

CS52 RECURSION

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CS 52 – Spring 2017

### Admin

- Final grade
- Assignment grading
- Assignment 4
- Reading

### Examples from this lecture

<http://www.cs.pomona.edu/~dkauchak/classes/cs52/examples/cs52machine/>

### CS52 machine

**CPU**

processor

registers

ic instruction counter (location in memory of the next instruction in memory)

r0 holds the value 0 (read only)

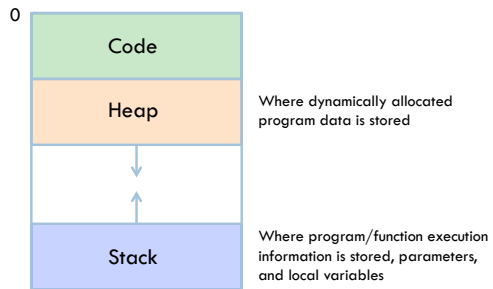
r1

r2

r3

- general purpose  
- read/write

## Memory layout



## Stack frame

Key unit for keeping track of a function call

- return address (where to go when we're done executing)
- parameters
- local variables

## CS52 function call conventions

r1: reserved for the stack pointer

r2: contains the return address

r3: contains the first parameter

additional parameters go on the stack (more on this)

the result (i.e. the return value) should go in r3

## Sum revisited

```

fun sum x =
  if x <= 0 then
    0
  else
    x + sum (x-1);

```

Note to future Dave from past Dave: write the function up on the board ☺

```

sum
  psh r2          ; save the return address on the stack
  bgt r3 r0 recurse ; check base case
  add r3 r0 0     ; if x <= 0, result is 0
  brs done

recurse
  psh r3          ; save n on the stack
  sub r3 r3 1    ; x = x-1

  lcw r2 sum     ; make recursive call
  cal r2 r2      ; sum (x-1), answer should be in r3

  pop r2         ; get n into r2
  add r3 r3 r2   ; r3 = n + sum (x-1)

done
  pop r2         ; get the return address
  jmp r2        ; go back to where we were called from
  
```

Function startup  
base cases

recursive call

recursive case

answer calculation

Function cleanup and return

```

sum
  psh r2          ; save the return address on the stack
  bgt r3 r0 recurse ; check base case
  add r3 r0 0     ; if n <= 0, result is 0
  brs done

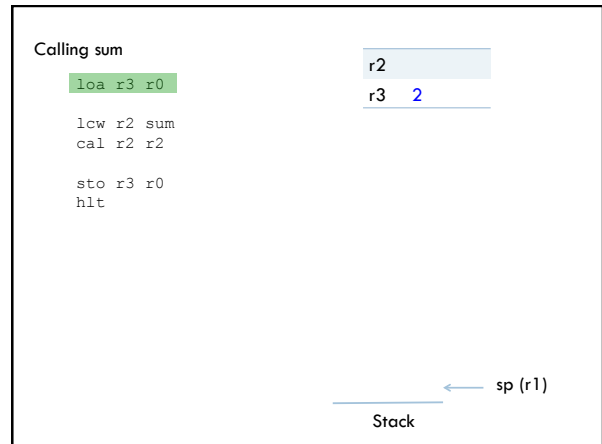
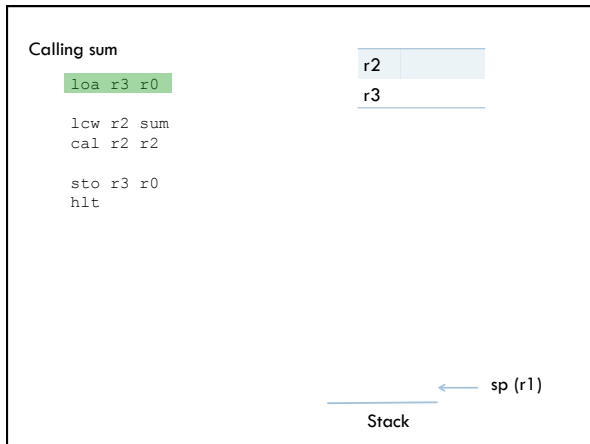
recurse
  psh r3          ; save n on the stack
  sub r3 r3 1    ; n = n-1

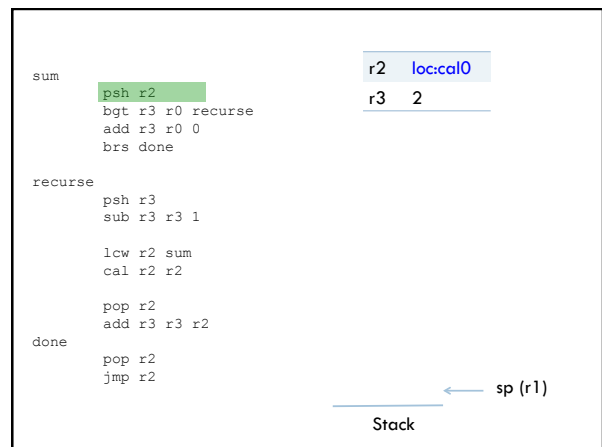
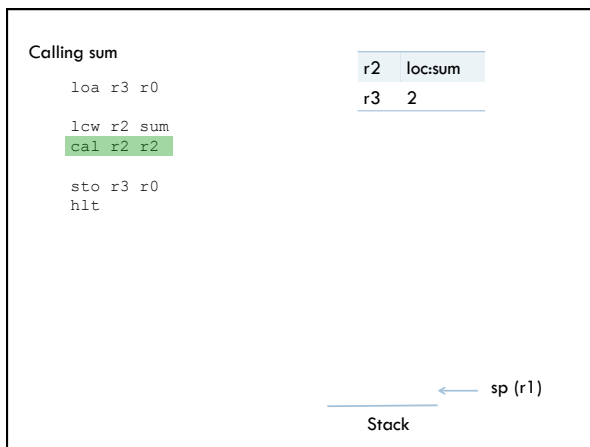
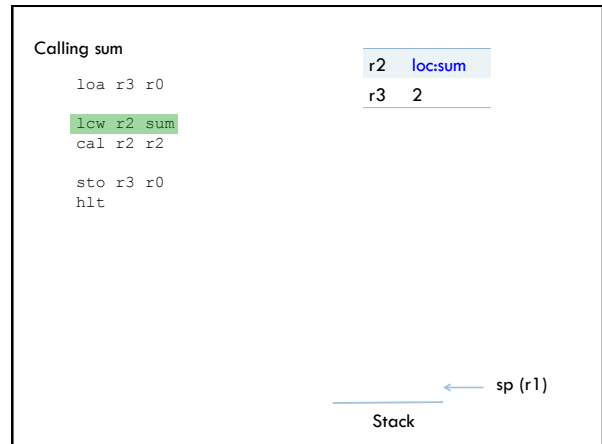
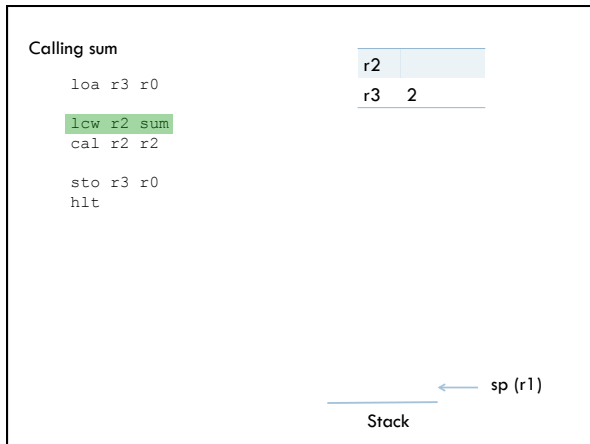
  lcw r2 sum     ; make recursive call
  cal r2 r2      ; sum(n-1), answer should be in r3

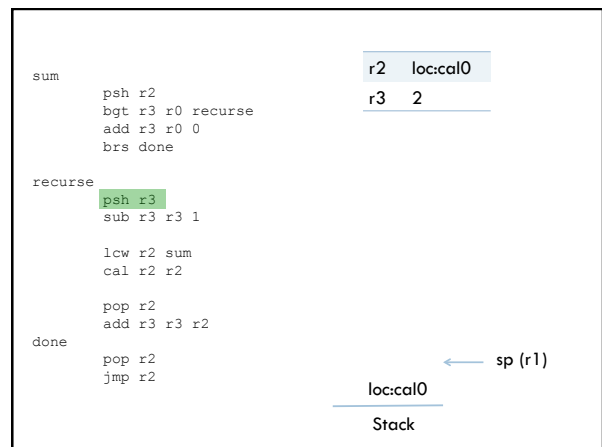
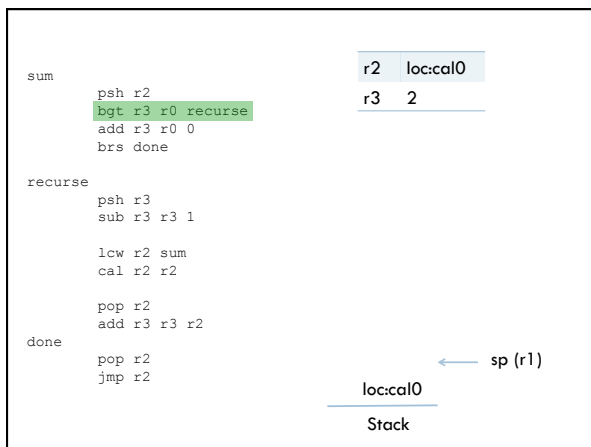
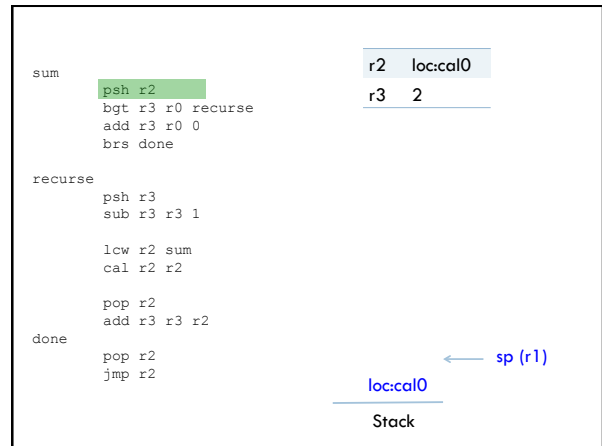
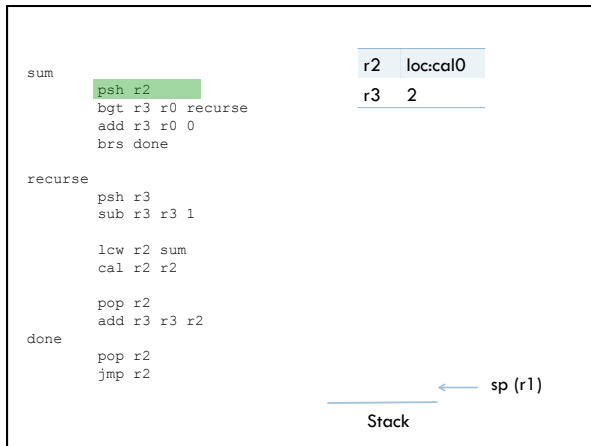
  pop r2         ; get n into r2
  add r3 r3 r2   ; r3 = n + sum(n-1)

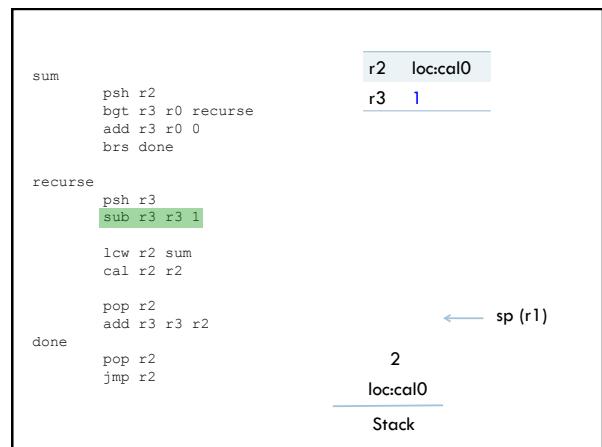
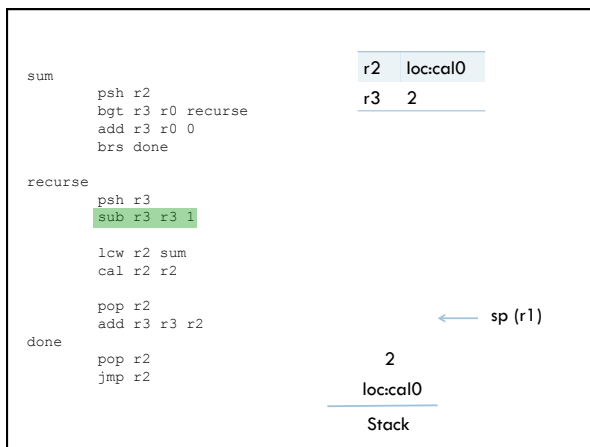
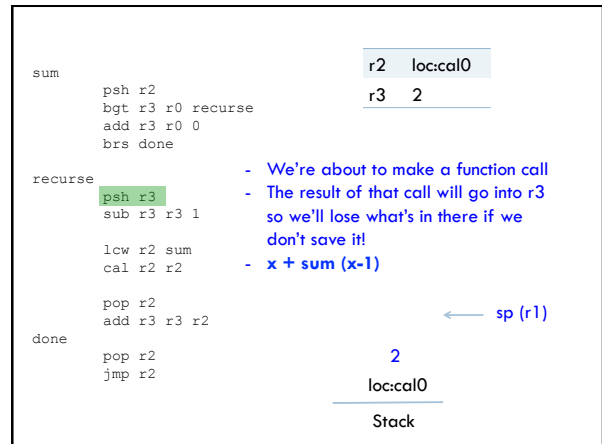
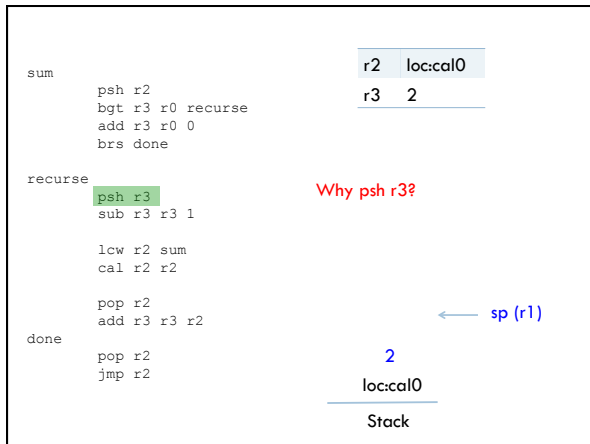
done
  pop r2         ; get the return address
  jmp r2        ; go back to where we were called from
  
```

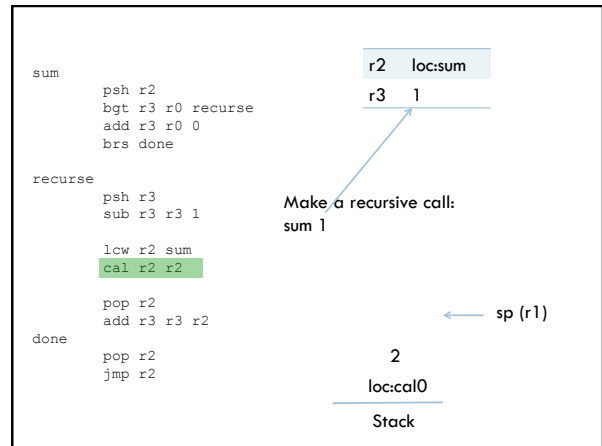
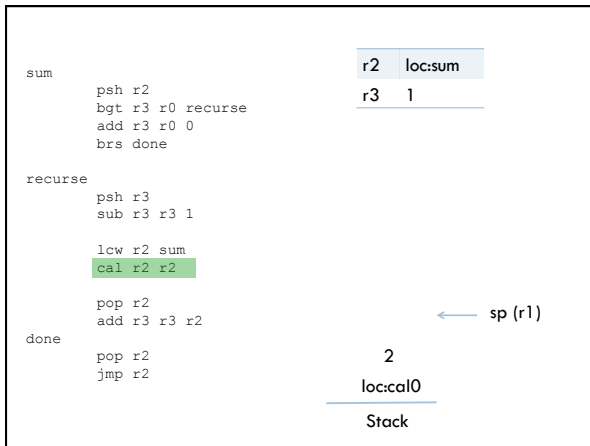
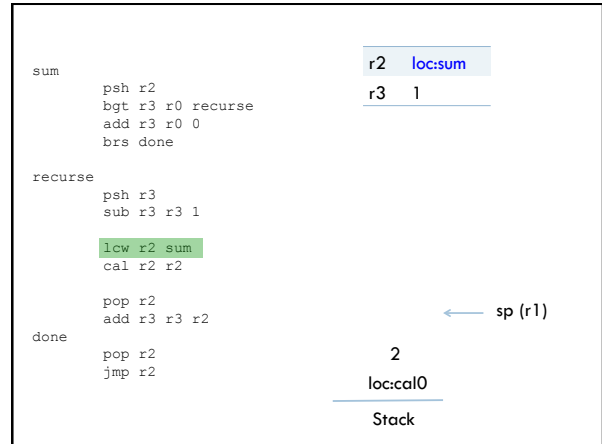
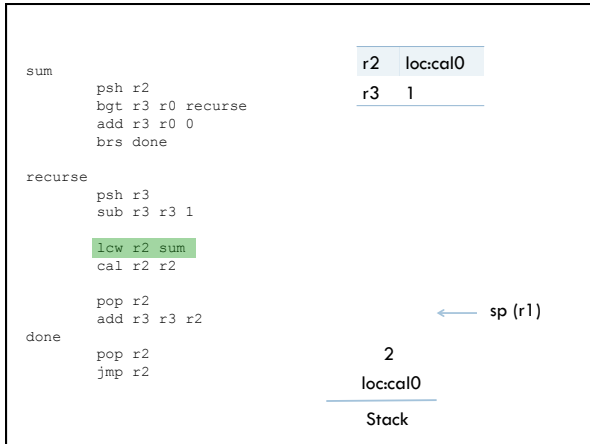
Notice symmetry of psh and pop

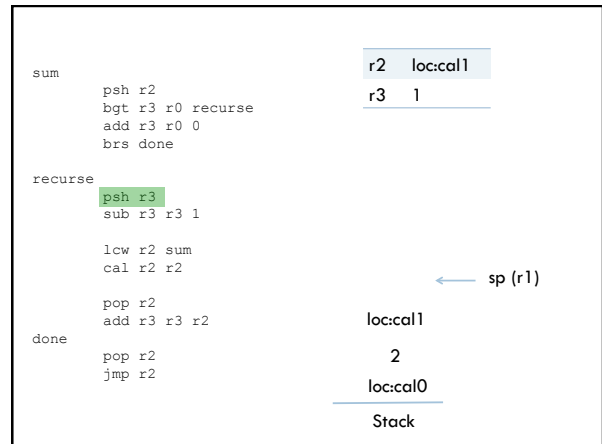
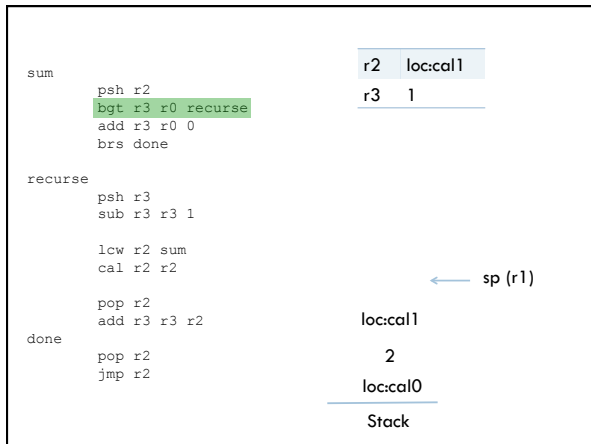
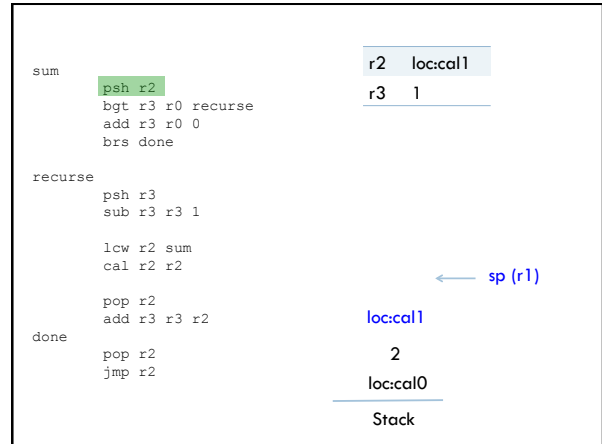
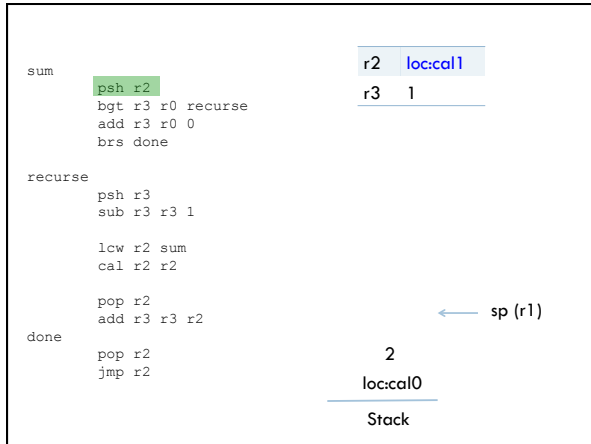




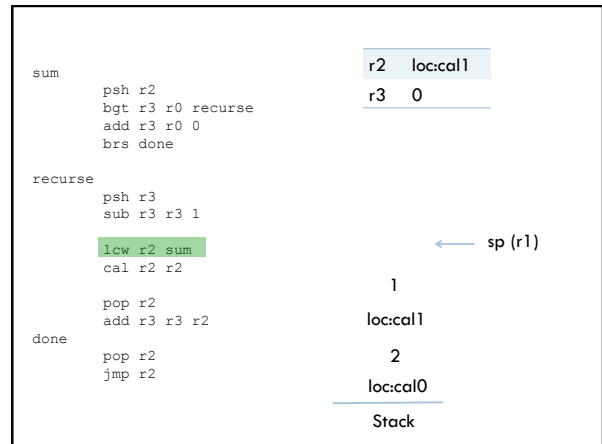
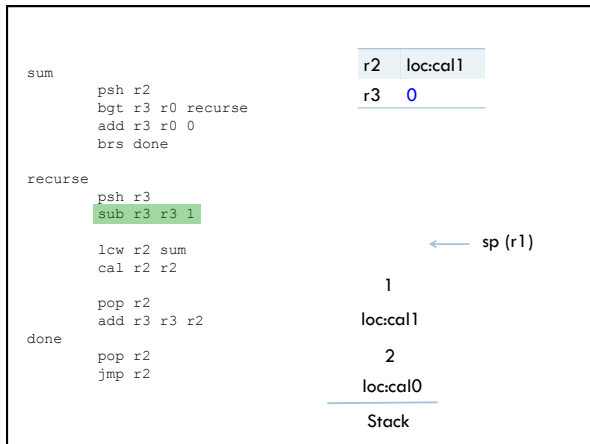
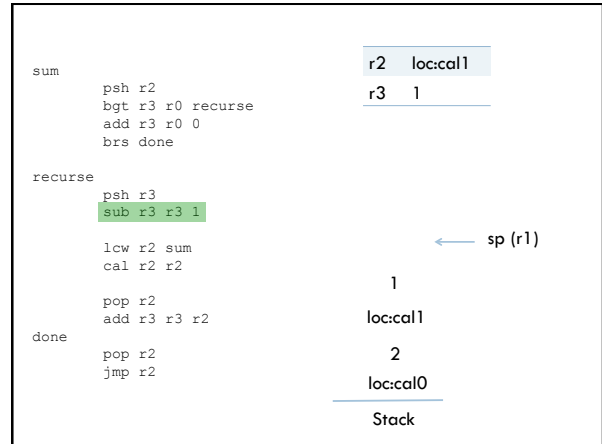
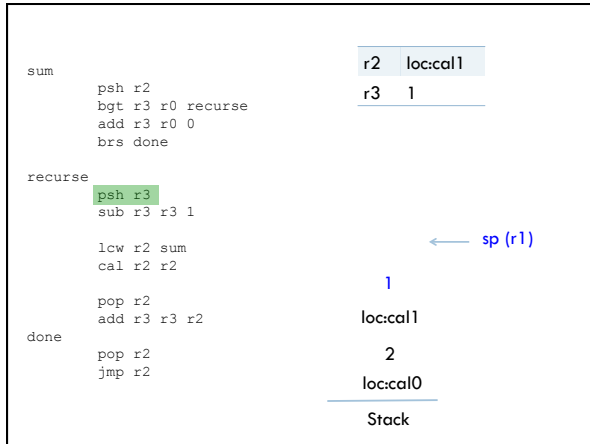


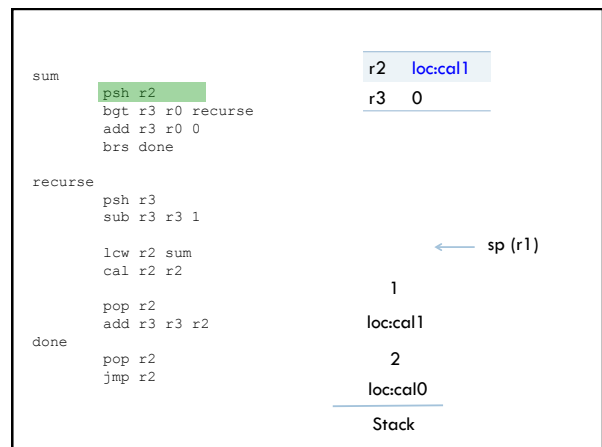
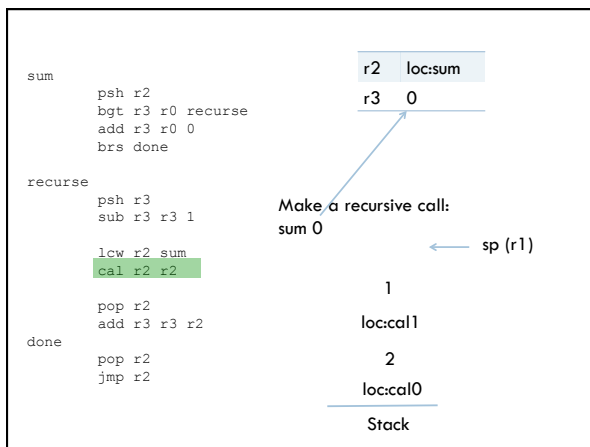
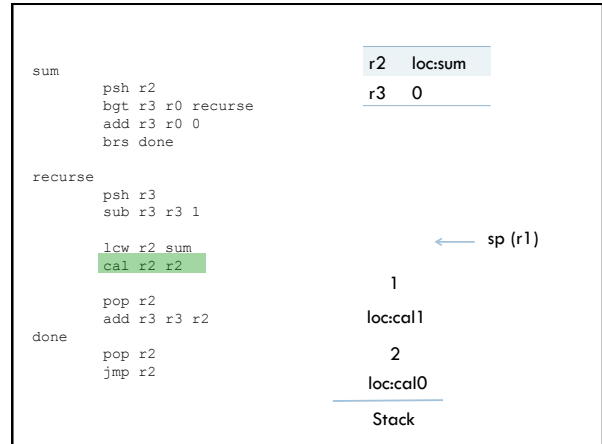
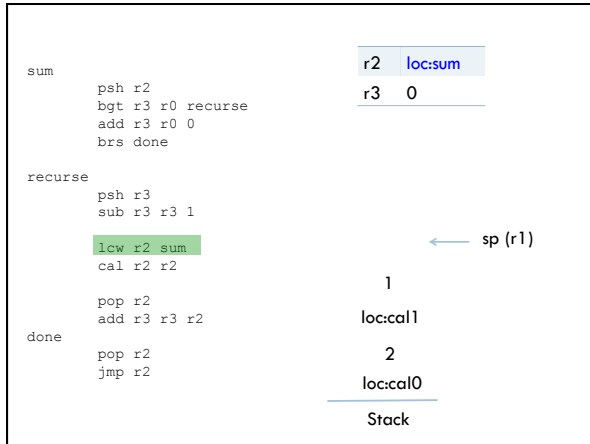


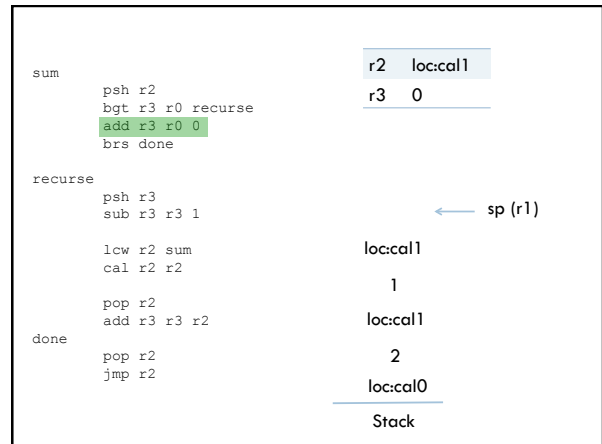
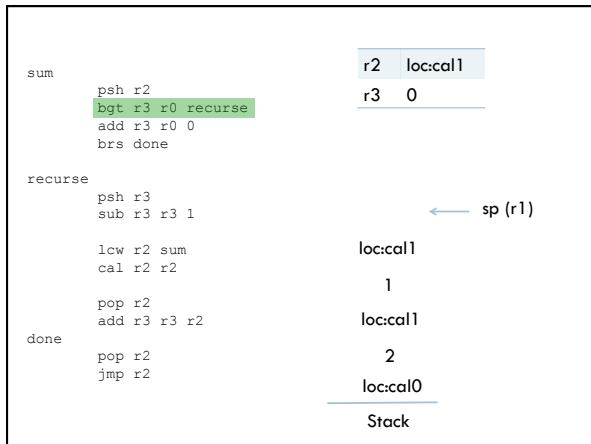
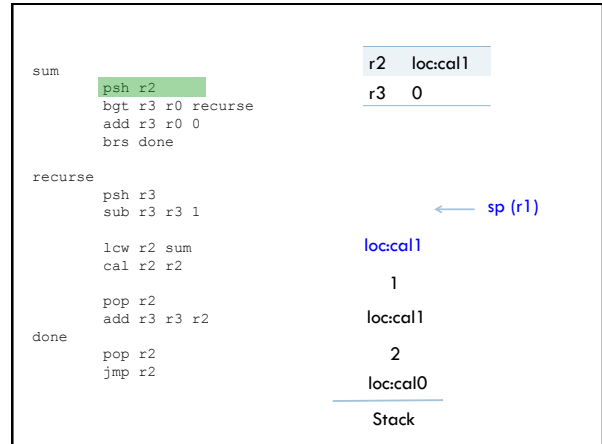
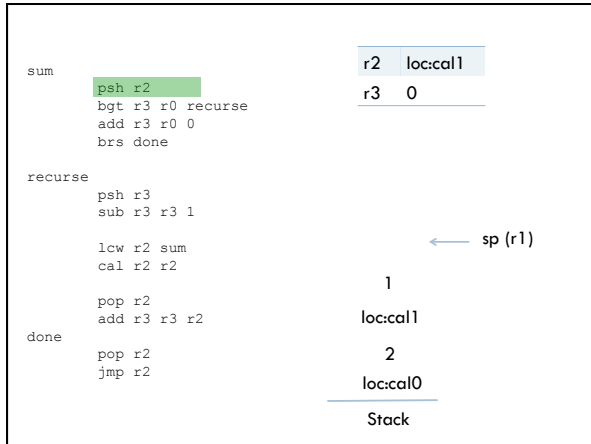


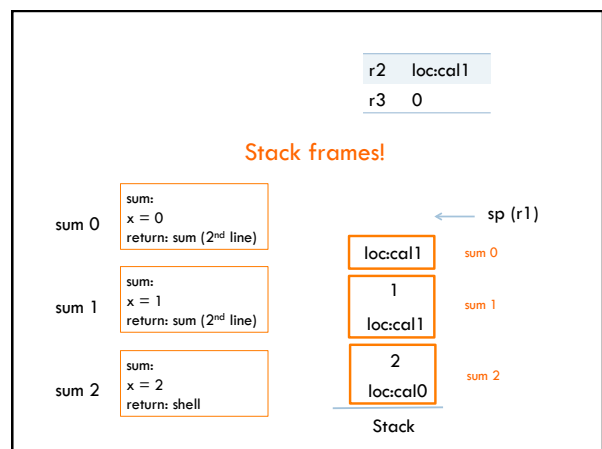
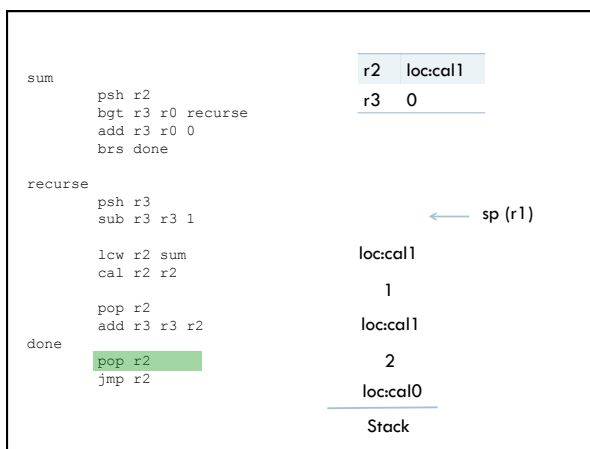
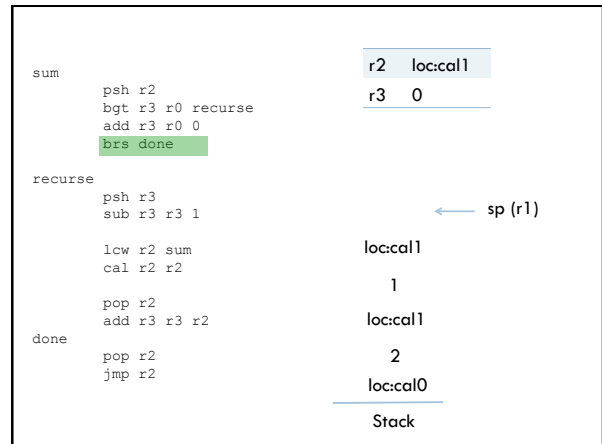
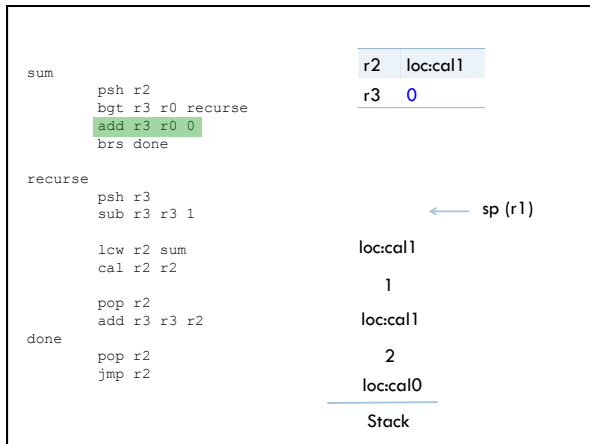


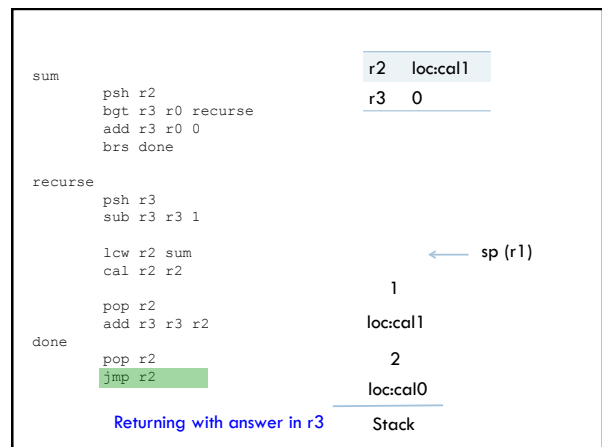
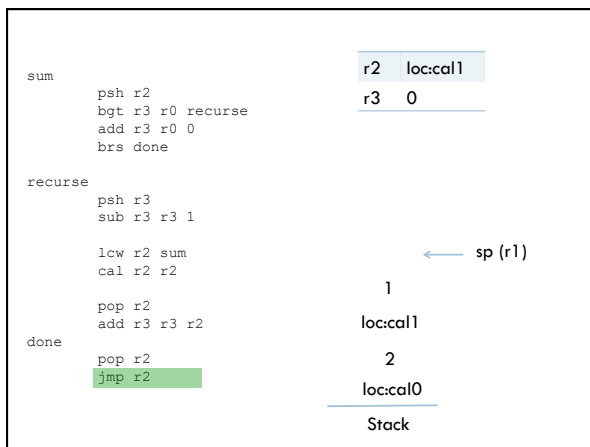
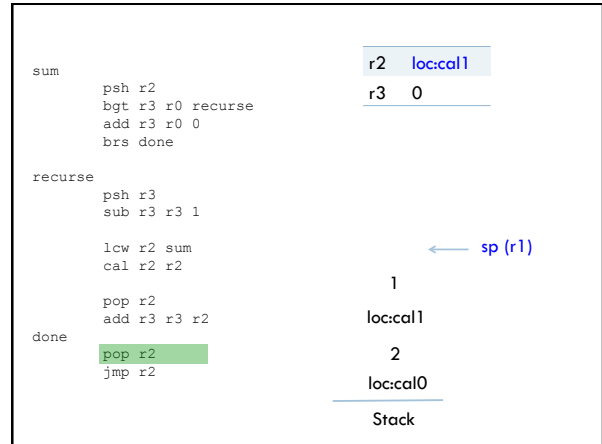
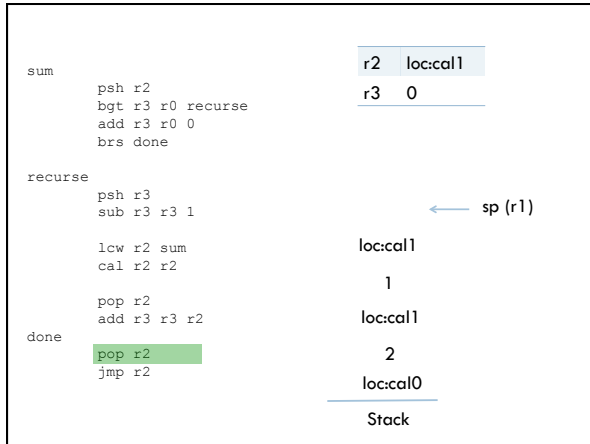


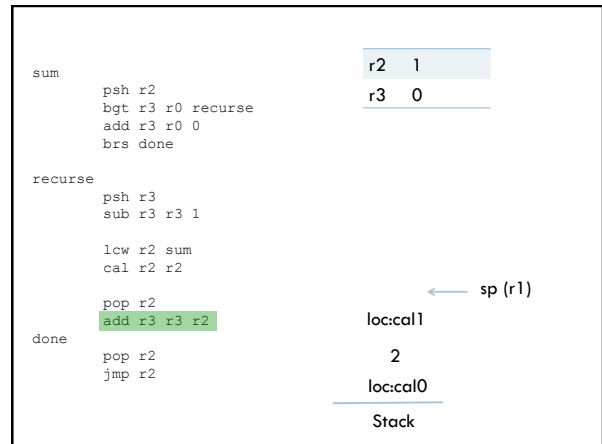
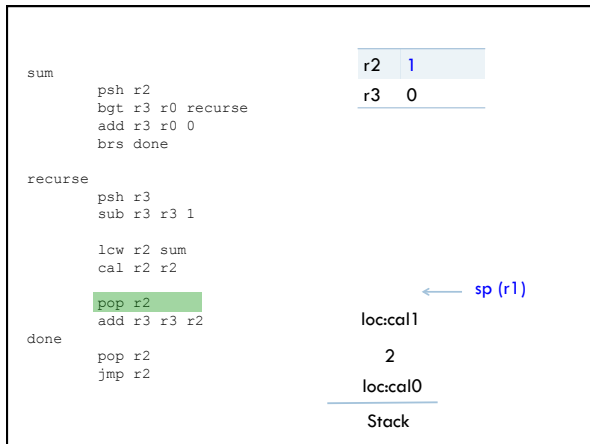
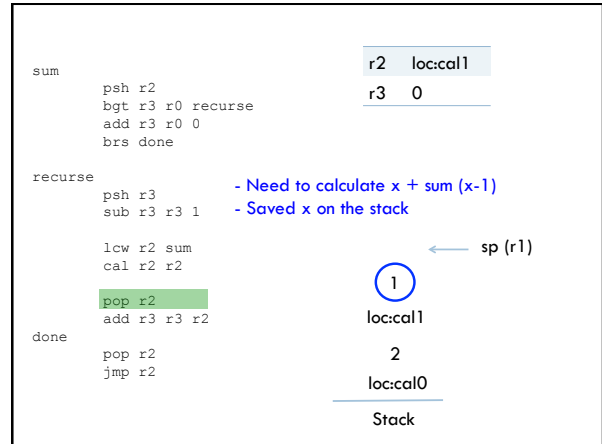
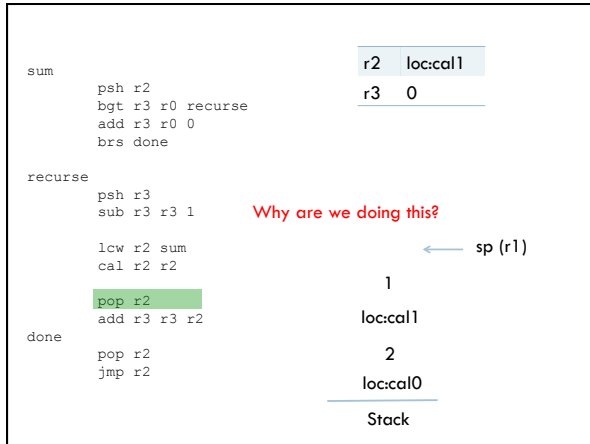


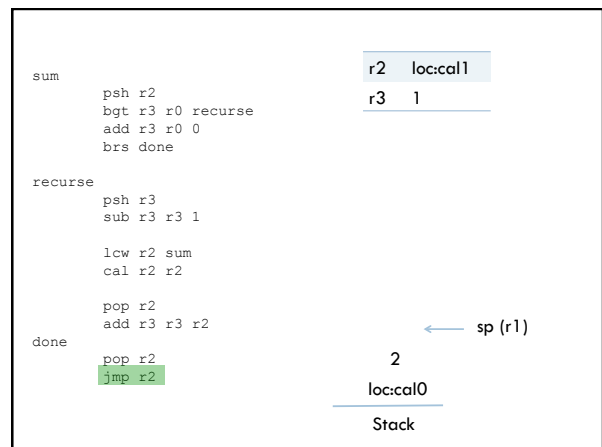
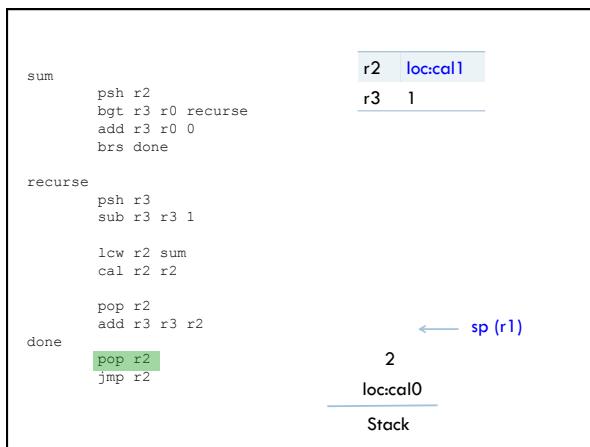
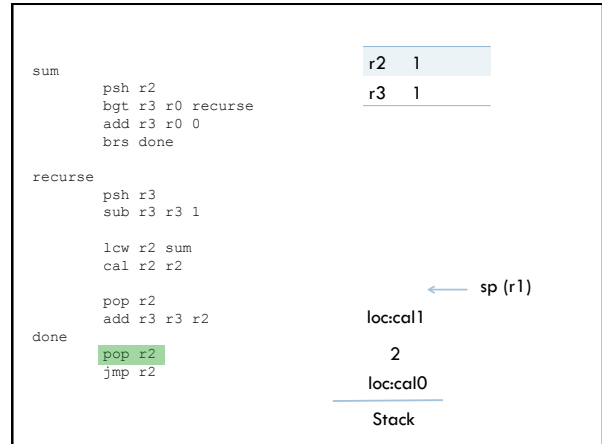
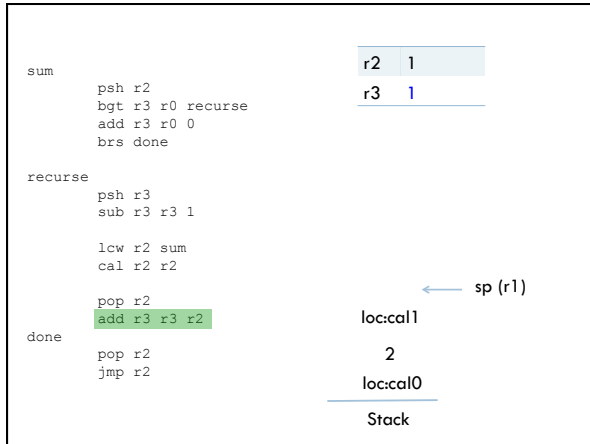


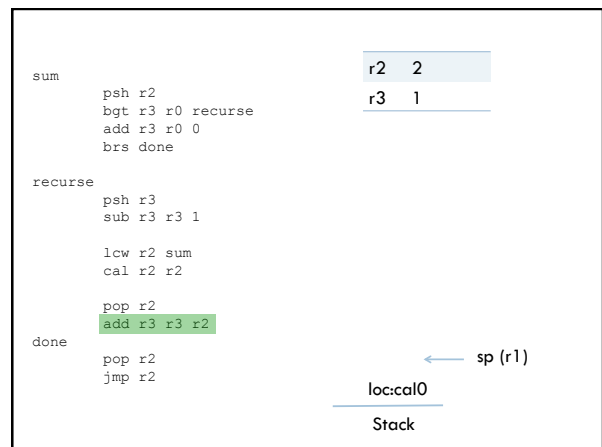
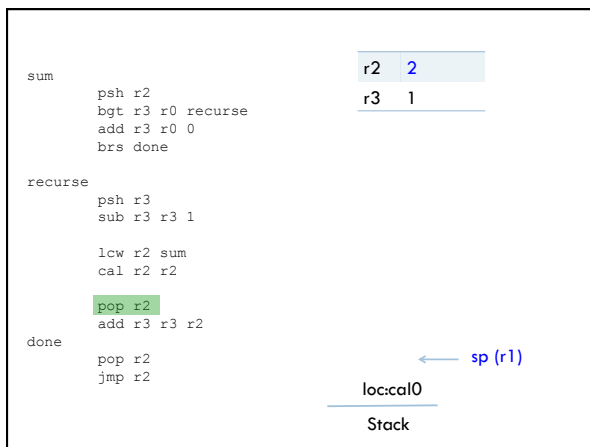
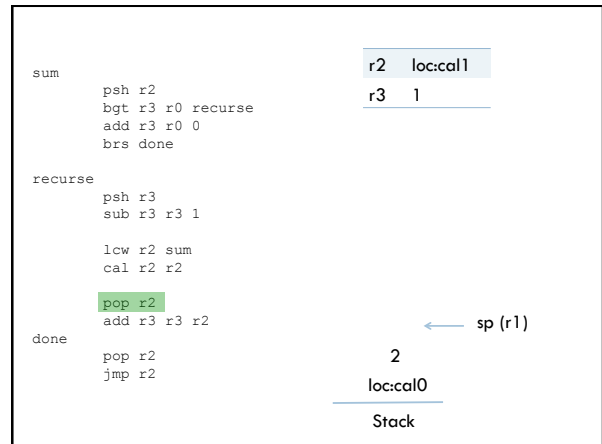
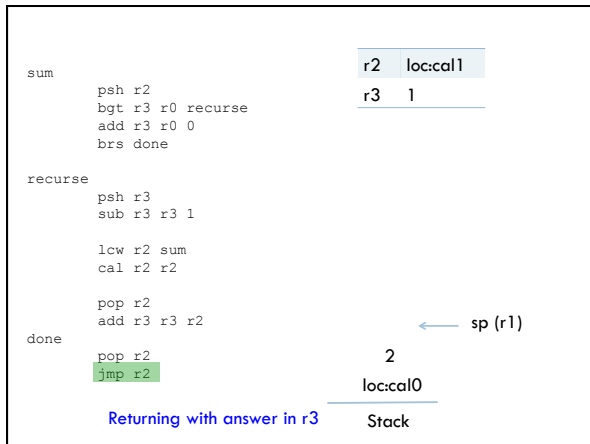




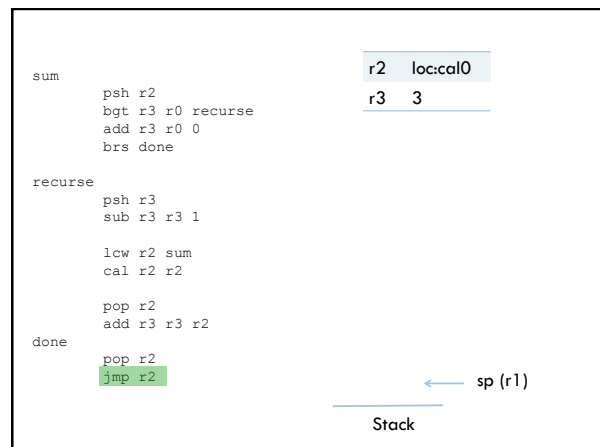
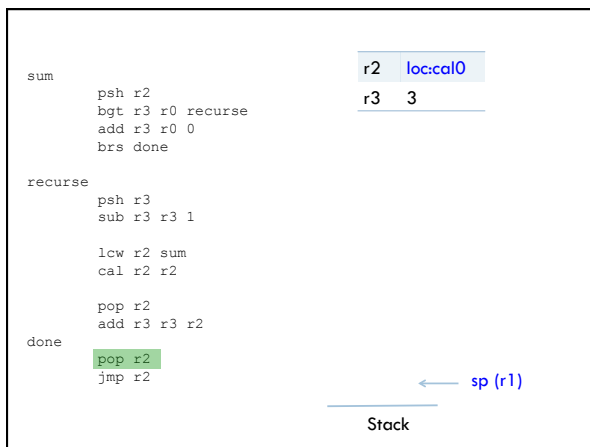
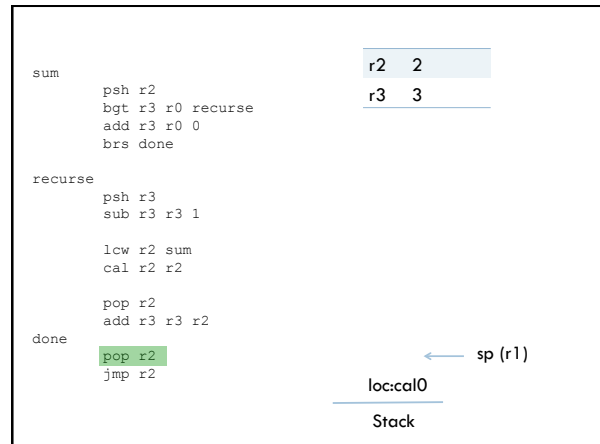
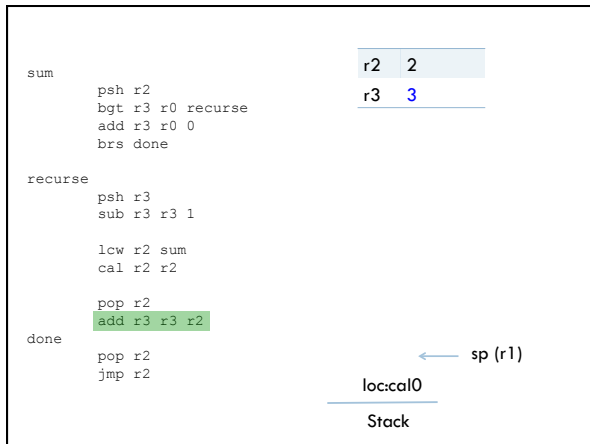












```

sum
  psh r2
  bgt r3 r0 recurse
  add r3 r0 0
  brs done

recurse
  psh r3
  sub r3 r3 1

  lcw r2 sum
  cal r2 r2

  pop r2
  add r3 r3 r2

done
  pop r2
  jmp r2
    
```

r2	loc:cal0
r3	3

← sp (r1)

Stack

Returning with answer in r3

Calling sum

```

  loa r3 r0

  lcw r2 sum
  cal r2 r2

  sto r3 r0
  hlt
    
```

r2	loc:cal0
r3	3

← sp (r1)

Stack

Calling sum

```

  loa r3 r0

  lcw r2 sum
  cal r2 r2

  sto r3 r0
  hlt
    
```

Print the answer: 3!

r2	loc:cal0
r3	3

← sp (r1)

Stack

Calling sum

```

  loa r3 r0

  lcw r2 sum
  cal r2 r2

  sto r3 r0
  hlt
    
```

r2	loc:cal0
r3	3

← sp (r1)

Stack

Calling sum

```

loa r3 r0
lcw r2 sum
cal r2 r2
sto r3 r0
hlt

```

r2	local0
r3	3

Notice that when we're all done, the stack is empty

### Real structure of CS52 program

```

; great comments at the top!
;
;   lcw r1 stack           Save address of highest end
;                           (highest address) of the stack in r1
;
;   instruction1          ; comment
;   instruction2          ; comment
;   ...
;
; ; stack area: 50 words
;
;   dat 100
;
stack

```

Reserve 50 words for the stack

### Structure of a single parameter function

```

fname
    psh r2          ; save return address on stack
    ...            ; do work using r3 as argument
                  ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2         ; return to caller

```

conventions:

- argument is in r3
- r1 is off-limits since it's used for the stack pointer
- return value goes in r3

### Functions with multiple arguments

```

fname
    psh r2          ; save return address on stack
    loa r2 r1 4     ; load the second parameter into r2
    ...            ; do work using r3 and r2 as arguments
                  ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2         ; return to caller

```

conventions:

- first argument is in r3
- r1 is off-limits since it's used for the stack pointer
- return value goes in r3

## Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4    ; load the second parameter into r2
    ...           ; do work using r3 and r2 as arguments
                  ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2         ; return to caller
```

loa  $R_a R_b$ :  $R_a = \text{mem}[R_b]$

loa  $R_a R_b S$ :  $R_a = \text{mem}[R_b + S]$

What does this operation do? What is the 4?

## Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4    ; load the second parameter into r2
    ...           ; do work using r3 and r2 as arguments
                  ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2         ; return to caller
```

loa  $R_a R_b$ :  $R_a = \text{mem}[R_b]$

loa  $R_a R_b S$ :  $R_a = \text{mem}[R_b + S]$

- r1 is the stack pointer and points at the top (next) slot
- stacks grow towards smaller memory values

## Functions with multiple arguments

```
fname
    psh r2          ; save return address on stack
    loa r2 r1 4    ; load the second parameter into r2
    ...           ; do work using r3 and r2 as arguments
                  ; put result in r3
    pop r2         ; restore return address from stack
    jmp r2         ; return to caller
```

- r1 is the stack pointer and points at the top (next) slot
- stacks grow towards smaller memory values
- r1+2 is then the top value of the stack
- r1+4 is the 2<sup>nd</sup> value of the stack

## Another recursive example

```
int mystery(int a, int b){
    if( b <= 0 ){
        return 0
    }
    else
        return a + mystery(a, b-1)
    }
}
```

What does this function do?

## Recursion

```

int mystery(int a, int b){
  if( b <= 0 ){
    return 0
  }
  else
    return a + mystery(a, b-1)
}
    
```

Multiplication...  $a*b$  (assuming  $b$  is positive)

Note to future Dave from past Dave: write the function up on the board ☺

```

mult
  psh r2      ; save the return address
  loa r2 r1 4 ; get at the 2nd argument, b
              ; a = r3, b = r2

  bgt r2 r0 else ; r2 > 0, i.e. recursive case
  add r3 r0 0   ; return 0
  brs endif

else
  sub r2 r2 1   ; r2 = b-1

  psh r3      ; save first argument, a, on stack
              ; (it's going to get overwritten by the return!)
  psh r2      ; add r2 as 2nd argument, r3 shouldn't have changed
  lw  r2 mult  ; call mult recursively
  cal r2 r2
  pop r0      ; pop 2nd argument off stack

  pop r2      ; pop 'a' into r2 off of the stack
  add r3 r3 r2 ; r3 = a + mult(a, b-1)

endif
  pop r2      ; get the return address
  jmp r2      ; return
    
```

Function startup

Base case

Recursive case

Recursive call

answer calculation

Function cleanup and return

```

mult
  psh r2      ; save the return address
  loa r2 r1 4 ; get at the 2nd argument, b
              ; a = r3, b = r2

  bgt r2 r0 else ; r2 > 0, i.e. recursive case
  add r3 r0 0   ; return 0
  brs endif

else
  sub r2 r2 1   ; r2 = b-1

  psh r3      ; save first argument, a, on stack
              ; (it's going to get overwritten by the return!)
  psh r2      ; add r2 as 2nd argument, r3 shouldn't have changed
  lw  r2 mult  ; call mult recursively
  cal r2 r2
  pop r0      ; pop 2nd argument off stack

  pop r2      ; pop 'a' into r2 off of the stack
  add r3 r3 r2 ; r3 = a + mult(a, b-1)

endif
  pop r2      ; get the return address
  jmp r2      ; return
    
```

Function startup

if( b <= 0 )  
return 0

mystery(a, b-1)

a + mystery(a, b-1)

Function cleanup and return

```

mult
  psh r2      ; save the return address
  loa r2 r1 4 ; get at the 2nd argument, b
              ; a = r3, b = r2

  bgt r2 r0 else ; r2 > 0, i.e. recursive case
  add r3 r0 0   ; return 0
  brs endif

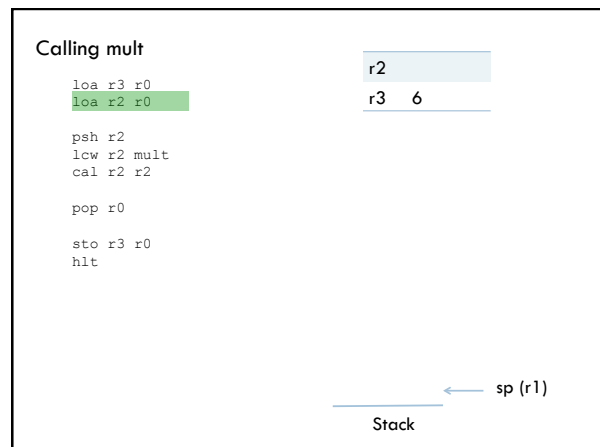
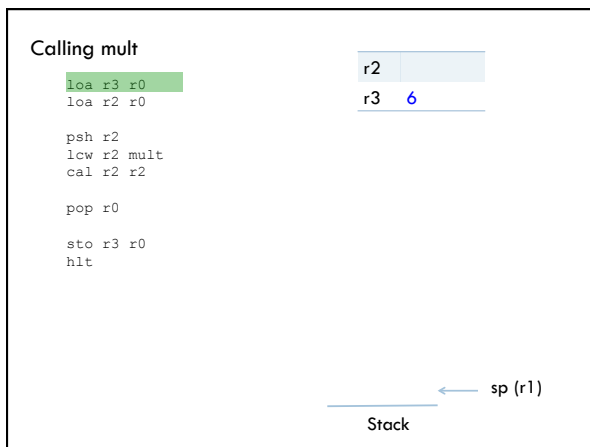
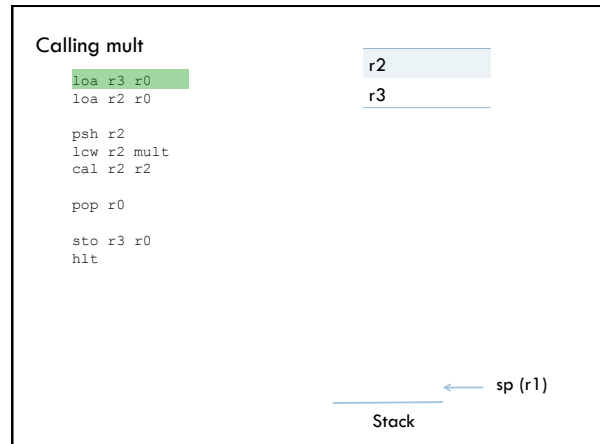
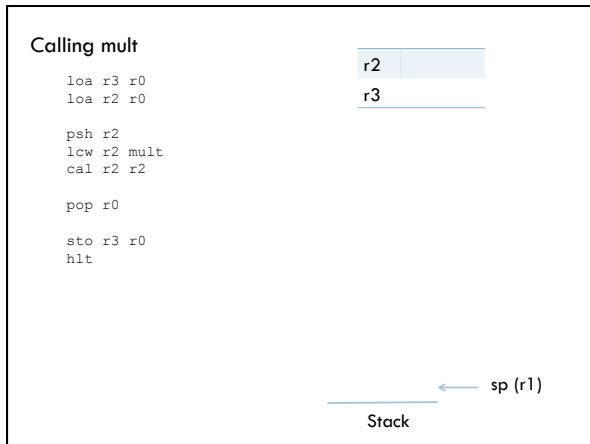
else
  sub r2 r2 1   ; r2 = a-1

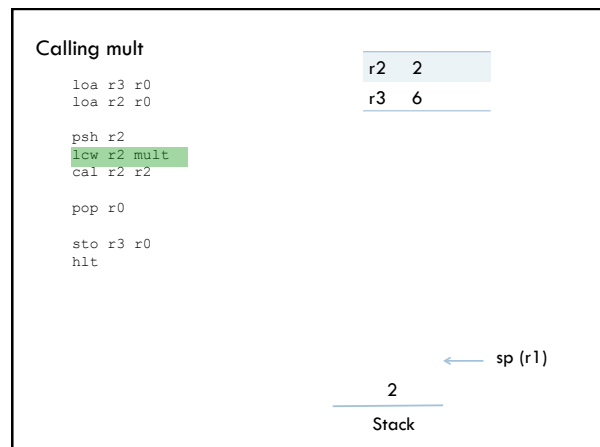
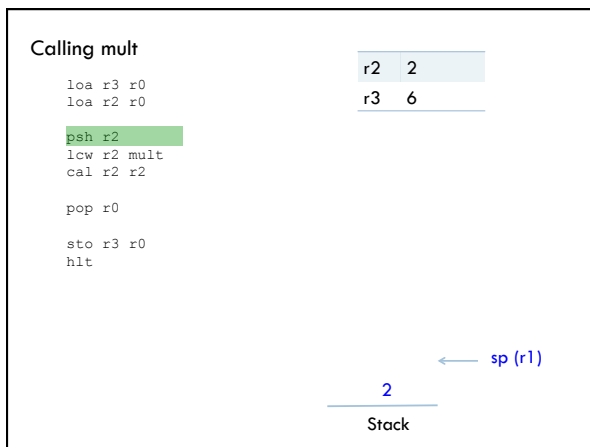
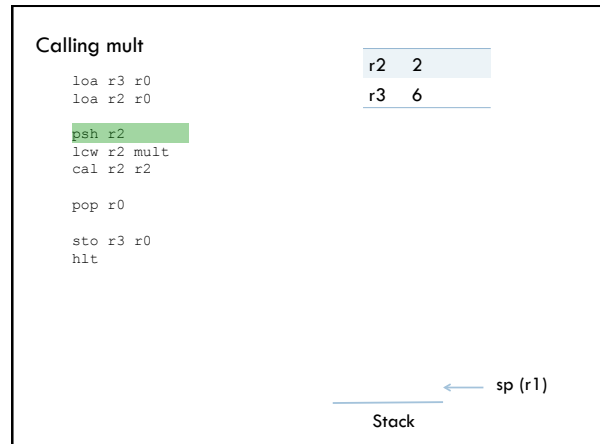
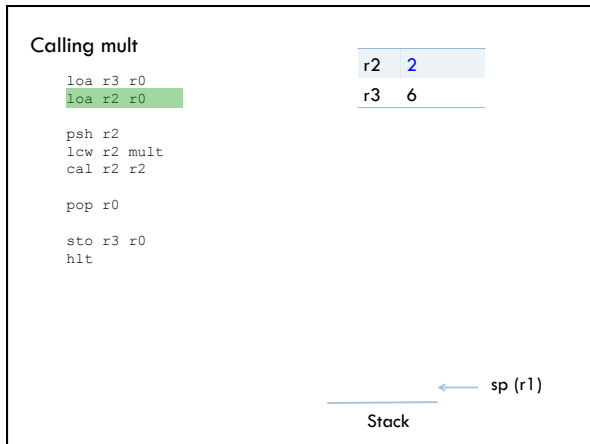
  psh r3      ; save first argument, a, on stack
              ; (it's going to get overwritten by the return!)
  psh r2      ; add r2 as 2nd argument, r3 shouldn't have changed
  lw  r2 mult  ; call mult recursively
  cal r2 r2
  pop r0      ; pop 2nd argument off stack

  pop r2      ; load a into r2 off of the stack
  add r3 r3 r2 ; r3 = a + mult(a, b-1)

endif
  pop r2      ; get the return address
  jmp r2      ; return
    
```

Notice symmetry of psh and pop





**Calling mult**

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hit
    
```

r2	loc: mult
r3	6

← sp (r1)

---

2

Stack

**Calling mult**

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hit
    
```

r2	loc: mult
r3	6

← sp (r1)

---

2

Stack

**mult**

```

psh r2
loa r2 r1 4

bgt r2 r0 else
add r3 r0 0
brs endif
else
sub r2 r2 1

psh r3
psh r2
lcw r2 mult
cal r2 r2
pop r0

pop r2
add r3 r3 r2
endif
pop r2
jmp r2
    
```

r2	loc: cal0
r3	6

← sp (r1)

---

2

Stack

**mult**

```

psh r2
loa r2 r1 4

bgt r2 r0 else
add r3 r0 0
brs endif
else
sub r2 r2 1

psh r3
psh r2
lcw r2 mult
cal r2 r2
pop r0

pop r2
add r3 r3 r2
endif
pop r2
jmp r2
    
```

r2	loc: cal0
r3	6

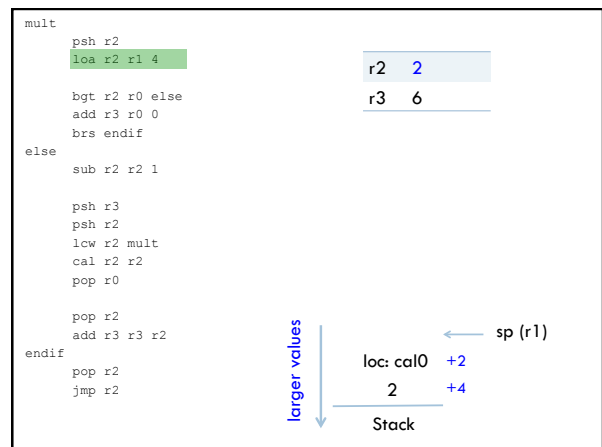
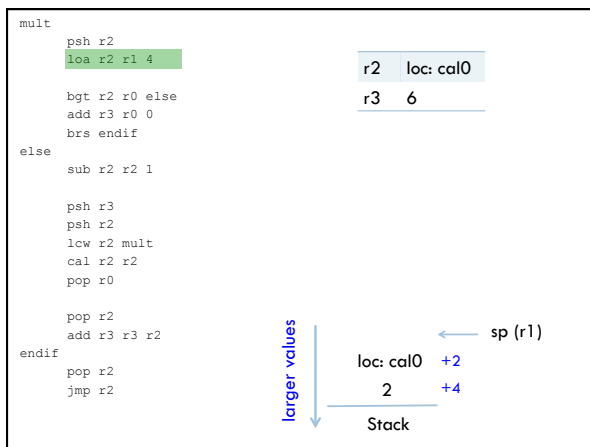
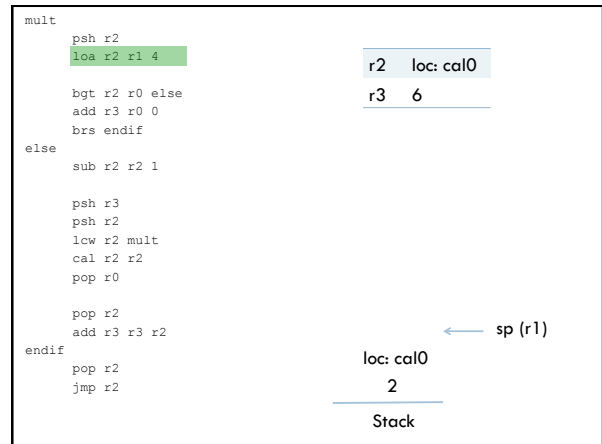
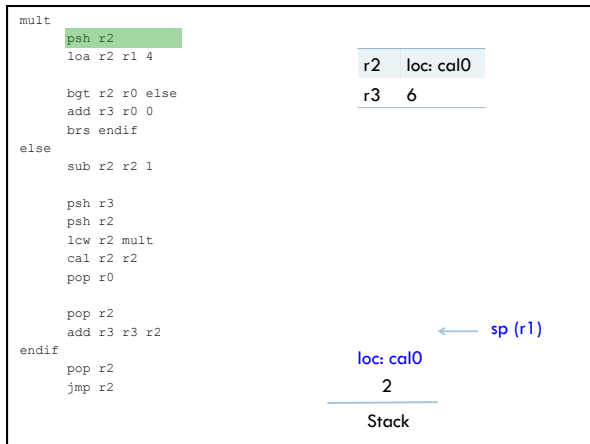
← sp (r1)

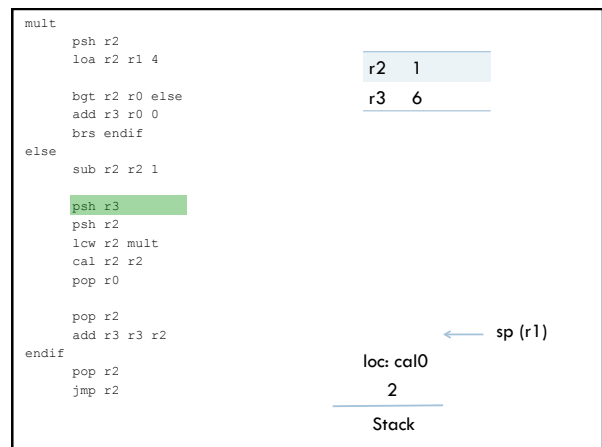
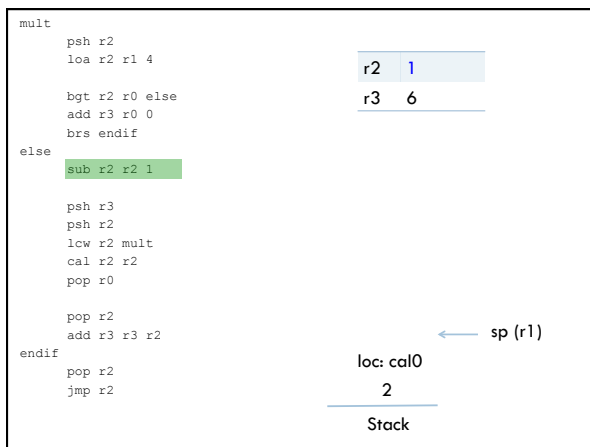
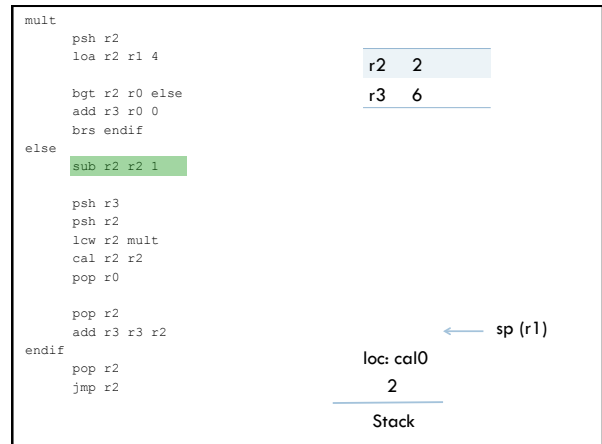
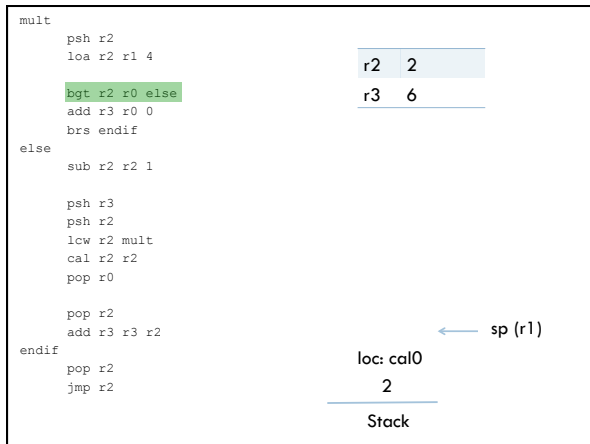
---

2

Stack







```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
  
```

r2	1
r3	6

Why psh r3?

← sp (r1)

6

loc: cal0

2

---

Stack

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
  
```

r2	1
r3	6

- We're about to make a function call
- The result of that call will go into r3 so we'll lose what's in there if we don't save it!

← sp (r1)

6

loc: cal0

2

---

Stack

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
  
```

r2	1
r3	6

← sp (r1)

6

loc: cal0

2

---

Stack

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
  
```

r2	1
r3	6

← sp (r1)

1

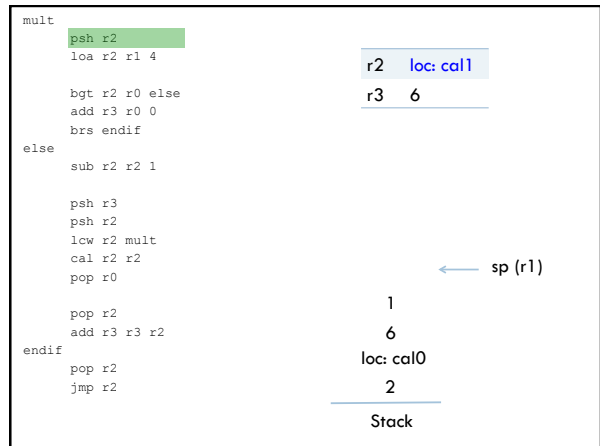
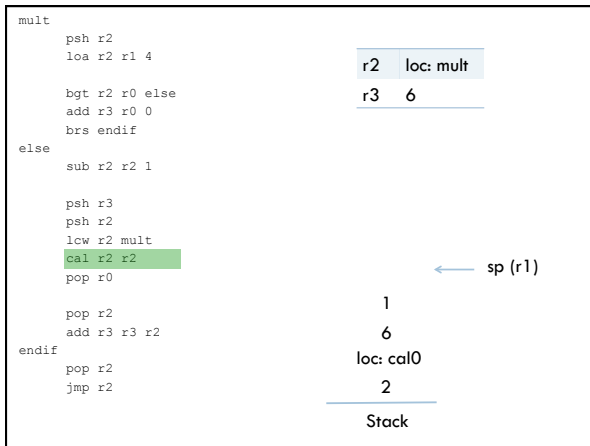
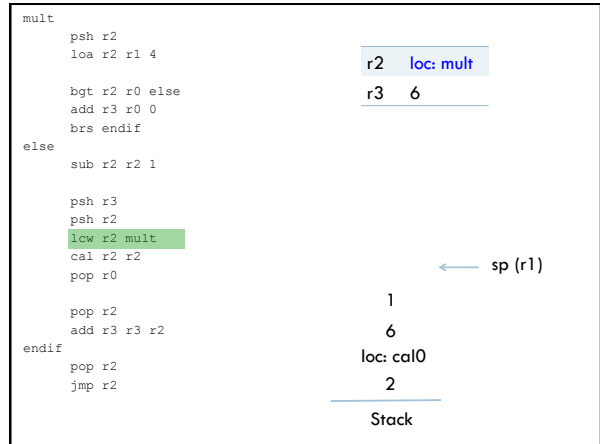
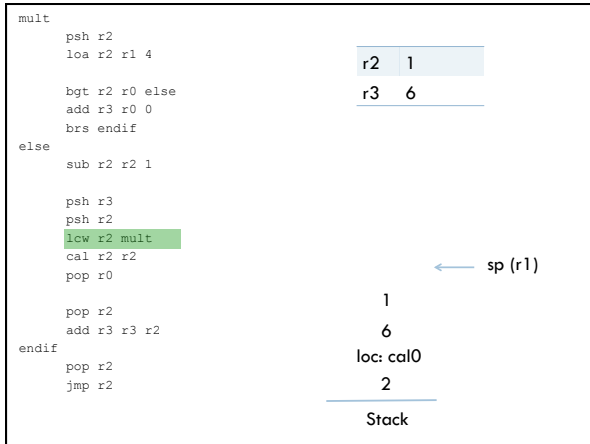
6

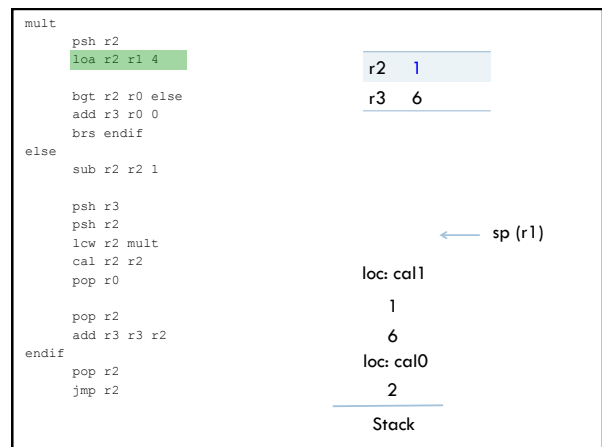
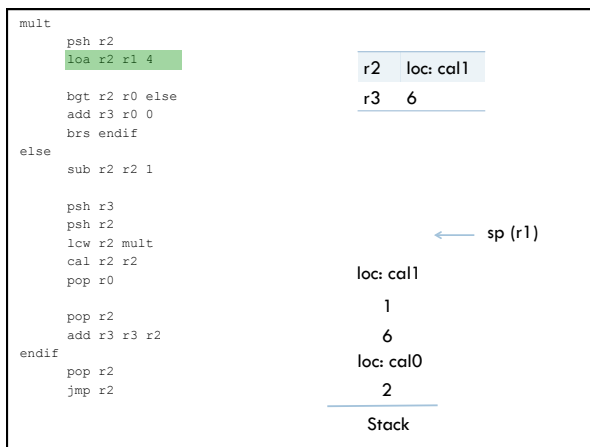
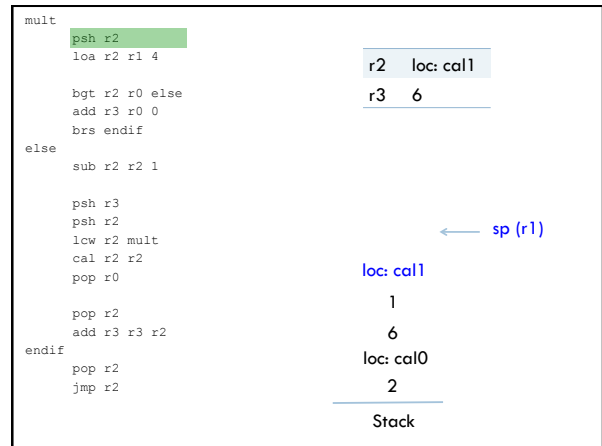
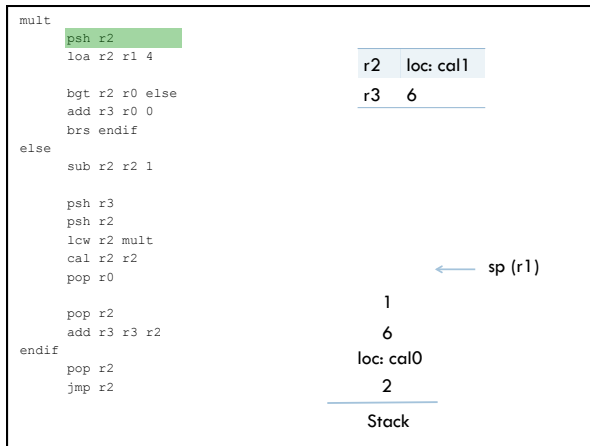
loc: cal0

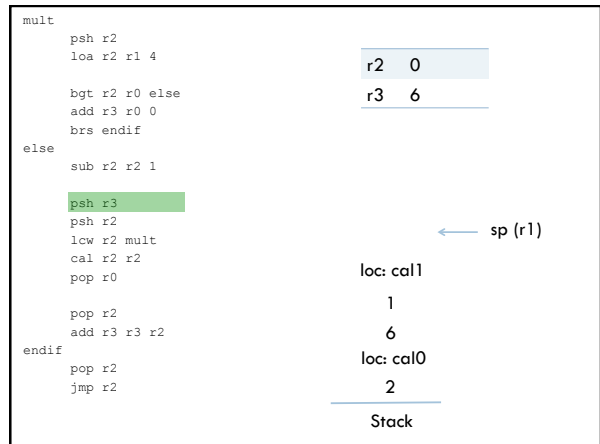
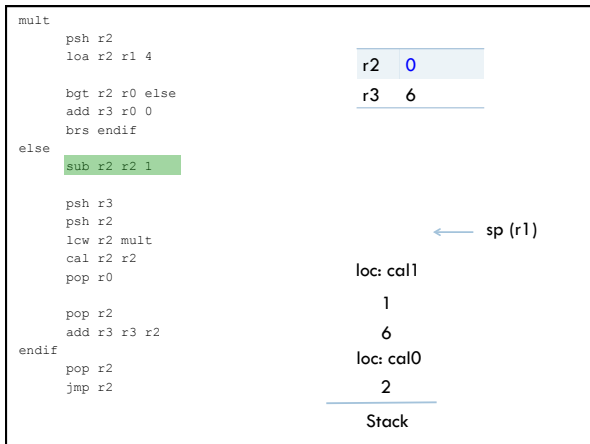
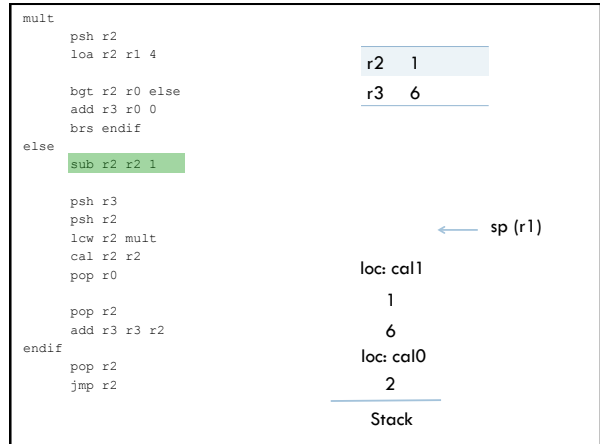
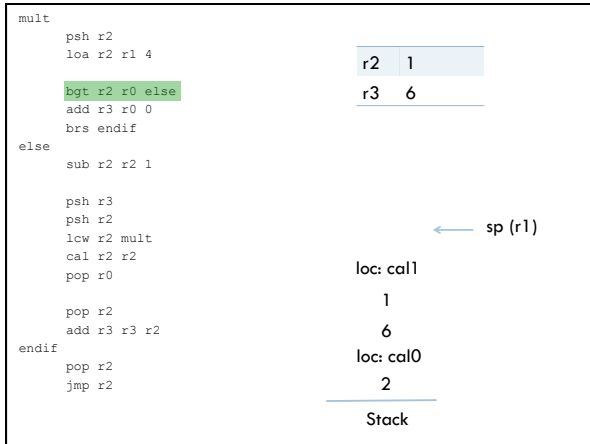
2

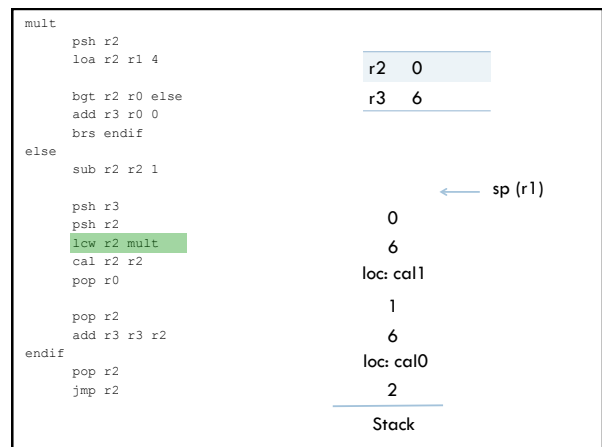
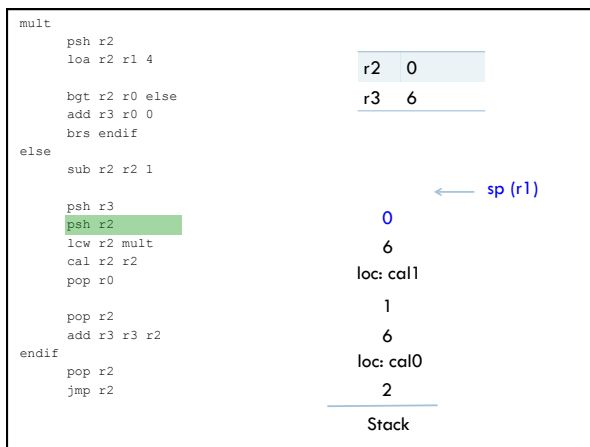
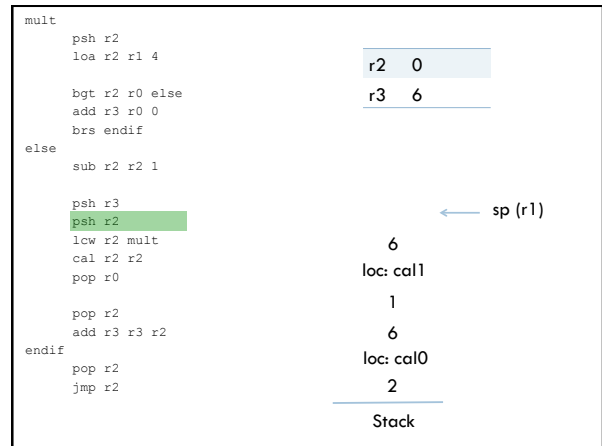
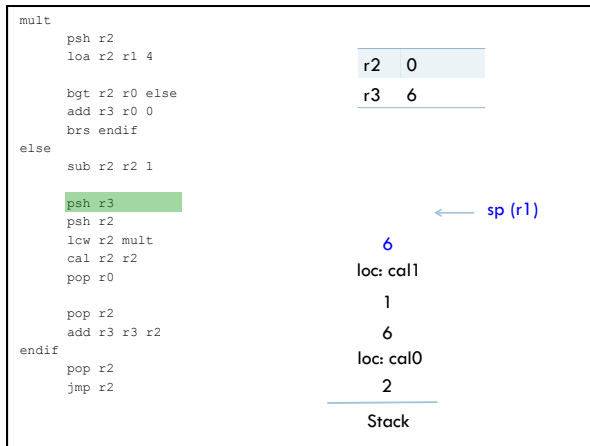
---

Stack









```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
  pop r2
  jmp r2

```

r2	loc: mult
r3	6
← sp (r1)	
	0
	6
	loc: cal1
	1
	6
	loc: cal0
	2
Stack	

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
  pop r2
  jmp r2

```

r2	loc: mult
r3	6
← sp (r1)	
	0
	6
	loc: cal1
	1
	6
	loc: cal0
	2
Stack	

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1

  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
  pop r2
  jmp r2

```

r2	loc: cal1
r3	6
← sp (r1)	
	0
	6
	loc: cal1
	1
	6
	loc: cal0
	2
Stack	

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1

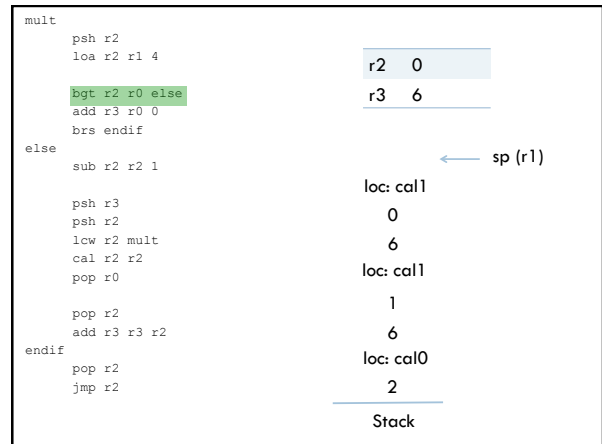
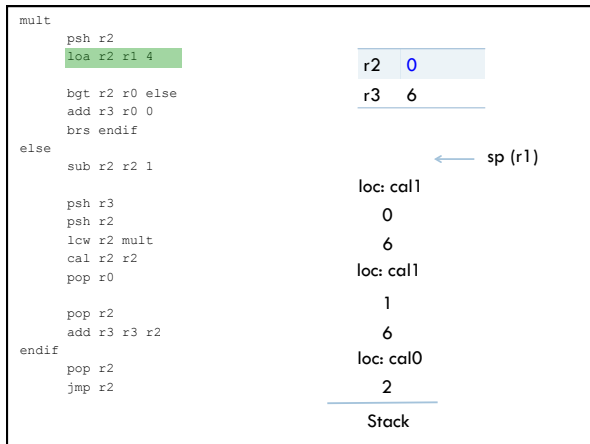
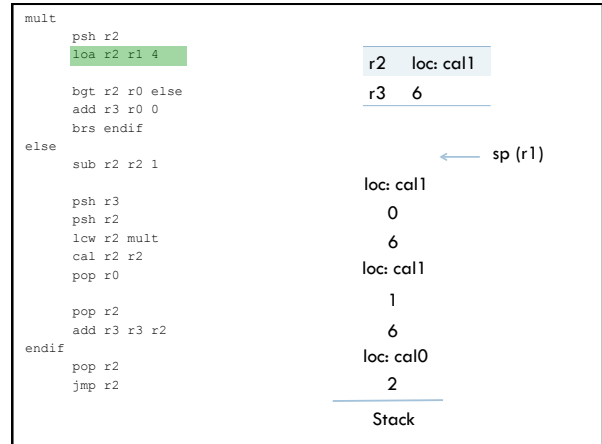
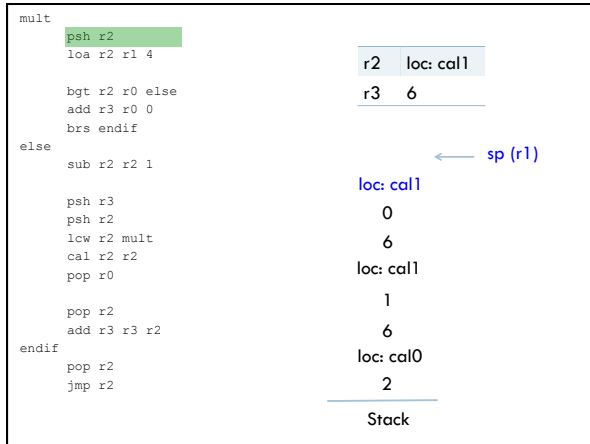
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0

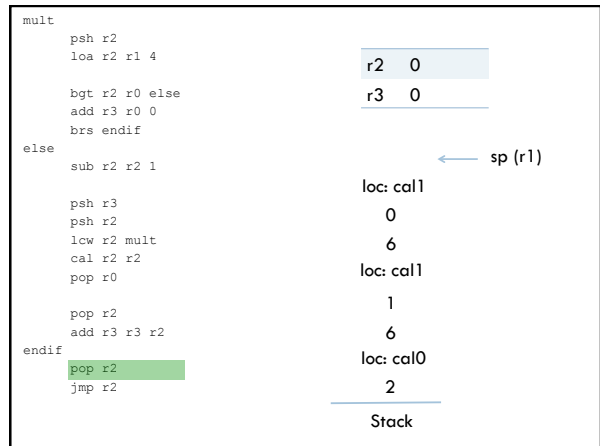
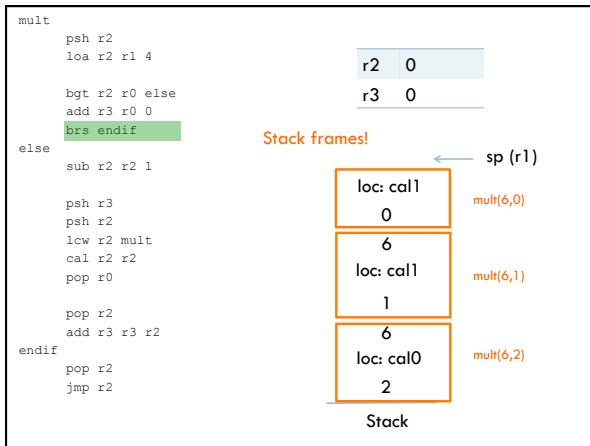
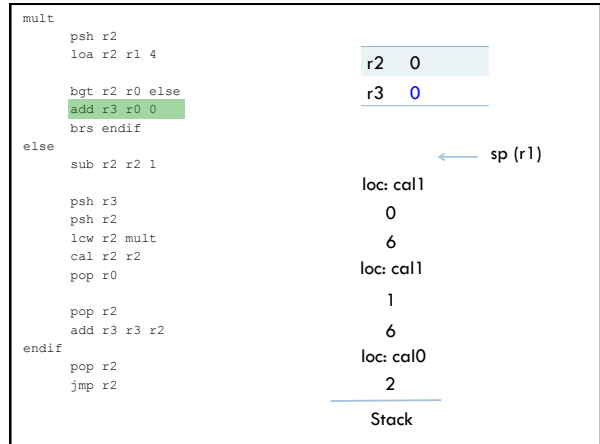
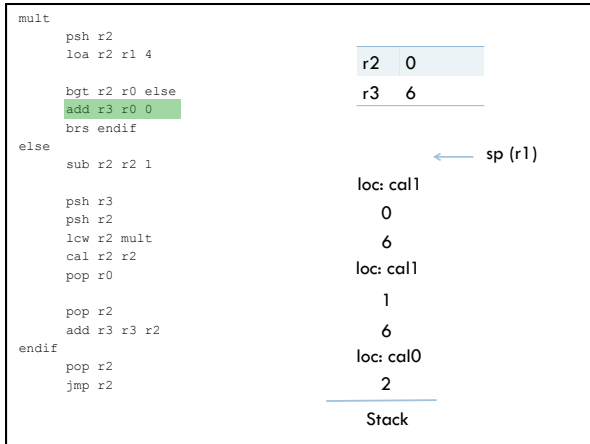
  pop r2
  add r3 r3 r2
endif
  pop r2
  jmp r2

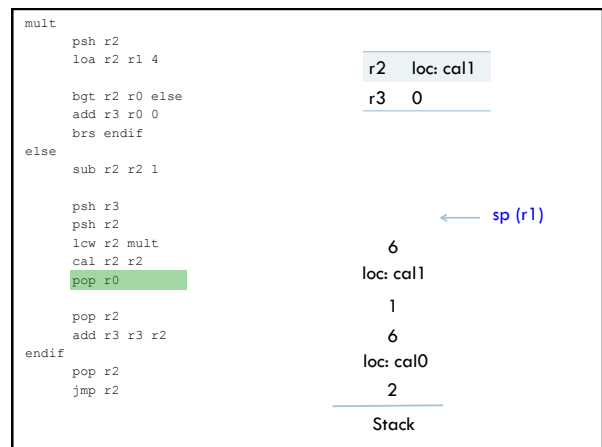
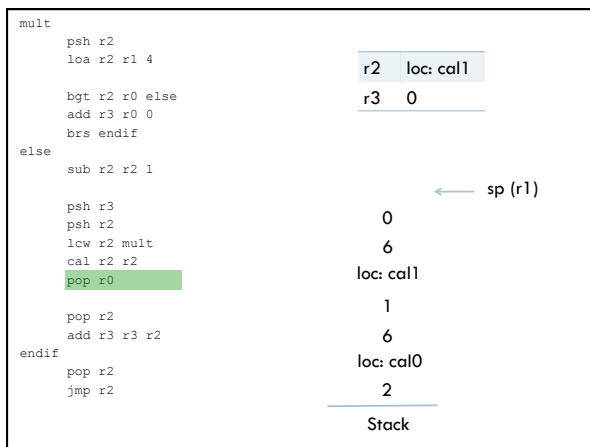
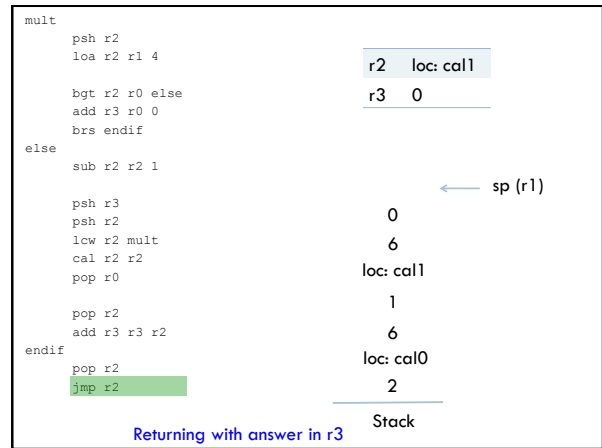
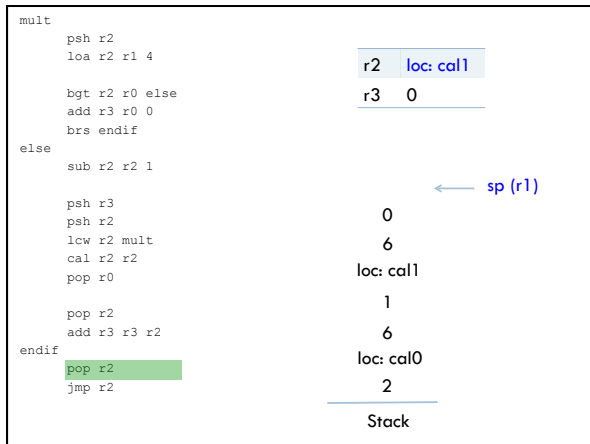
```

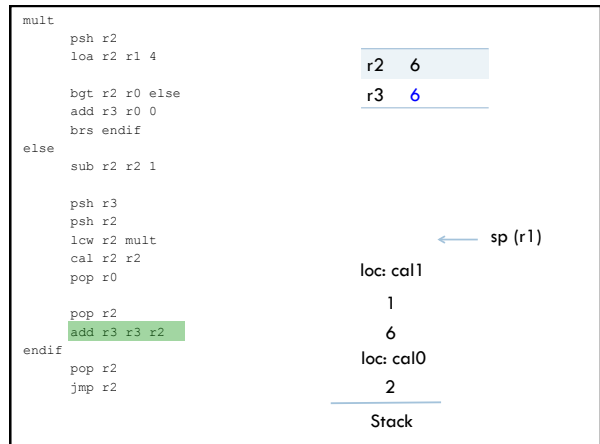
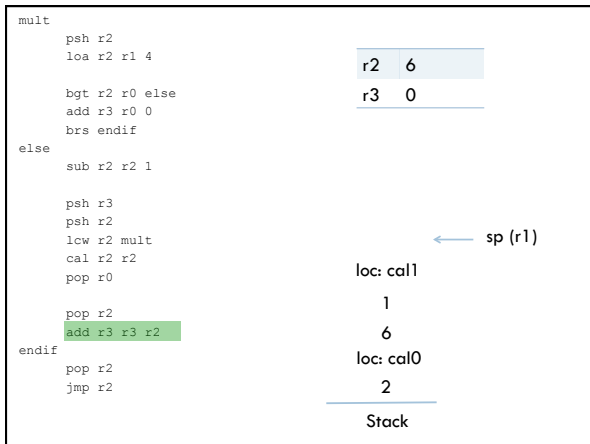
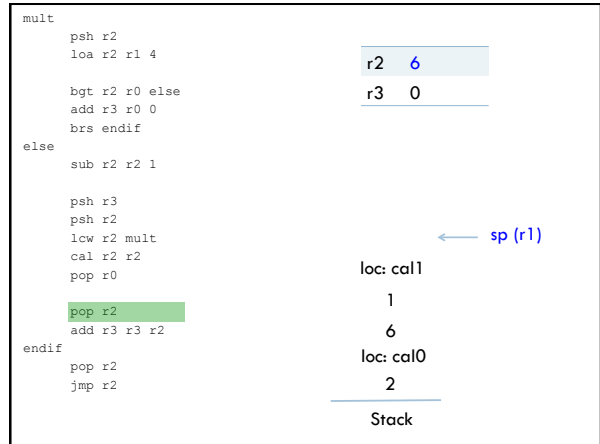
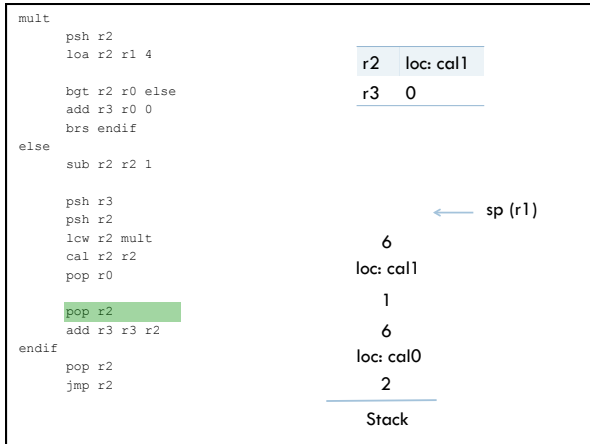
r2	loc: cal1
r3	6
← sp (r1)	
	0
	6
	loc: cal1
	1
	6
	loc: cal0
	2
Stack	

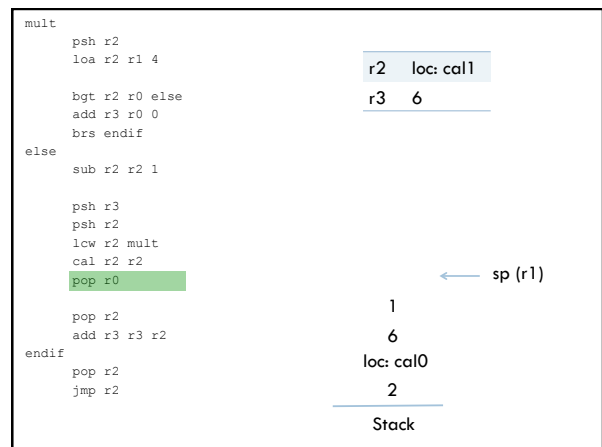
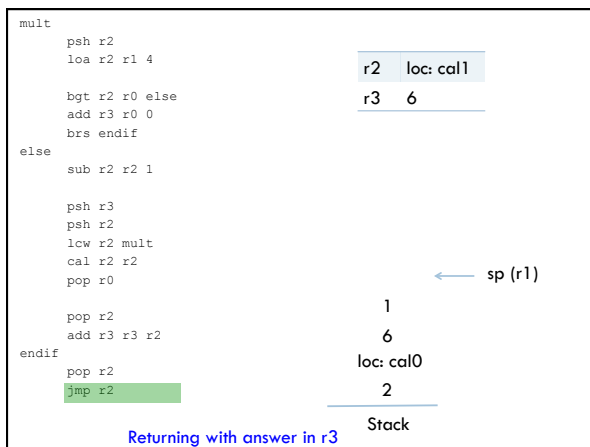
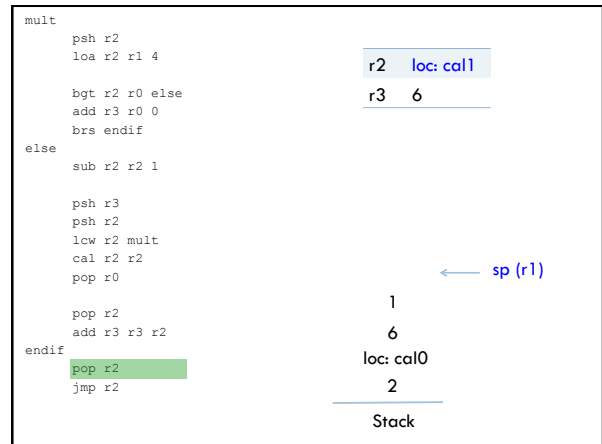
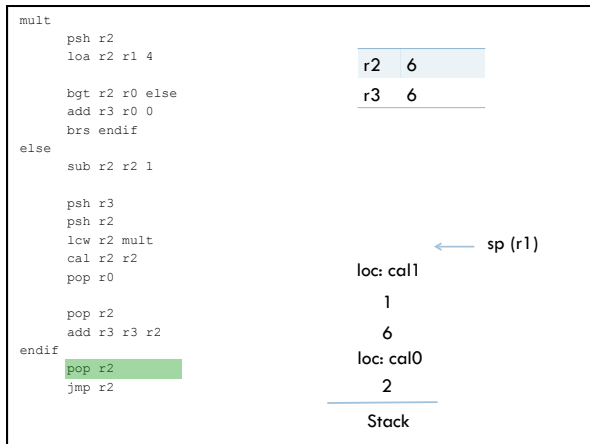












```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1

  psh r3
  psh r2
  lclw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
  
```

Stack: 6, loc: cal0 (2)

Registers: r2 (loc: cal1), r3 (6)

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1

  psh r3
  psh r2
  lclw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
  
```

Stack: 6, loc: cal0 (2)

Registers: r2 (loc: cal1), r3 (6)

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1

  psh r3
  psh r2
  lclw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
  
```

Stack: 6, loc: cal0 (2)

Registers: r2 (6), r3 (6)

```

mult
  psh r2
  loa r2 r1 4

  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1

  psh r3
  psh r2
  lclw r2 mult
  cal r2 r2
  pop r0

  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2
  
```

Stack: 6, loc: cal0 (2)

Registers: r2 (6), r3 (6)

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0
  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2

```

r2	6
r3	12

← sp (r1)

loc: cal0

---

2

Stack

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0
  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2

```

r2	6
r3	12

← sp (r1)

loc: cal0

---

2

Stack

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0
  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2

```

r2	loc: cal0
r3	12

← sp (r1)

2

Stack

```

mult
  psh r2
  loa r2 r1 4
  bgt r2 r0 else
  add r3 r0 0
  brs endif
else
  sub r2 r2 1
  psh r3
  psh r2
  lcw r2 mult
  cal r2 r2
  pop r0
  pop r2
  add r3 r3 r2
endif
pop r2
jmp r2

```

r2	loc: cal0
r3	12

← sp (r1)

2

Stack

Returning with answer in r3

Calling mult

r2	loc: cal0
r3	12

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hit
    
```

← sp (r1)

2

---

Stack

Calling mult

r2	loc: cal0
r3	12

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0
hit
    
```

← sp (r1)

---

Stack

Calling mult

r2	loc: cal0
r3	12

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0

hit
    
```

← sp (r1)

---

Stack

Calling mult

r2	loc: cal0
r3	12

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0

sto r3 r0

hit
    
```

← sp (r1)

Print the answer: 12!

---

Stack



**Calling mult**

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 mult
cal r2 r2

pop r0
sto r3 r0
hit

```

Print the answer: 12!

r2	loc: cal0
r3	12

← sp (r1)

Stack

## multiply the easy way

look at mult\_easy.a52 code

we can import libraries (really just functions in other files) using the "inc" command

the included files must be in the same directory as the .a52 file running

## A final aside

$10101_2$   
 $10_2$   
 $11_2$   
 $1_2$   
 $1010_2$   
 $101010010010010010000111111101001_2$

Which of these binary numbers is even?

## A quick aside

$10101_2 = 21$   
 $10_2 = 2$   
 $11_2 = 3$   
 $1_2 = 1$   
 $1010_2 = 10$   
 $101010010010010010000111111101001_2 = 5,675,487,209$

Is there an easier way to tell than just calculating the value?

## A quick aside

The last digit represents the  $2^0 = 1$ s digit

*All other digits represent even values since they are powers of 2*

Therefore:

- If the rightmost digit is 1 = odd number
- If the rightmost digit is 0 = even number

## Bitwise logical operators

instruction name    arguments

add	}	RRR or RRS
sub		
and		
orr		
xor		

Perform **and**, **or** and **xor** per bit of the number

## Bitwise logical operators

```
add r2 r0 3
add r3 r0 6
and r3 r3 r2
```

## Bitwise logical operators

```
add r2 r0 3
add r3 r0 6
and r3 r3 r2
```

What are r2 and r3 in binary after these instructions?

## Bitwise logical operators

```
add r2 r0 3
add r3 r0 6
and r3 r3 r2
```

```
r2:  11
r3: 110
```

## Bitwise logical operators

```
add r2 r0 3
add r3 r0 6
and r3 r3 r2 Perform and, or and xor per bit of the number
```

```
r2:  11
r3: 110
```

## Bitwise logical operators

```
add r2 r0 3
add r3 r0 6
and r3 r3 r2 Perform and, or and xor per bit of the number
```

```
r2:  11
r3: 110
```

## Bitwise logical operators

```
add r2 r0 3
add r3 r0 6
and r3 r3 r2 Perform and, or and xor per bit of the number
```

```
r2:  11
r3: 110
    0
```

## Bitwise logical operators

add r2 r0 3

add r3 r0 6

**and r3 r3 r2** Perform **and**, **or** and **xor** per bit of the number

r2:	11
r3:	110
0	

## Bitwise logical operators

add r2 r0 3

add r3 r0 6

**and r3 r3 r2** Perform **and**, **or** and **xor** per bit of the number

r2:	11
r3:	110
10	

## Bitwise logical operators

add r2 r0 3

add r3 r0 6

**and r3 r3 r2** Perform **and**, **or** and **xor** per bit of the number

r2:	11
r3:	110
10	

## Bitwise logical operators

add r2 r0 3

add r3 r0 6

**and r3 r3 r2** Perform **and**, **or** and **xor** per bit of the number

r2:	011
r3:	110
10	

## Bitwise logical operators

```
add r2 r0 3
```

```
add r3 r0 6
```

```
and r3 r3 r2
```

Perform **and**, **or** and **xor** per bit of the number

r2:	0	1	1
r3:	1	1	0
010			

## CS52 programming advice

1. Match your psh and pops
2. Follow the register conventions
3. Develop code incrementally
4. Debugging: write out stack, registers, etc. on paper and compare against system execution

## Examples from this lecture

<http://www.cs.pomona.edu/~dkauchak/classes/cs52/examples/cs41b/>

## Another example

I didn't have time to cover the next example in class, but left it in the notes as another example of a function that takes two parameters

### Another example

```

max
  psh r2
  loa r2 r1 4

  bge r3 r2 endif
  add r3 r2 0
endif
  pop r2
  jmp r2

```

What does this code do?

### Another example

```

max
  psh r2
  loa r2 r1 4

  bge r3 r2 endif
  add r3 r2 0
endif
  pop r2
  jmp r2

```

max, as a function!

### Calling max

```

  loa r3 r0
  loa r2 r0

  psh r2
  lcw r2 max
  cal r2 r2
  pop r0

  sto r3 r0
  hlt

```

Anything different?

### Calling max

```

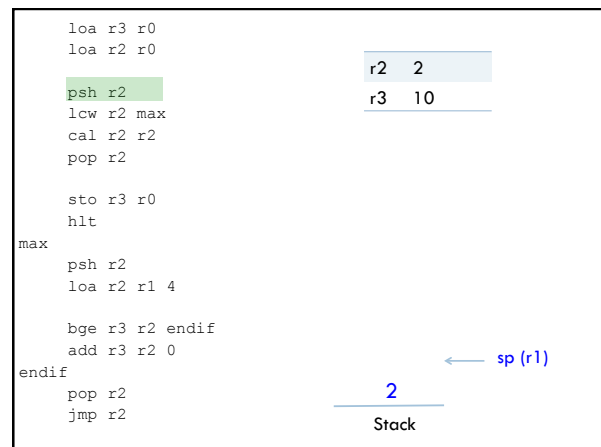
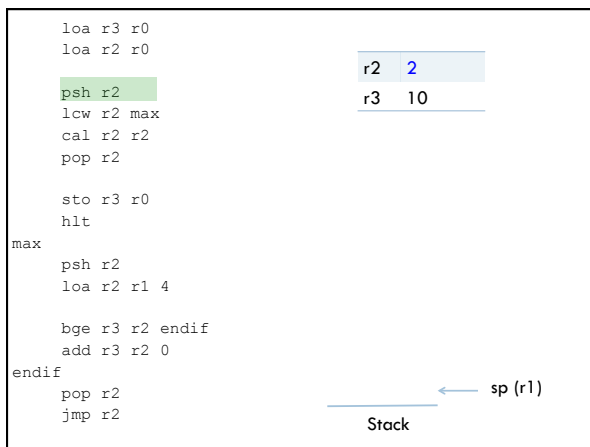
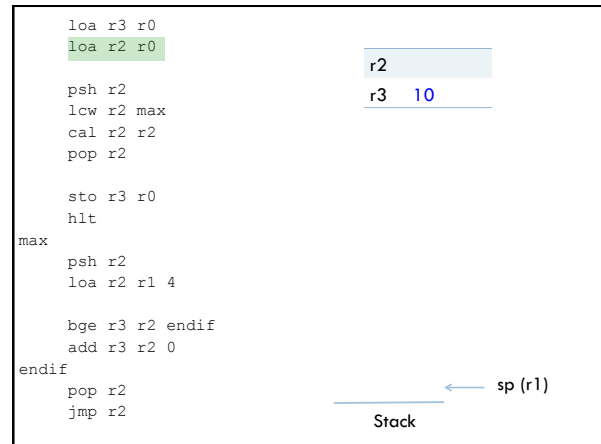
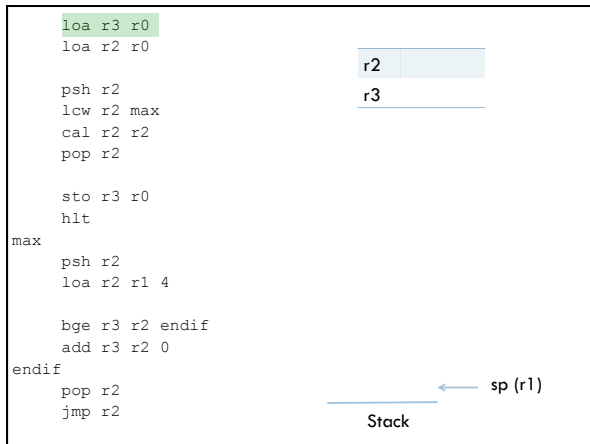
  loa r3 r0
  loa r2 r0

  psh r2
  lcw r2 max
  cal r2 r2
  pop r0

  sto r3 r0
  hlt

```

For the second argument, psh it on the stack



```

loa r3 r0
loa r2 r0

psh r2
lcw r2 max
cal r2 r2
pop r2

sto r3 r0
hlt
max
psh r2
loa r2 r1 4

bge r3 r2 endif
add r3 r2 0
endif
pop r2
jmp r2

```

r2	2
r3	10

← sp (r1)

2  
Stack

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 max
cal r2 r2
pop r2

sto r3 r0
hlt
max
psh r2
loa r2 r1 4

bge r3 r2 endif
add r3 r2 0
endif
pop r2
jmp r2

```

r2	max
r3	10

← sp (r1)

2  
Stack

Notice that we overwrite the value in r2  
If we hadn't saved it on the stack, it would have been lost

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 max
cal r2 r2
pop r2

sto r3 r0
hlt
max
psh r2
loa r2 r1 4

bge r3 r2 endif
add r3 r2 0
endif
pop r2
jmp r2

```

r2	max
r3	10

← sp (r1)

2  
Stack

```

loa r3 r0
loa r2 r0

psh r2
lcw r2 max
cal r2 r2
pop r2

sto r3 r0
hlt
max
psh r2
loa r2 r1 4

bge r3 r2 endif
add r3 r2 0
endif
pop r2
jmp r2

```

r2	loc: cal
r3	10

← sp (r1)

2  
Stack



