ] | 93 | 0135 | 0x5d | \| \} | 125 | 0175 | 0x7d |
| > | 62 | 0076 | 0x3e | $\wedge$ | 94 | 0136 | 0x5e | \| ~ | 126 | 0176 | 0x7e |
| ? | 63 | 0077 | 0x3f | - | 95 | 0137 | 0x5f |  |  |  |  |

## x86 instructions

Machine code bytes

| B8 | 22 | 11 | 00 | FF |
| :--- | :--- | :--- | :--- | :--- |
| 01 | $C A$ |  |  |  |
| 31 | $F 6$ |  |  |  |
| 53 |  |  |  |  |
| 8B | $5 C$ | 24 | 04 |  |
| 8D | 34 | 48 |  |  |
| 39 | $C 3$ |  |  |  |
| 72 | EB |  |  |  |
| C3 |  |  |  |  |

## Assembly

foo:
movl \$0xFF001122, \%eax addl \%ecx, \%edx
xorl \%es1, \%es1
pushl \%ebx movl 4(\%esp), \%ebx
leal (\%eax,\%ecx,2), \%exi
cmpl \%eax, \%ebx
jnae foo
retl

## Bits and Bytes Require Interpretation

00000000001101010011000000110001 (or 0x00353031) might be interpreted as

- The integer 3,485,74510
- A floating point number close to $4.884569 \times 10^{-39}$
- The string " 105 "
- A portion of an image or video
- An address in memory


## Information is Bits + Context

```
#include<stdio.h>
int main(int argc, char** argv) {
    printf("Hello world!\n");
    return 0;
}
```


## Example Data Representations

| C Data Type |  |
| :--- | :---: |
| char | x86-64 |
| short | 1 |
| int | 2 |
| long | 4 |
| float | 8 |
| double | 4 |
| pointer | 8 |

## Memory Access in C

```
int x; // an integer
int *p; // a pointer to an integer
// normal initialization:
x = 0;
// silly, but illustrative:
p = &x; // & means "address of"
*p = 0; // * means "memory at address"
```

- \& and * are inverses of one another
- prefix vs infix operators
- x occupies 4 bytes in memory; p occupies 8


## Arrays

- Contiguous block of memory
- Pointer to start, then indexed by element size
- Indices start at zero
- ary [k] is the same as * (ary+k)
- Location of ary +k depends on the type of array elements


## Arrays and Pointers Combined

int *p[47];

- Array of pointers ... or ... pointer to an array??
- It's an array of 47 pointers
- p [3] is the fourth pointer in the array p
- $\mathrm{p}[3]$ is the base of an array
- p [3] [6] is the integer at position 6 in the array $p$ [3]


## What is printed?

```
int a[100];
int *p[47];
p[3] = a+12;
for (int i = 0; i < 100; i++){
    a[i] = i;
}
printf("%d\n", p[3][4]);
```


## Structs

- Heterogeneous records, like objects
- Typical linked list declaration:

```
typedef struct cell {
    int value;
    struct cell *next;
} cell_t;
```

- Usage:

```
cell_t c;
c.value = 42;
c.next = NULL;
```

How many bytes are allocated for $c$ ?
for $p$ ?

- Usage with pointers:

```
cell_t *p;
p->value = 42;
p->next = NULL;
```

$\mathrm{p}->$ next is an
bbreviation for
(*p). next

## Compilation

## - gcc -o hello hello.c



```
#include<stdio.h>
int main(int argc,
    char ** argv) {
    printf("Hello
            world!\n");
    return 0;
}
```

```
int printf(const char *
    restrict,
    ...)
__(__printf__,}(\frac{_}{1,}2)))
int main(int argc,
        char ** argv) {
    printf("Hello
    world!\n");
    return 0;
}
```

```
55
48 89 e5
48 83 ec 20
48 8d 05 25 00 00 00
c7 45 fc 00 00 00 00
89 7d f8
48 89 75 f0
48 89 c7
b0 00
e8 00 00 00 00
31 c9
89 45 ec
89 c8
48 83 c4 20
5d
c3
```


## Running a Program

- ./hello


## A Computer System



## LOGISTICS

## Course staff



## Prof. Eleanor Birrell

Edmunds 221
Research in security and privacy
OH: T 1:30-3:30pm, W 7-9pm


Jenna
Brandt


Joe
Brennan


Gabriel da Motta


Adam Lininger-White


Douglas Webster

## The Course in a Nutshell

- Textbook
- Bryant and O'Halloran, Computer Systems: A Programmer's Perspective, third edition, Pearson, 2016 (Recommended)
- Classes
- Monday and Wednesday, 1:15-2:30 or 2:45-4pm in Edmunds 101
- Labs
- Mondays 7-8:15 in Edmunds 229/219
- Start Monday! Be sure to have an account and password

Mentor Session Schedule (Edmunds 227)

| Monday | Tuesday | Wednesday | Thursday | Friday | Saturday | Sunday |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| LAB | $7-9 p m$ | $7-9 p m$ | $7-9 p m$ | $2-4 p m$ | $2-4 p m$ | $2-4 \mathrm{pm}$ |

## Grading

- Assignments
- Introduced during labs, Due Fridays at 5pm
- Tremendous fun, work in pairs
- $45 \%$ of the grade
- Seven late days
- Midterm exam
- March 11
- $20 \%$ of the grade each
- Final exam
- Thursday, May 7 or Tuesday, May 12 or Friday, May 15 (2-5pm)
- $30 \%$ of the grade
- Participation
- $5 \%$ of the grade


## Course website

## https://www.cs.pomona.edu/classes/cs105

- All information is on the course website
- All course materials get posted on the course website
- Links from the course page:
- Piazza, for questions and discussion
- Lab assistants and mentors, schedule
- Submission site
- Grading site

