

CS52 MACHINE

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CS 52 – Spring 2017

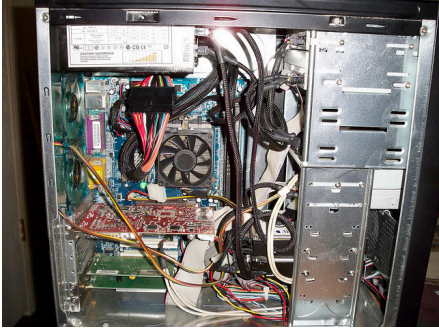
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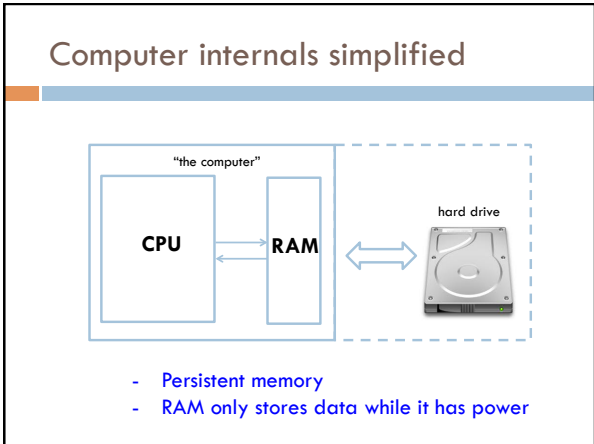
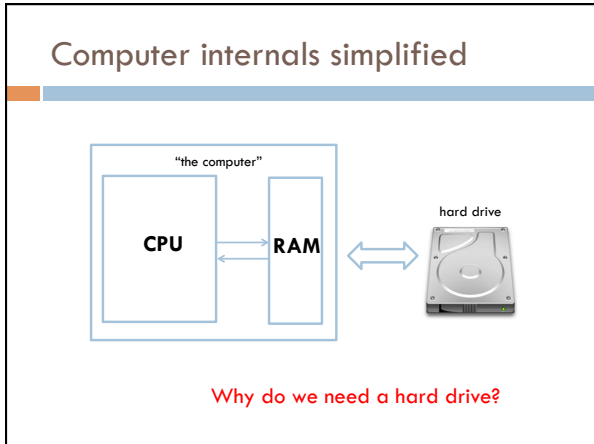
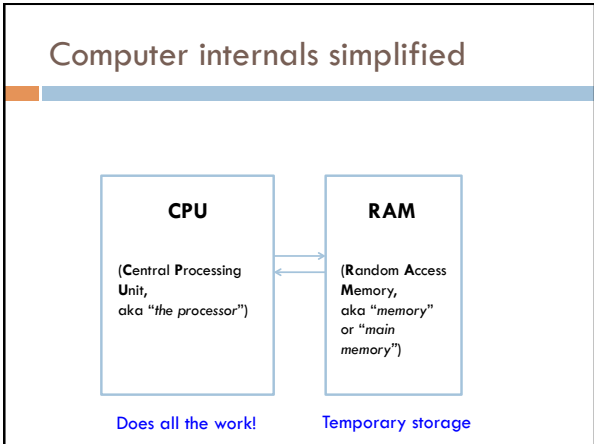
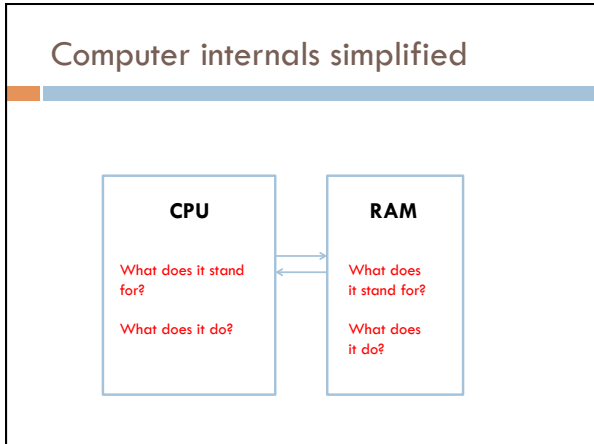
- Midterm 1
- Assignment 3
- Assignment 4

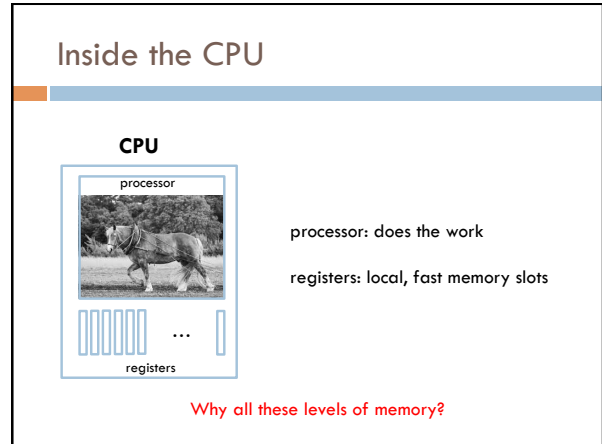
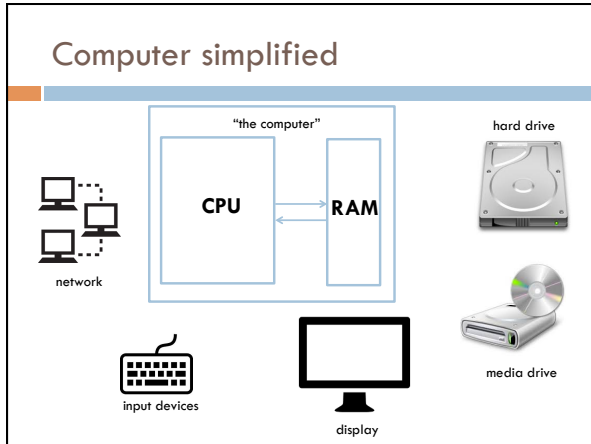
Examples from this lecture

<http://www.cs.pomona.edu/~dkauchak/classes/cs52/examples/cs52machine/>

Computer internals

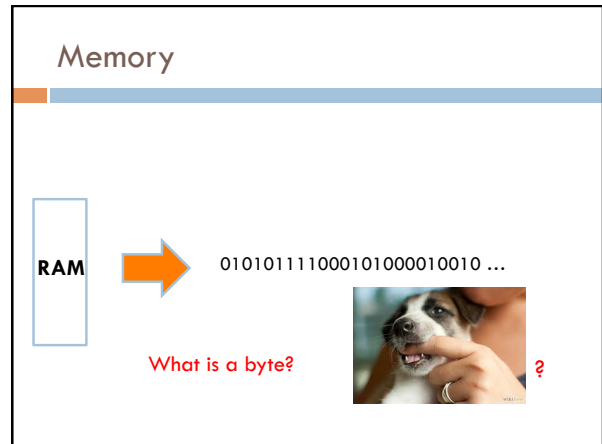






Memory speed

operation	access time	times slower than register access	for comparison ...
register	0.3 ns	1	1 s
RAM	120 ns	400	6 min
Hard disk	1ms	~million	1 month
google.com	0.4s	~billion	30 years



Memory

RAM → 01010111 10001010 00010010 ...

byte = 8 bits
byte is abbreviated as B

My laptop has 16GB (gigabytes) of memory. How many bits is that?

Memory sizes

	bits
byte	8
kilobyte (KB)	2^{10} bytes = ~8,000
megabyte (MB)	2^{20} = ~ 8 million
gigabyte (GB)	2^{30} = ~8 billion

My laptop has 16GB (gigabytes) of memory. How many bits is that?

Memory sizes

	bits
byte	8
kilobyte (KB)	2^{10} bytes = ~8,000
megabyte (MB)	2^{20} = ~ 8 million
gigabyte (GB)	2^{30} = ~8 billion

~128 billion bits!

Memory

RAM →

address	bits
0	01010111
1	10001010
2	00010010
3	01011010
...	...

Memory is byte addressable

Memory

address	
0	01010111
1	10001010
2	00010010
3	01011010
...	...

Memory is organized into "words", which is the most common functional unit

Memory

address	32-bit words
0	10101011 10001010 00010010 01011010
4	11001011 00001110 01010010 01010110
8	10111011 10010010 00000000 01110100
...	...

Most modern computers use 32-bit (4 byte) or 64-bit (8 byte) words

Memory in the CS52 Machine

address	16-bit words
0	10101011 10001010
2	00010010 01011010
4	11001011 00001110
...	...

We'll use 16-bit words for our model (the CS52 machine)

CS52 machine

CPU

- processor (includes a picture of a horse)
- registers (four vertical bars)
- ic: instruction counter (location in memory of the next instruction in memory)
- r0: holds the value 0 (read only)
- r1, r2, r3: general purpose read/write

In executing a program, the CS52 Machine follows a simple loop:

- The machine fetches the value at `mem[ic]` for use as an instruction.
- The machine increments the value in `ic` by 2.
- The machine decodes and carries out the instruction.

ic

r0

r1

r2

r3

instruction counter
(location in memory of the next instruction in memory)

holds the value 0 (read only)


— general purpose

— read/write

CS52 machine instructions

CPU

processor



registers

What types of operations might we want to do (think really basic)?

CS52 machine code

Four main types of instructions

1. math
2. branch/conditionals
3. memory
4. control the machine (e.g. stop it)

instruction name	arguments
add	} RRR or RRS
sub	
and	
orr	
xor	

instruction name arguments

```

add }
sub }
and } RRR or RRS
orr }
xor }
    
```

instruction/operation name
(always three characters)

instruction name arguments

```

add }
sub }
and } RRR or RRS
orr }
xor }
    
```

operation arguments
R = register (e.g. r0)
S = signed number (byte)

instruction name arguments

```

add }
sub }
and } RRR or RRS
orr }
xor }
    
```

1st R: register where the answer will go
2nd R: register of first operand
3rd S/R: register/value of second operand

add r1 r2 r3

What does this do?

1st R: register where the answer will go
2nd R: register of first operand
3rd S/R: register/value of second operand

add r1 r2 r3

$$r1 = r2 + r3$$

Add contents of registers r2 and r3 and store the result in r1

1st R: register where the answer will go
 2nd R: register of first operand
 3rd S/R: register/value of second operand

add r2 r1 10

What does this do?

1st R: register where the answer will go
 2nd R: register of first operand
 3rd S/R: register/value of second operand

add r2 r1 10

$$r2 = r1 + 10$$

Add 10 to the contents of register r1 and store in r2

1st R: register where the answer will go
 2nd R: register of first operand
 3rd S/R: register/value of second operand

add r1 r0 8
 neg r2 r1
 sub r2 r1 r2

What number is in r2?

1st R: register where the answer will go
 2nd R: register of first operand
 3rd S/R: register/value of second operand


```

add r1 r0 8    r1 = 8
neg r2 r1      r2 = -8, r1 = 8
sub r2 r1 r2   r2 = 16

```

1st R: register where the answer will go
 2nd R: register of first operand
 3rd S/R: register/value of second operand

Accessing memory

```

sto }
loa } RRS

```

sto = save data in register TO memory
 loa = put data FROM memory into a register

sto r1 r2 ; store the contents of r1 to mem[r2]
 loa r1 r2 ; get data from mem[r2] and put into r1

Accessing memory

```

sto }
loa } RRS

```

sto = save data in register TO memory
 loa = put data FROM memory into a register

Special cases:

- saving TO (sto) address 0 prints
- reading from (loa) address 0 gets input from user

Basic structure of CS52 program

```

; great comments at the top!
;
instruction1      ; comment
instruction2      ; comment
...
hlt

```

 whitespace before operations/instructions

Running the CS52 machine

Look at subtract.a52

- load two numbers from the user
- subtract
- print the result

CS52 simulator

Different windows

- ▣ Memory (left)
- ▣ Instruction execution (right)
- ▣ Registers
- ▣ I/O and running program

```

brs B
beq
bne } RRB
blt
bge
bgt
ble

```

1st R: first register for comparison
 2nd R: second register in comparison
 3rd B: label

```
beq r3 r0 done
```

What does this do?

1st R: first register for comparison
 2nd R: second register in comparison
 3rd B: label

beq r3 r0 done

If $r3 = 0$, branch to the label "done"
if not (else) ic is incremented as normal to
the next instruction

1st R: first register for comparison
2nd R: second register in comparison
3rd B: label

ble r2 r3 done

What does this do?

1st R: first register for comparison
2nd R: second register in comparison
3rd B: label

ble r2 r3 done

If $r2 \leq r3$, branch to the label done

1st R: first register for comparison
2nd R: second register in comparison
3rd B: label

brs B
beq }
bne }
b1t } RRB
bge }
bgt }
ble }

- Conditionals
- Loops
- Change the order that instructions are executed

CS52 machine execution

A *program* is simply a sequence of instructions stored in a block of contiguous words in the machine's memory. In executing a program, the CS52 Machine follows a simple loop:

- The machine fetches the value at `mem[ic]` for use as an instruction.
- The machine increments the value in `ic` by 2.
- The machine decodes and carries out the instruction.

Basic structure of CS52 program

```
; great comments at the top!
;
    instruction1    ; comment
    instruction2    ; comment
    ...
label1
    instruction     ; comment
    instruction     ; comment
label2
    ...
    hlt
    end
```

- whitespace before operations/instructions
- labels go here

More CS52 examples

Look at `max_simple.a52`

- Get two values from the user
- Compare them
- Use a branch to distinguish between the two cases
 - Goal is to get largest value in `r3`
- print largest value