

# Lecture 9: Use and Abuse of the Stack (cont'd)

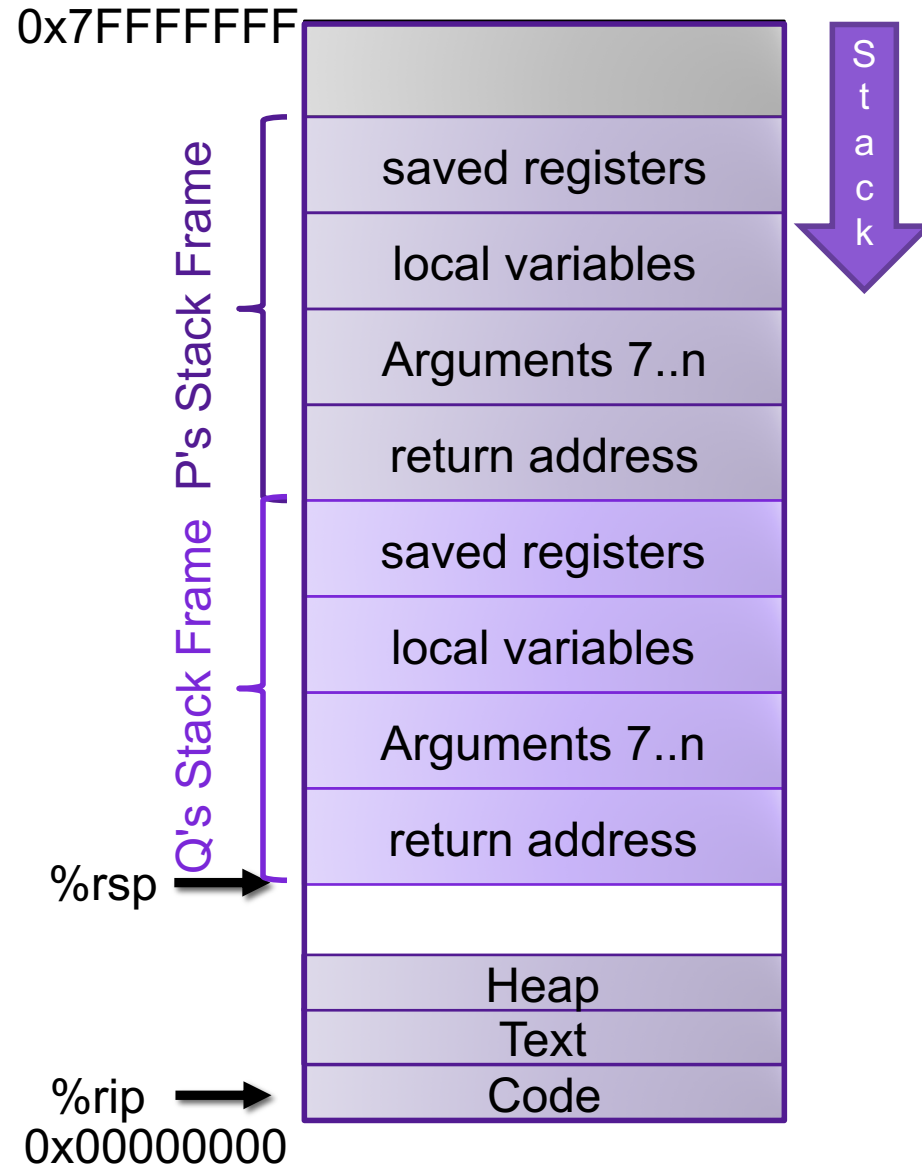
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CS 105

September 26, 2019

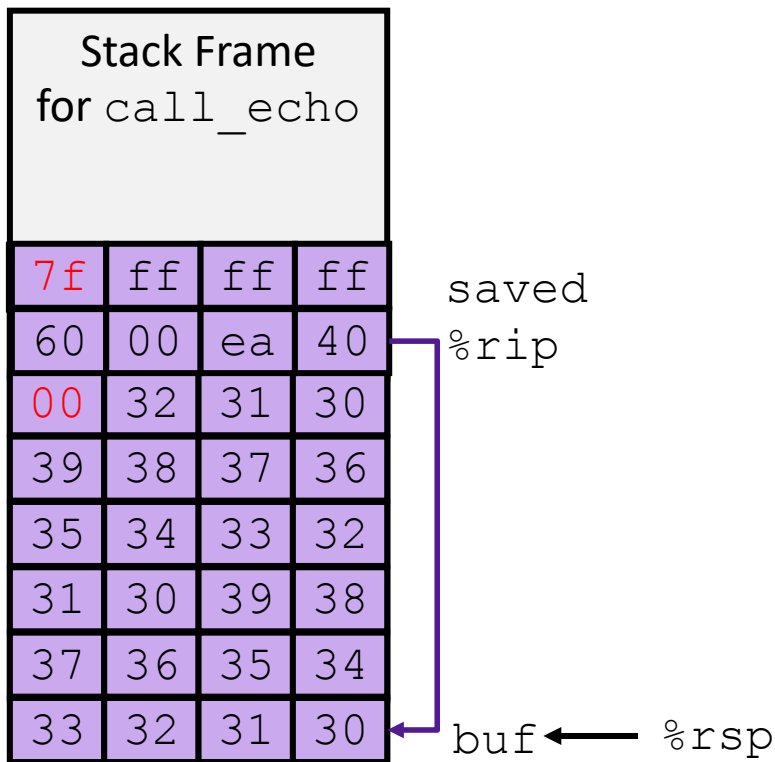
# Review: Stack Frames

- Each function called gets a stack frame
- Passing data:
  - calling procedure P uses registers (and stack) to provide parameters to Q.
  - Q uses register `%rax` for return value
- Passing control:
  - **call <proc>**
    - Pushes return address (current `%rip`) onto stack
    - Sets `%rip` to first instruction of proc
  - **ret**
    - Pops return address from stack and places it in `%rip`
- Local storage:
  - allocate space on the stack by decrementing stack pointer, deallocate by incrementing



# Review: Buffer Overflow

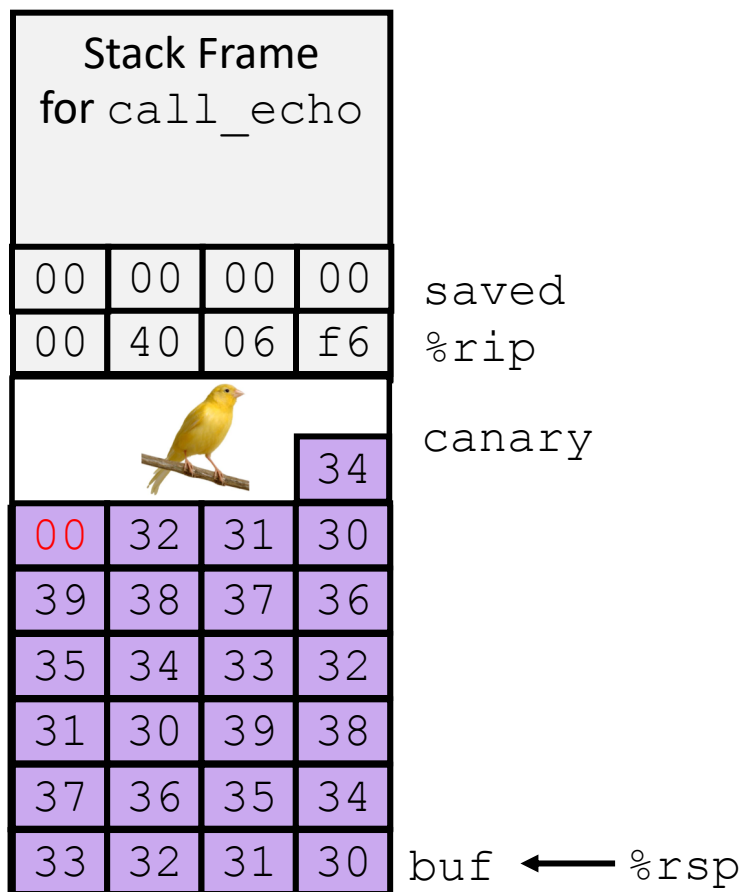
- Most common form of memory reference bug
  - Unchecked lengths on string inputs
  - Particularly for bounded character arrays on the stack
    - sometimes referred to as stack smashing



```
/* Echo Line */
void echo()
{
    char buf[4];
    gets(buf);
    puts(buf);
}
```

```
echo:
    subq   $18, %rsp
    movq   %rsp, %rdi
    call   gets
    call   puts
    addq   $18, %rsp
    ret
```

# Review: Stack Canaries



```

/* Echo Line */
void echo()
{
    char buf[4];
    gets(buf);
    puts(buf);
}

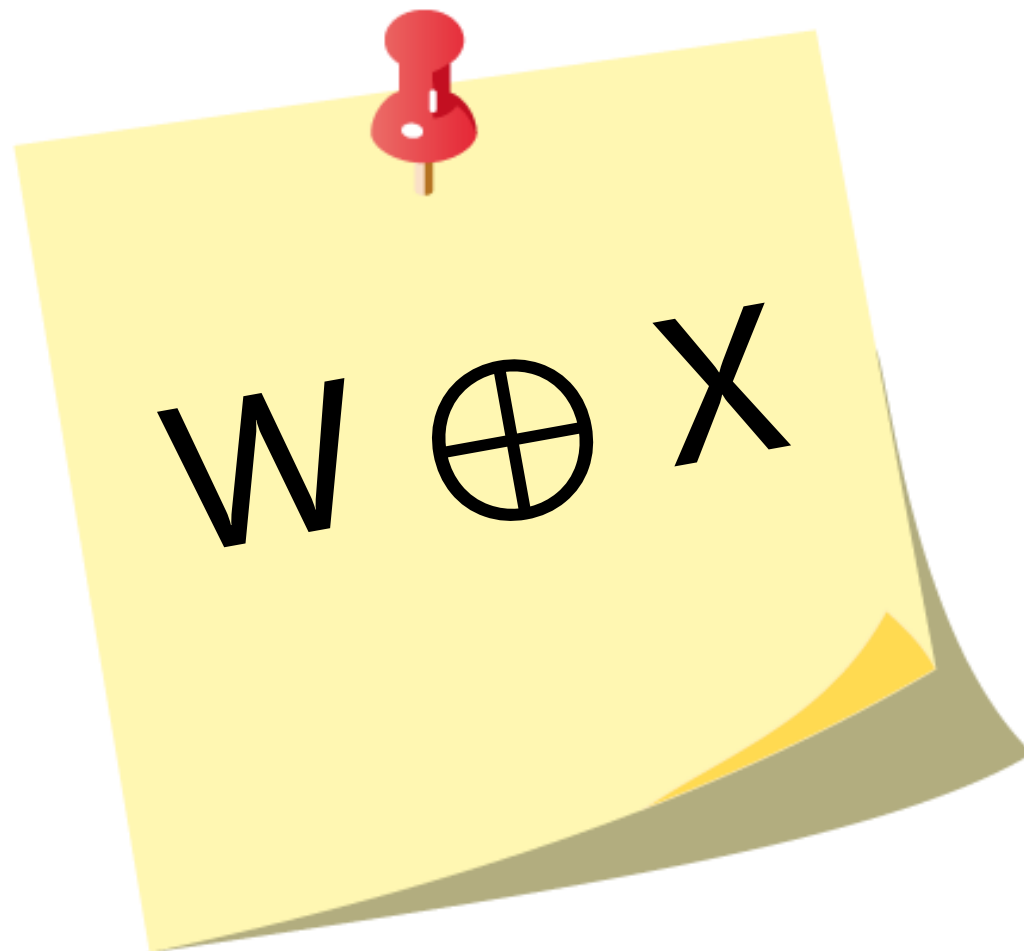
```

```

echo:
    subq    $24, %rsp
    movq   %rsp, %rdi
    call   gets
    call   puts
    movq   24(%rsp), %rdx
    xorq   %fs:40, %rdx
    je     .L3
    call   __stack_chk_fail
.L3:
    addq   $24, %rsp
    ret

```

# Review: Memory Tagging



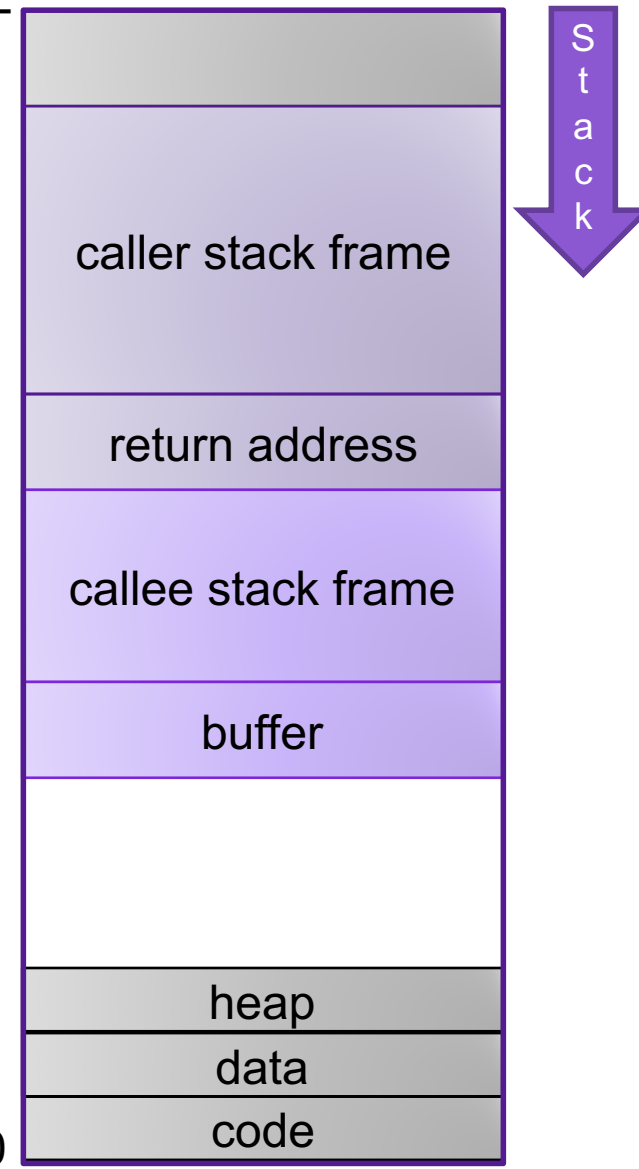
# Code Reuse Attacks

- Key idea: execute instructions that already exist
- Defeats memory tagging defenses
- Examples:
  1. return to a function in the current program
  2. return to a library function (e.g., return-into-libc)
  3. return to some other instruction (return-oriented programming)

# Returning to a function

0x7FFFFFFF

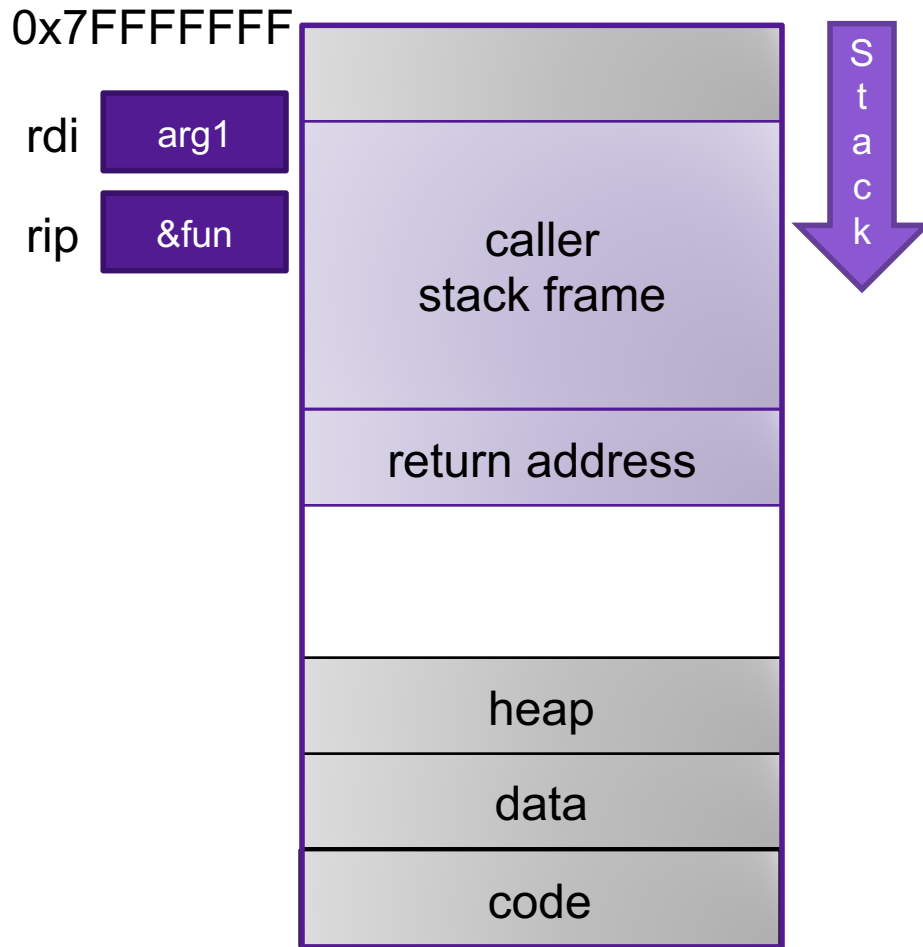
- Overwrite the saved return address with the location of a function in the current program



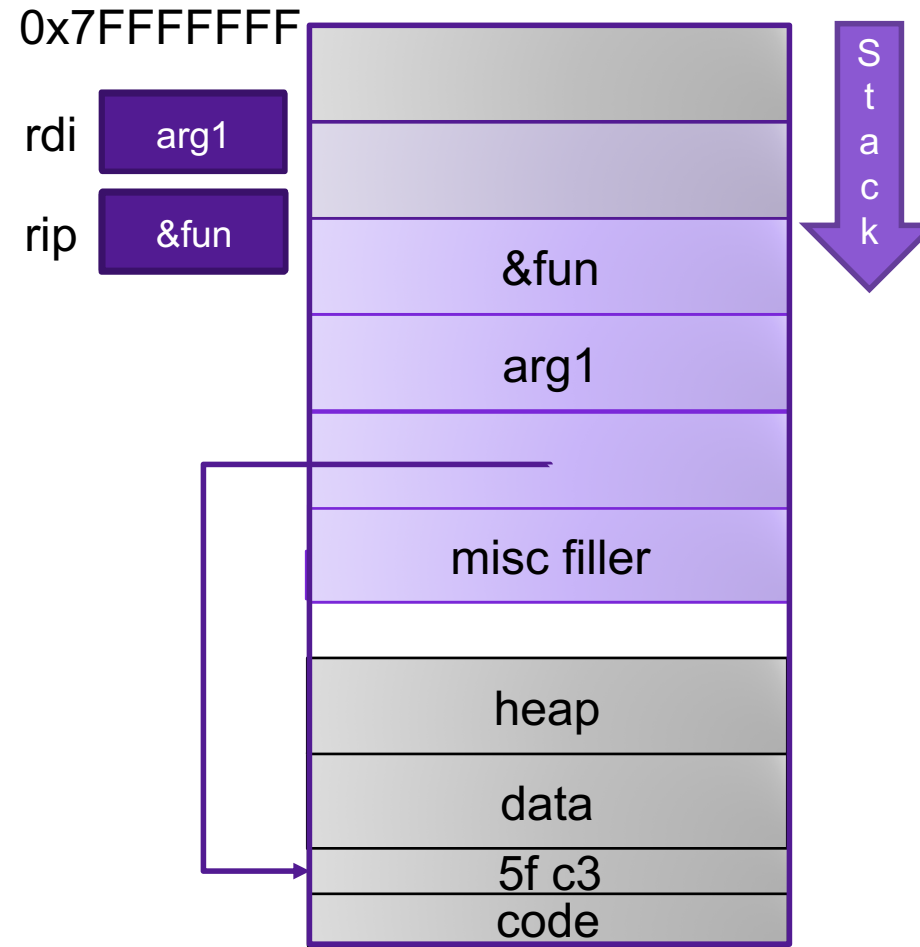
0x00000000

# Handling Arguments

what function expects  
when it is called...



overflow with argument





# Return-into-libc

Sr.No.	Function & Description
1	<b>double atof(const char *str)</b> <a href="#">↗</a> Converts the string pointed to, by the argument <i>str</i> to a floating-point number (type double).
2	<b>int atoi(const char *str)</b> <a href="#">↗</a> Converts the string pointed to, by the argument <i>str</i> to an integer (type int).
3	<b>long int atol(const char *str)</b> <a href="#">↗</a> Converts the string pointed to, by the argument <i>str</i> to a long integer (type long int).
8	<b>void free(void *ptr)</b> <a href="#">↗</a> Deallocates the memory previously allocated by a call to <i>calloc</i> , <i>malloc</i> , or <i>realloc</i> .
9	<b>void *malloc(size_t size)</b> <a href="#">↗</a> Allocates the requested memory and returns a pointer to it.
10	<b>void *realloc(void *ptr, size_t size)</b> <a href="#">↗</a> Attempts to resize the memory block pointed to by <i>ptr</i> that was previously allocated with a call to <i>malloc</i> or <i>calloc</i> .
15	<b>int system(const char *string)</b> <a href="#">↗</a> The command specified by <i>string</i> is passed to the host environment to be executed by the command processor.
16	<b>void *bsearch(const void *key, const void *base, size_t nitems, size_t size, int (*compar)(const void *, const void *))</b> <a href="#">↗</a> Performs a binary search.
17	<b>void qsort(void *base, size_t nitems, size_t size, int (*compar)(const void *, const void*))</b> <a href="#">↗</a> Sorts an array.
18	<b>int abs(int x)</b> <a href="#">↗</a> Returns the absolute value of <i>x</i> .
22	<b>int rand(void)</b> <a href="#">↗</a> Returns a pseudo-random number in the range of 0 to <i>RAND_MAX</i> .
23	<b>void srand(unsigned int seed)</b> <a href="#">↗</a> This function seeds the random number generator used by the function <b>rand</b> .

# ASCII Armoring

- Make sure all system library addresses contain a null byte (0x00).
- Can be done by placing this code in the first 0x01010101 bytes of memory

# Properties of x86 Assembly

- variable length instructions
- not word aligned
- dense instruction set

# Gadgets

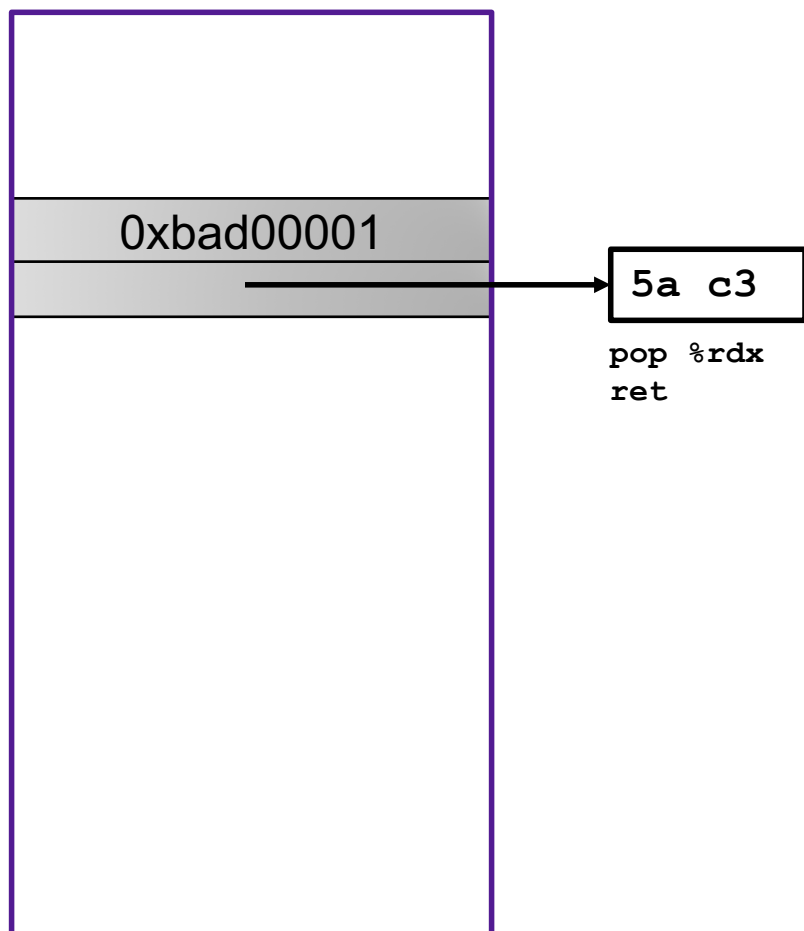
```
void setval(unsigned *p) {  
    *p = 3347663060u;  
}
```

```
<setval>:  
4004d9: c7 07 d4 48 89 c7 movl $0xc78948d4, (%rdi)  
4004df: c3                ret
```

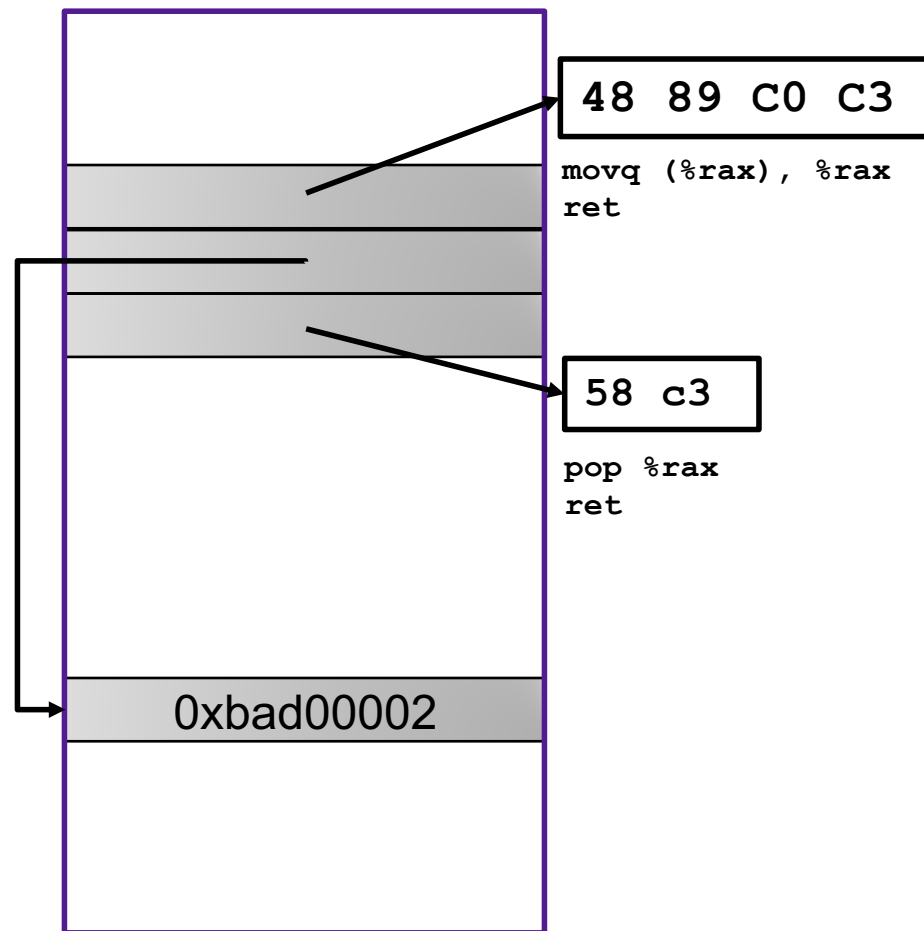
gadget address: 0x4004dc  
encodes: movq %rax, %rdi  
ret  
executes: %rdi <- %rax

# Example Gadgets

Load Constant



Load from memory

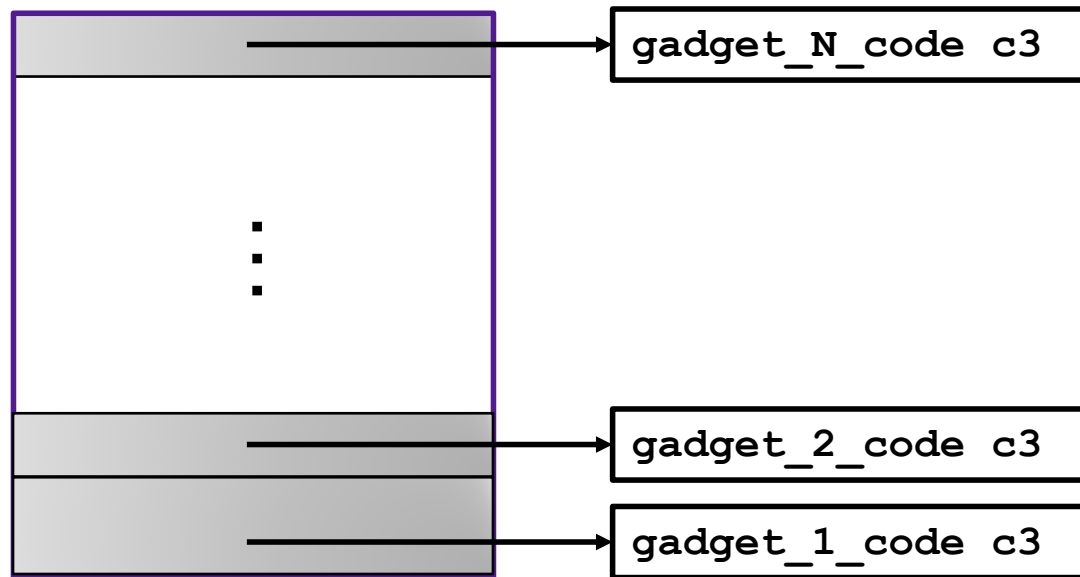


# Return-oriented Programming

Return-Oriented  
Programming

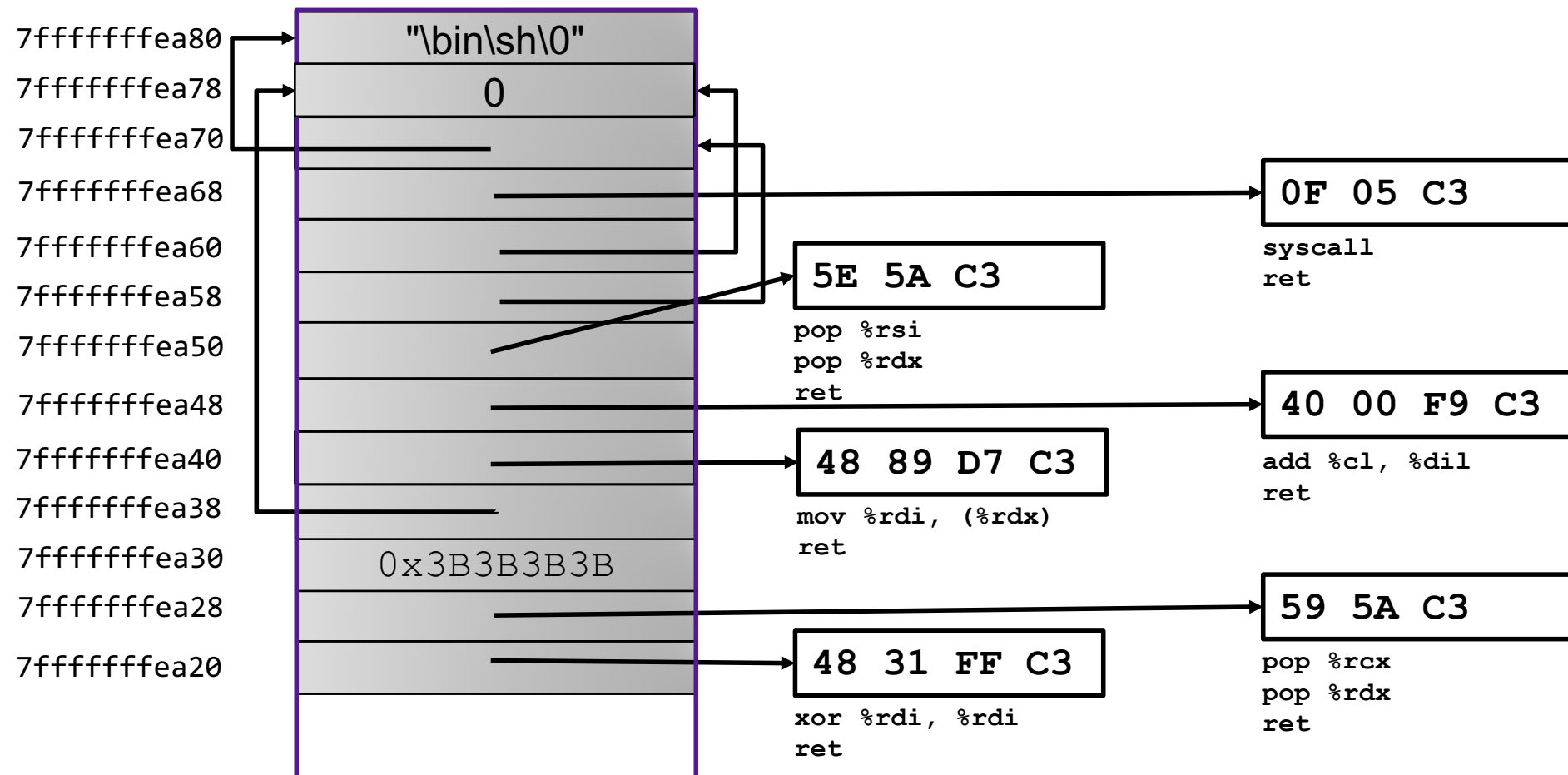
is a lot like a ransom  
note, BUT instead of cutting  
out letters from magazines,  
YOU ARE cutting out  
instructions from text  
segments

# Return-oriented Programming



Final ret in each gadget sets pc (%rip) to beginning of next gadget code

# Return-Oriented Shellcode





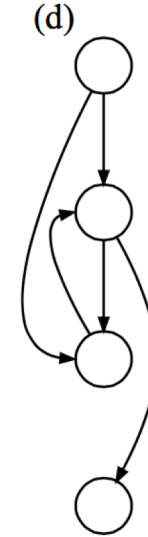
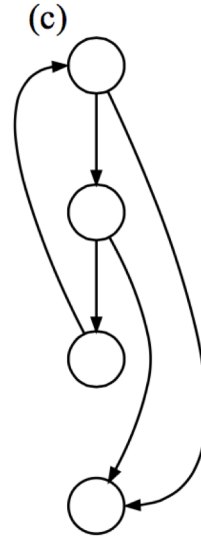
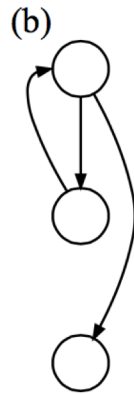
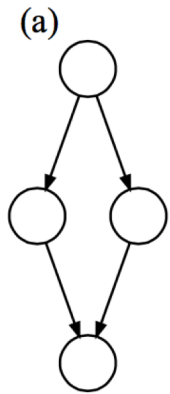
# Address Space Layout Randomization



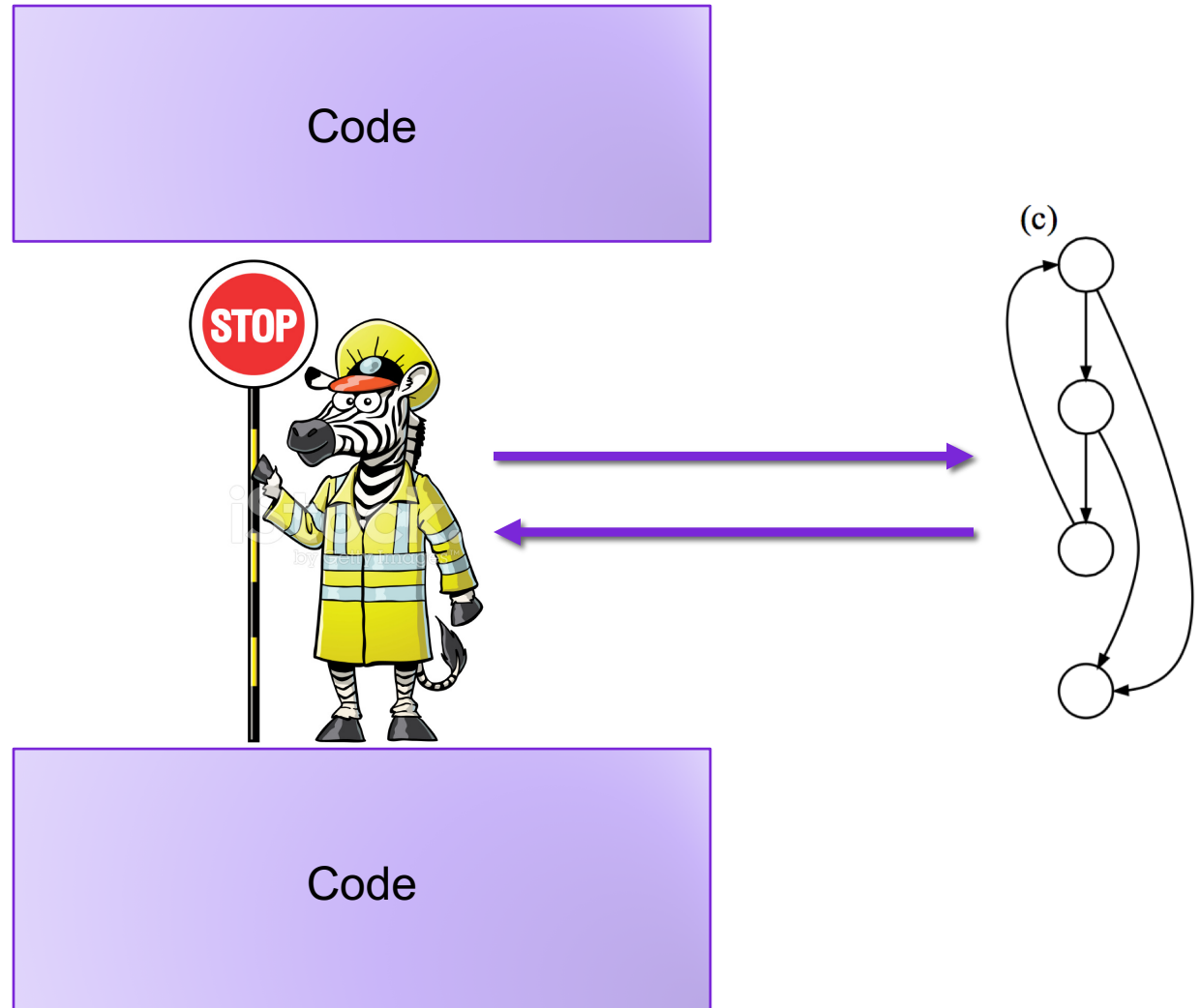
# Gadget Elimination



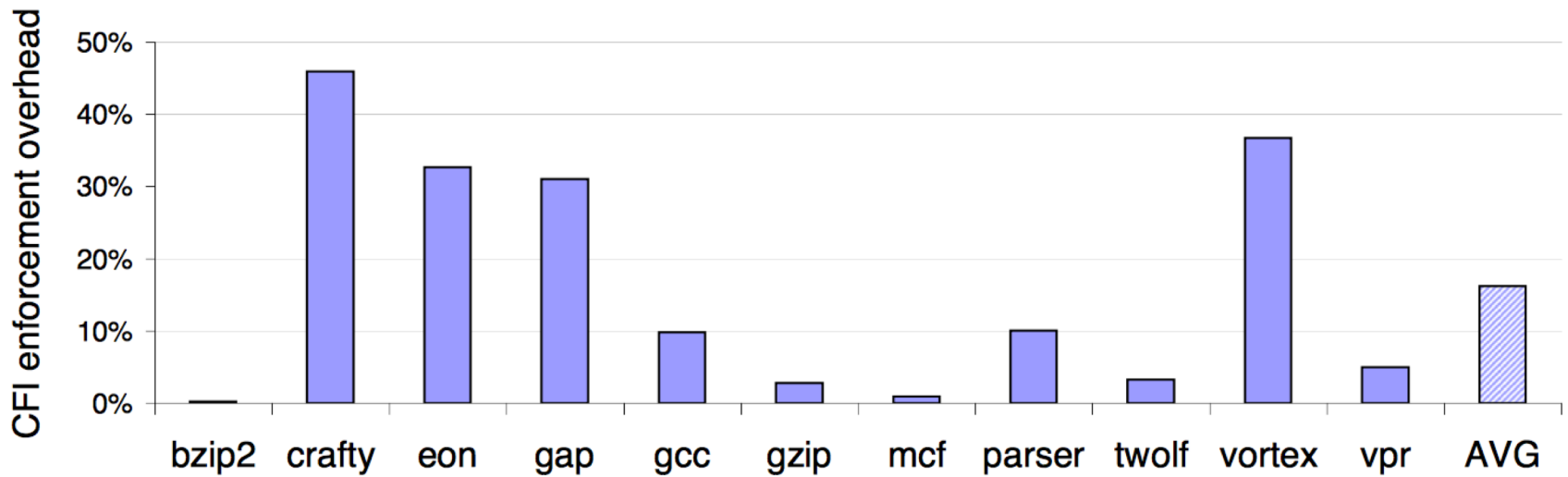
# Control Flow Integrity



# CFI = Insert Monitors

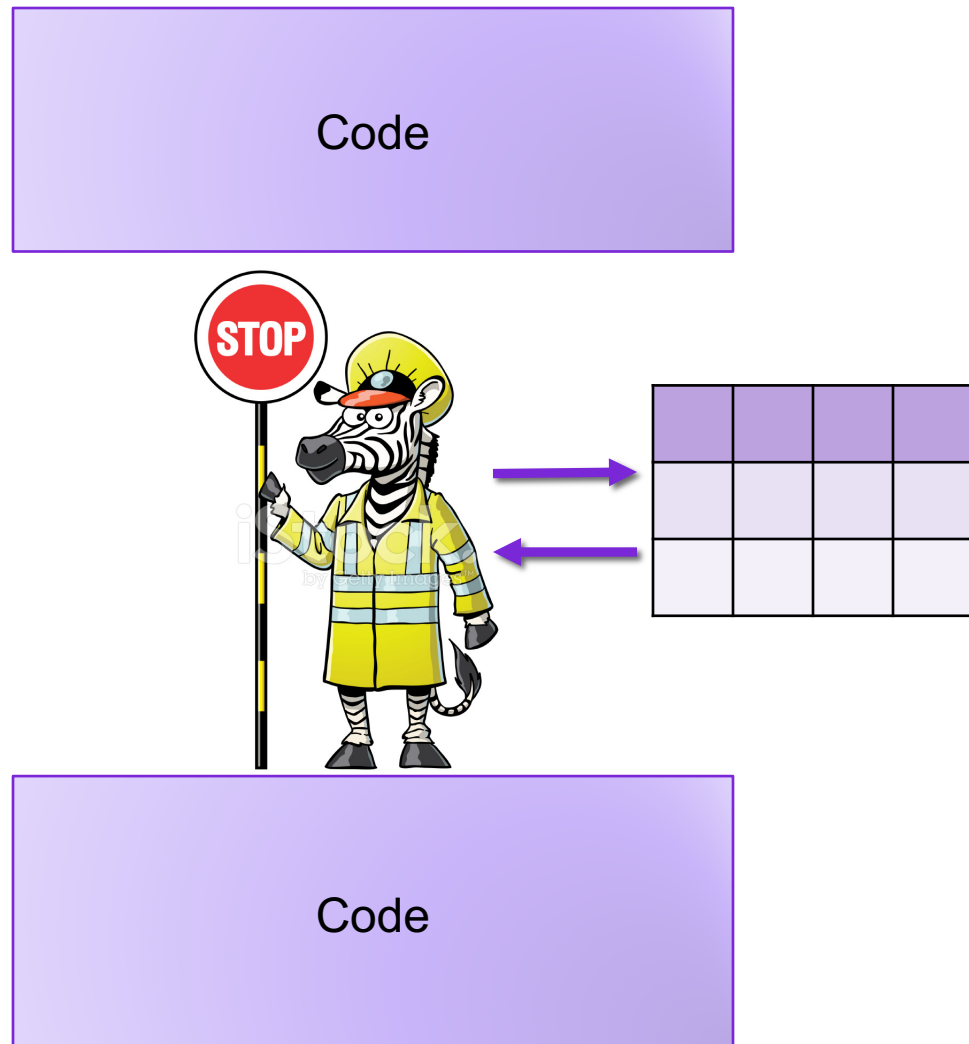


# CFI Overhead



# Control Flow Guard

- Approximate CFI implementation in Windows 8.1, 10
- Jump is valid if it begins at the beginning of a function
  - Granularity: 8 bytes
- Check implemented as a bitmap



# The state of the world

## Defenses:

- high-level languages
- Stack Canaries
- Memory tagging
- ASLR
- continuing research and development...

But all they aren't perfect!



# The state of the world



*"Security is lax on this side."*

CN  
COLLECTION