CS062

DATA STRUCTURES AND ADVANCED PROGRAMMING

23: Shortest Paths



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Lecture 23: Shortest Paths

- Introduction to Shortest Paths
- API
- Properties
- Dijkstra's Algorithm

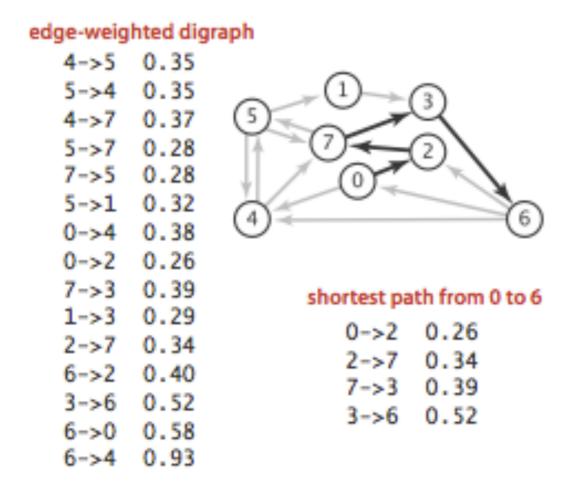
Edge-weighted digraph

Edge-weighted digraph: a digraph where we associate weights or costs with each edge.

edge-weighted digraph 4->5 0.35 5->4 0.35 4->7 0.37 5->7 0.28 7->5 0.28 5->1 0.32 0->4 0.380->2 0.26 7->3 0.39 1->3 0.29 2->7 0.34 6->2 0.40 3->6 0.52 6->0 0.58 6->4 0.93

Shortest Paths

Shortest path from vertex S to vertex t: a directed path from S to t with the property that no other such path has a lower weight (total weight sum of edges it consists).



An edge-weighted digraph and a shortest path

Shortest Path variants

- Single source: from one vertex S to every other vertex.
- Single sink: from every vertex to one vertex t.
- Source-sink: from one vertex S to another vertex t.
- All pairs: from every vertex to every other vertex.

What version is there in your navigation app?

Shortest Paths Assumptions

- Not all vertices need to be reachable.
 - We will assume so in this lecture.
- Weights are non-negative.
 - There are algorithms that can handle negative weights.
- Shortest paths are not necessarily unique but they are simple.

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Weighted directed edge API

- public class DirectedEdge
 - DirectedEdge(int v, int w, double weight)
 - ▶ Constructs a weighted edge from v to w (v->w) with the provided weight.
 - int from()
 - Returns vertex source of this edge.
 - int to()
 - Returns vertex destination of this edge.
 - double weight()
 - Returns weight of this edge.
 - > String toString()
 - ▶ Returns the string representation of this edge.

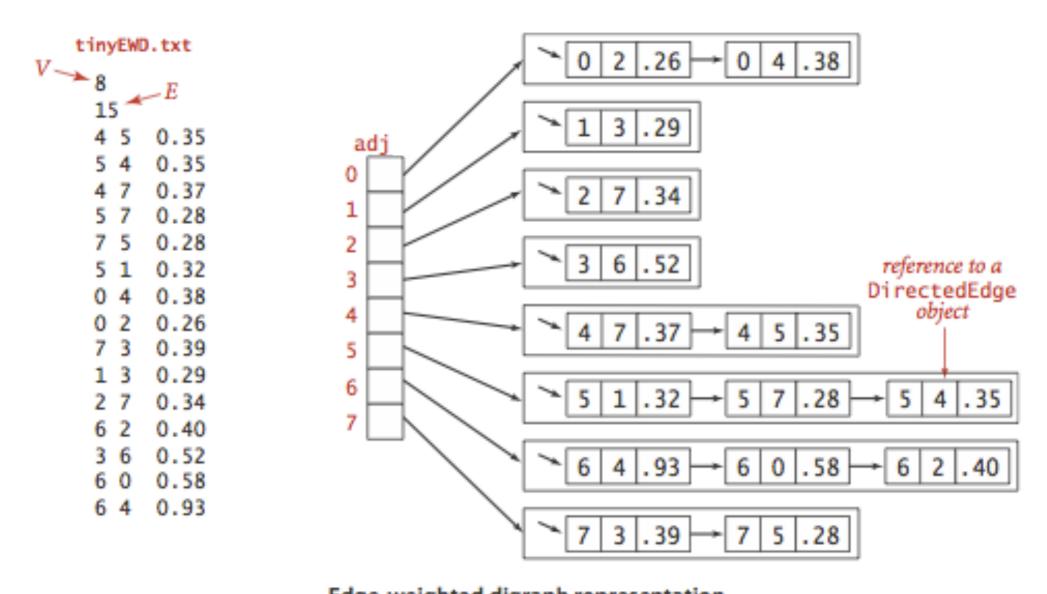
Weighted directed edge in Java

```
public class DirectedEdge {
    private final int v;
    private final int w;
    private final double weight;
   public DirectedEdge(int v, int w, double weight) {
        this.v = v;
        this.w = w;
        this.weight = weight;
    }
   public int from() {
        return v;
    public int to() {
        return w;
    }
    public double weight() {
        return weight;
    }
```

Edge-weighted digraph API

- public class EdgeWeightedDigraph
 - EdgeWeightedDigraph(int v)
 - > Constructs an edge-weighted digraph with v vertices.
 - void addEdge(DirectedEdge e)
 - Add weighted directed edge e.
 - Iterable<DirectedEdge> adj(int v)
 - ▶ Returns edges adjacent from V.
 - int V()
 - ▶ Returns number of vertices.
 - int E()
 - ▶ Returns number of edges.
 - Iterable<DirectedEdge> edges()
 - > Returns all edges.

Edge-weighted digraph adjacency list representation



Edge-weighted digraph representation

Edge-weighted digraph in Java

```
public class EdgeWeightedDigraph {
    private final int V;
                                        // number of vertices in this digraph
    private int E;
                                         // number of edges in this digraph
    private ArrayList<ArrayList<DirectedEdge>> adj; // adj.get(v) = adjacency list for v
    public EdgeWeightedDigraph(int V) {
       this.V = V;
       this.E = 0;
       adj = new ArrayList<ArrayList<DirectedEdge>>(V);
        for (int v = 0; v < V; v++)
            adj.add(new ArrayList<DirectedEdge>());
    public void addEdge(DirectedEdge e) {
        int v = e.from();
        int w = e.to();
        adj.get(v).add(e);
        E++;
    }
   public Iterable<DirectedEdge> adj(int v) {
       return adj.get(v);
```

Single-source shortest path API

- ▶ Goal: find shortest path from S to every other vertex in the digraph.
- public class SP
 - SP(EdgeWeightedDigraph G, int s)
 - Shortest paths from s in digraph G.
 - double distTo(int v)
 - Length of shortest path from S to V.
 - Iterable<DirectedEdge> pathTo(int v)
 - \blacktriangleright Returns edges along the shortest path from S to \lor .
 - boolean hasPathTo(int v)
 - \blacktriangleright Returns whether there is a path from S to \lor .

Lecture 23: Shortest Paths

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Data structures for single-source shortest paths

- Goal: find shortest path from S to every other vertex in the digraph.
- Shortest-paths tree (SPT): a subgraph which will be a directed tree rooted at S which will contain all the vertices reachable from S and every tree path in the SPT is a shortest path in the digraph.
- Representation of shortest paths with two vertex-indexed arrays.
 - Edges on the shortest-paths tree: edgeTo[v] is the last edge on a shortest path from S to V.
 - Distance to the source: distTo[v] is the length of the shortest path from S to V.

```
public Iterable<DirectedEdge> pathTo(int v) {
    Stack<DirectedEdge> path = new Stack<DirectedEdge>();
    for (DirectedEdge e = edgeTo[v]; e != null; e = edgeTo[e.from()]) {
        path.push(e);
    }
    return path;
}
```

```
4->5 0.35

5->4 0.35

4->7 0.37

5->7 0.28

7->5 0.28

5->1 0.32

0->4 0.38

0->2 0.26

7->3 0.39

1->3 0.29

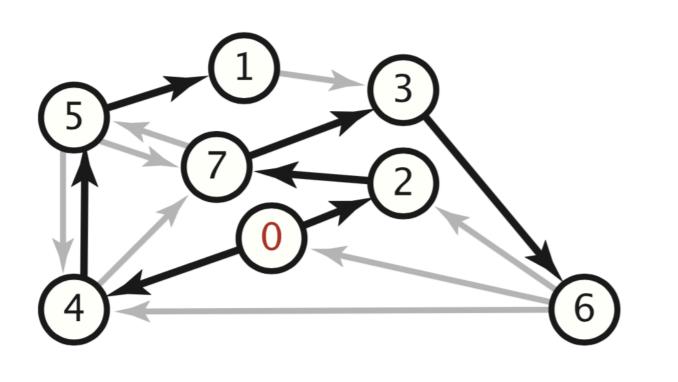
2->7 0.34

6->2 0.40

3->6 0.52

6->0 0.58

6->4 0.93
```

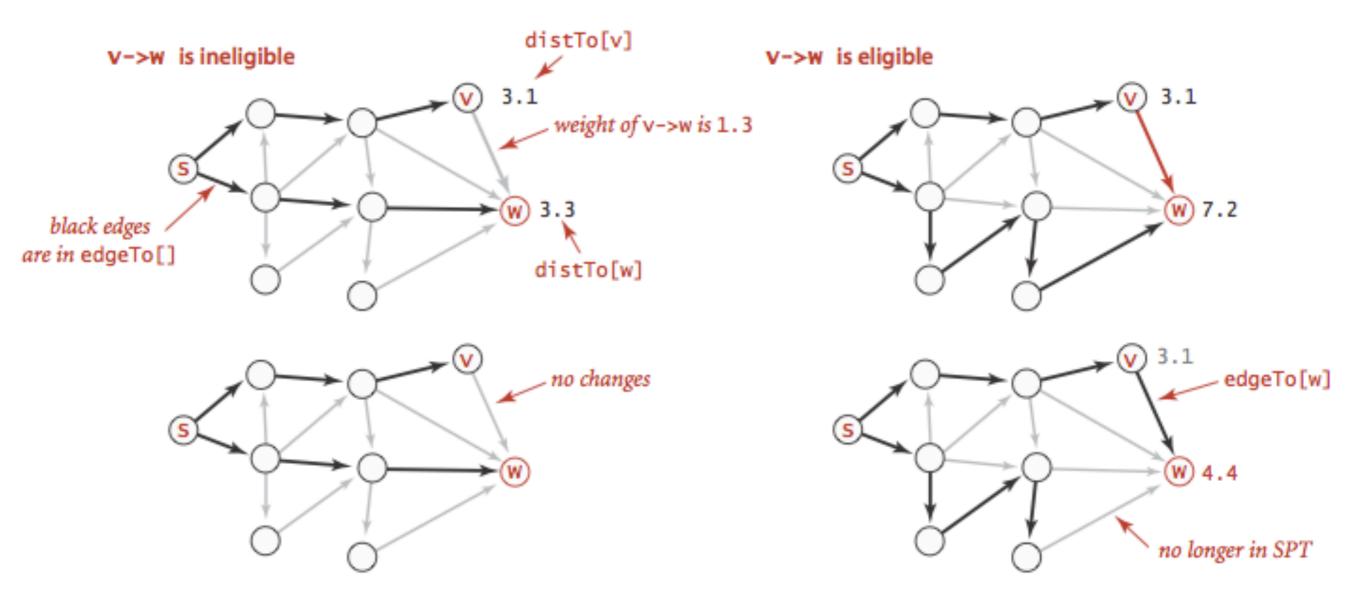


edgeTo[] distTo[] nu11 0 0 5->1 0.32 1.05 2 $0 \rightarrow 2 \quad 0.26$ 0.26 3 7->3 0.39 0.99 4 0.38 $0 - > 4 \ 0.38$ 5 0.73 4->5 0.35 6 $3 - > 6 \ 0.52$ 1.51 0.60 $2 \rightarrow 7 \quad 0.34$

Edge relaxation

- ▶ Relax edge e = V->W
 - distTo[v] is the length of the shortest known path from S to v.
 - distTo[w] is the length of the shortest known path from S to w.
 - edgeTo[w] is the last edge on shortest known path from S to W.
 - If e = v->w yields shorter path to w, update distTo[w] and edgeTo[w].

Edge relaxation



Edge relaxation implementation

```
private void relax(DirectedEdge e) {
  int v = e.from(), w = e.to();
  if (distTo[w] > distTo[v] + e.weight()) {
     distTo[w] = distTo[v] + e.weight();
     edgeTo[w] = e;
  }
}
```

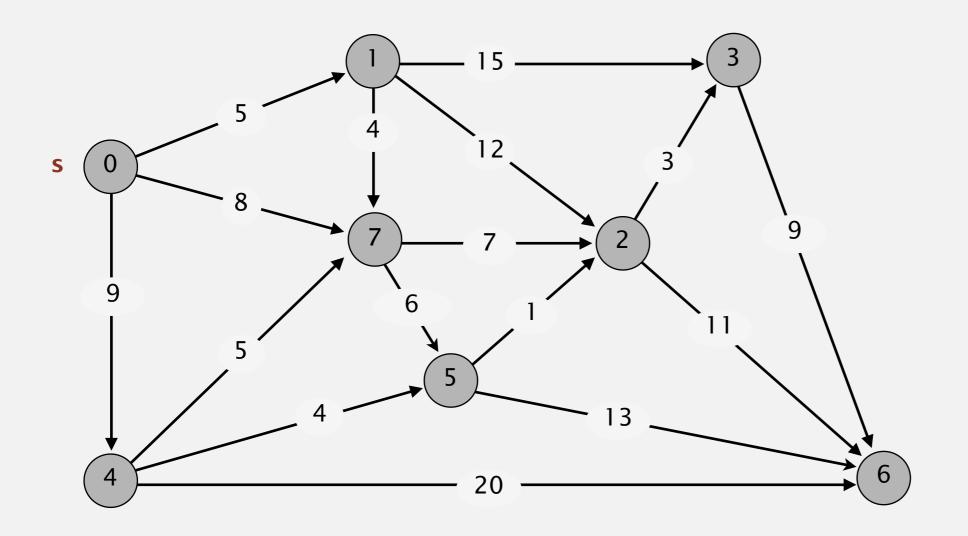
Framework for shortest-paths algorithm

- Generic algorithm to compute a SPT from S
 - ▶ $distTo[v] = \infty$ for each vertex v.
 - edgeTo[v]=null for each vertex v.
 - b distTo[s]=0.
 - Repeat until done:
 - Relax any edge.
- distTo[v] is the length of a simple path from S to V.
- b distTo[v] does not increase.

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- Consider vertices in increasing order of distance from s
 (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



0→7 8.0 1→2 12.0 15.0 1→3 1→7 4.0 2→3 3.0 2→6 11.0 3→6 9.0 4→5 4.0 4→6 20.0 4→7 5.0 1.0 5→2 5→6 13.0 6.0 7→5 7.0 7→2

5.0

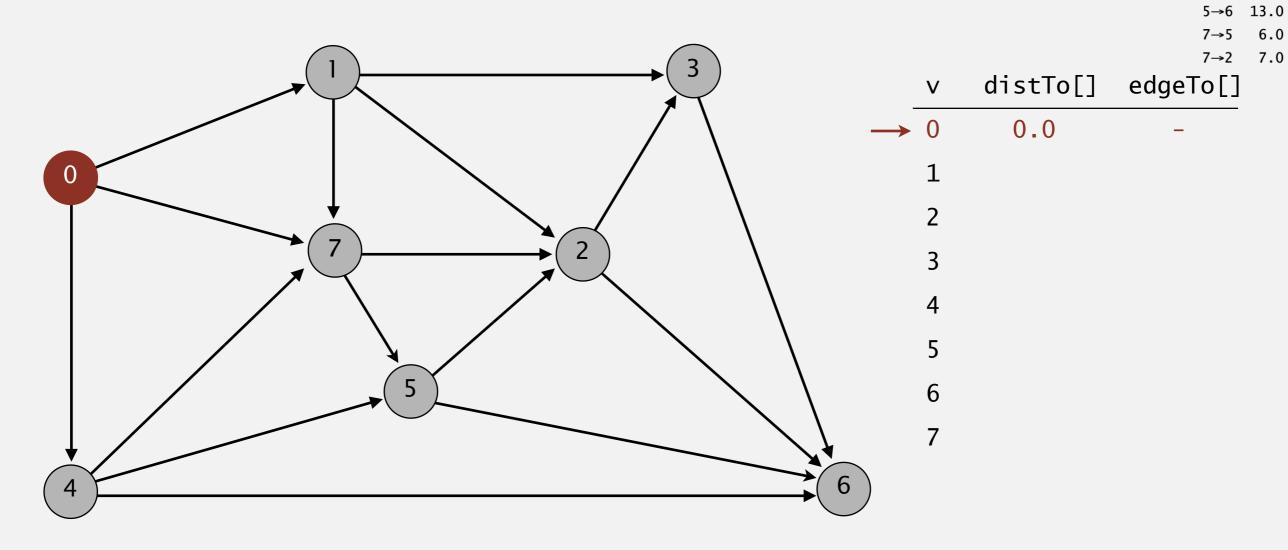
9.0

0→1

0→4

an edge-weighted digraph

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

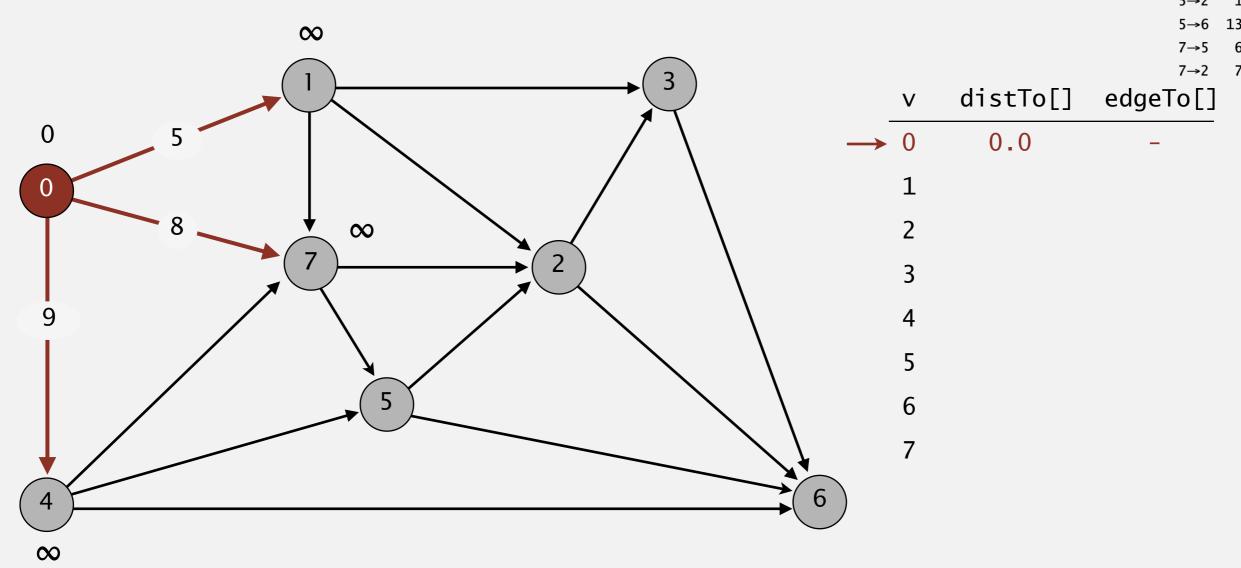


choose source vertex 0

5.0

0→1

- Consider vertices in increasing order of distance from s
 (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.

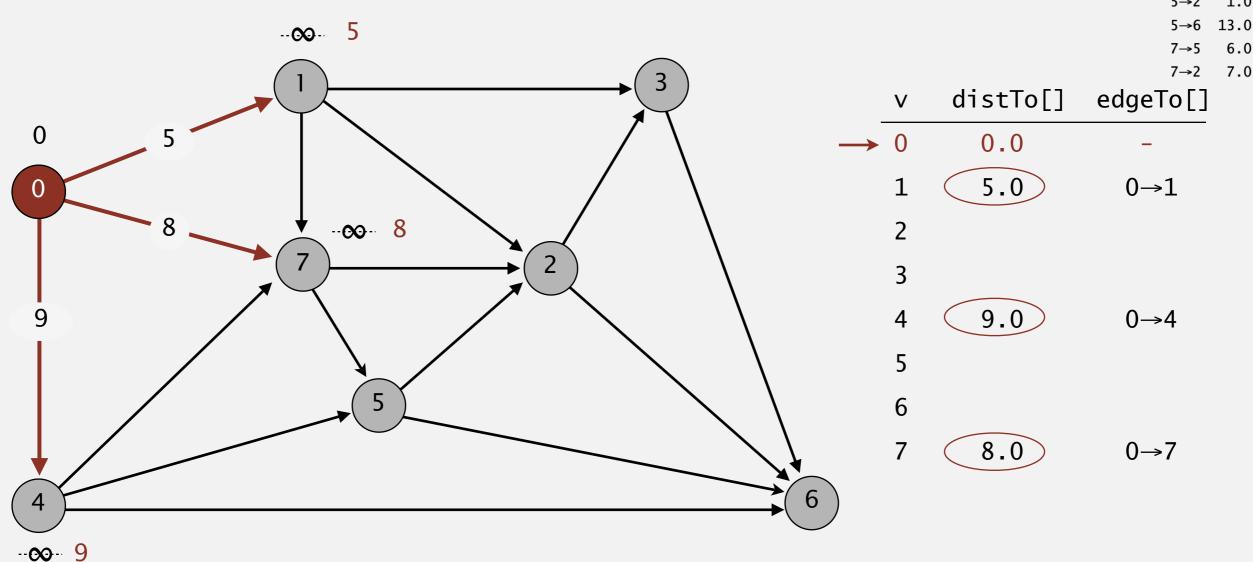


relax all edges adjacent from 0

5.0

0→1

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 0

0→1

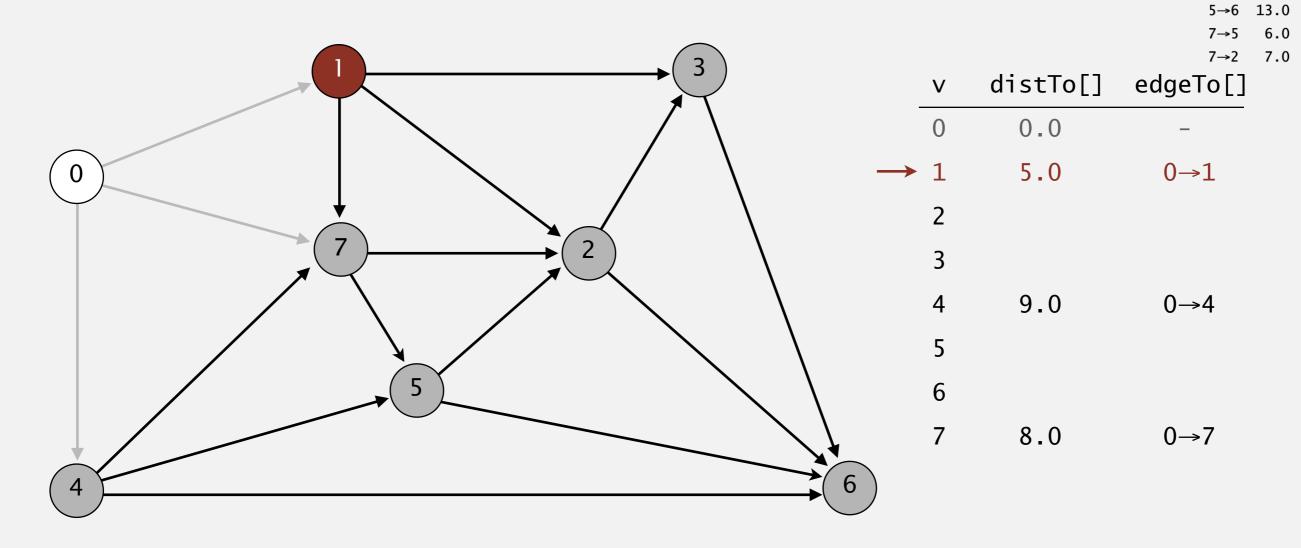
1→2 12.0

5.0

3.0

11.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



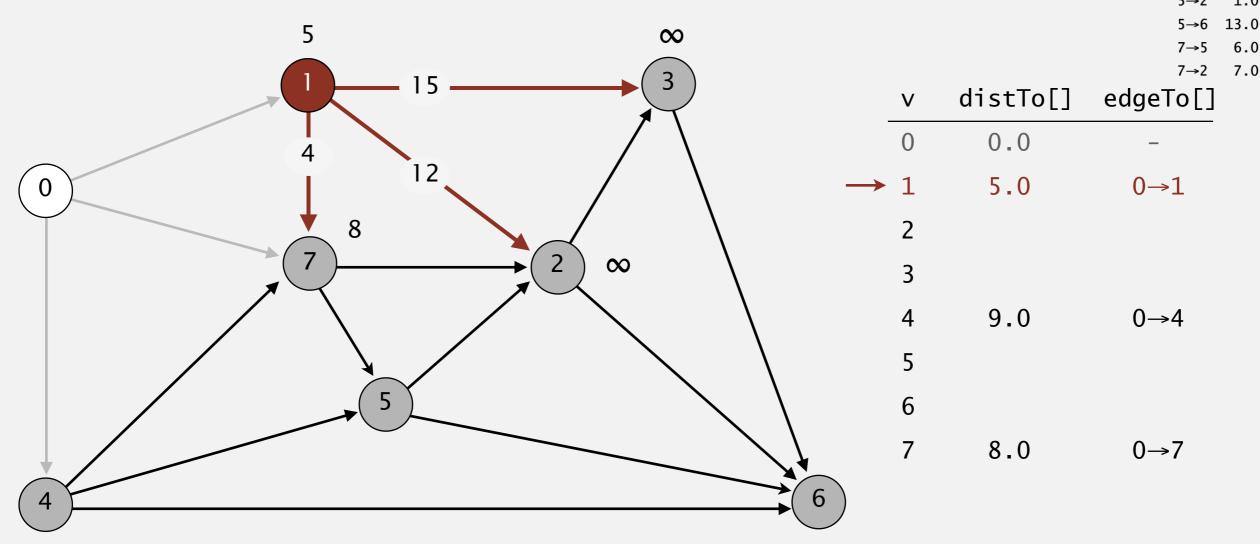
choose vertex 1

0→1

1→2 12.0

5.0

- Consider vertices in increasing order of distance from s
 (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 1

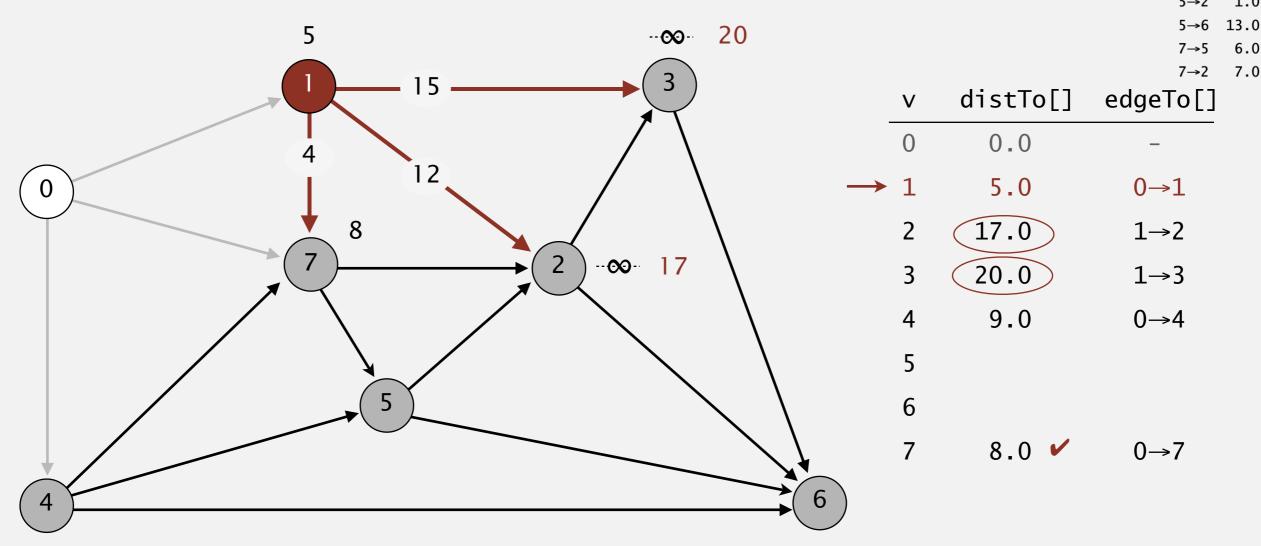
0→1

1→2 12.0

5.0

3.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 1

0→1

5.0

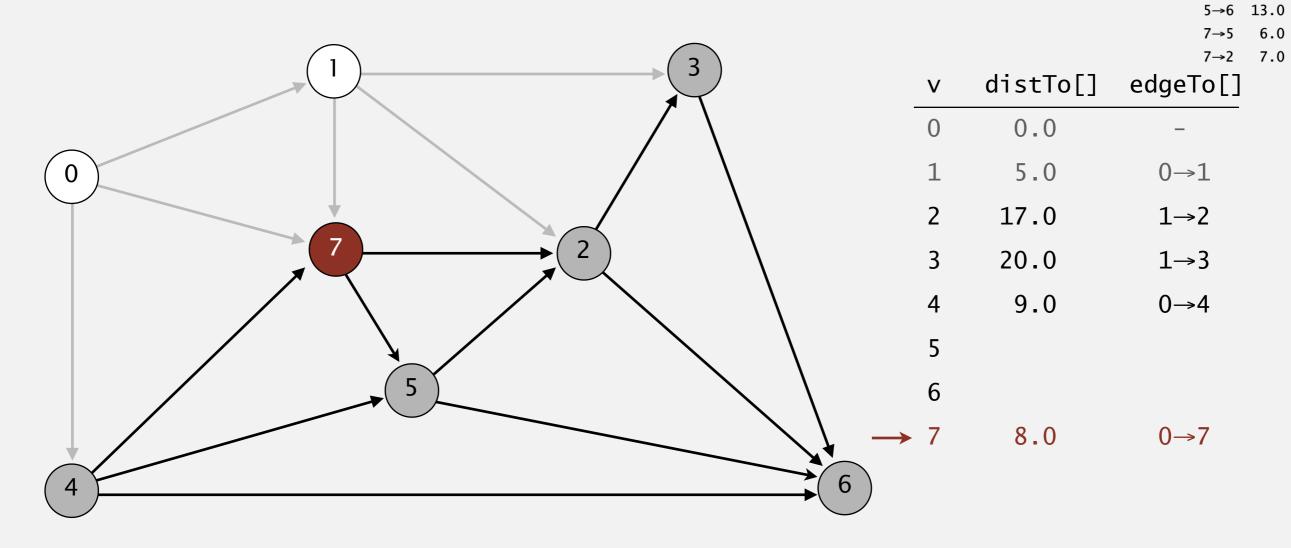
3.0

11.0

20.0

1→2 12.0

- Consider vertices in increasing order of distance from s
 (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



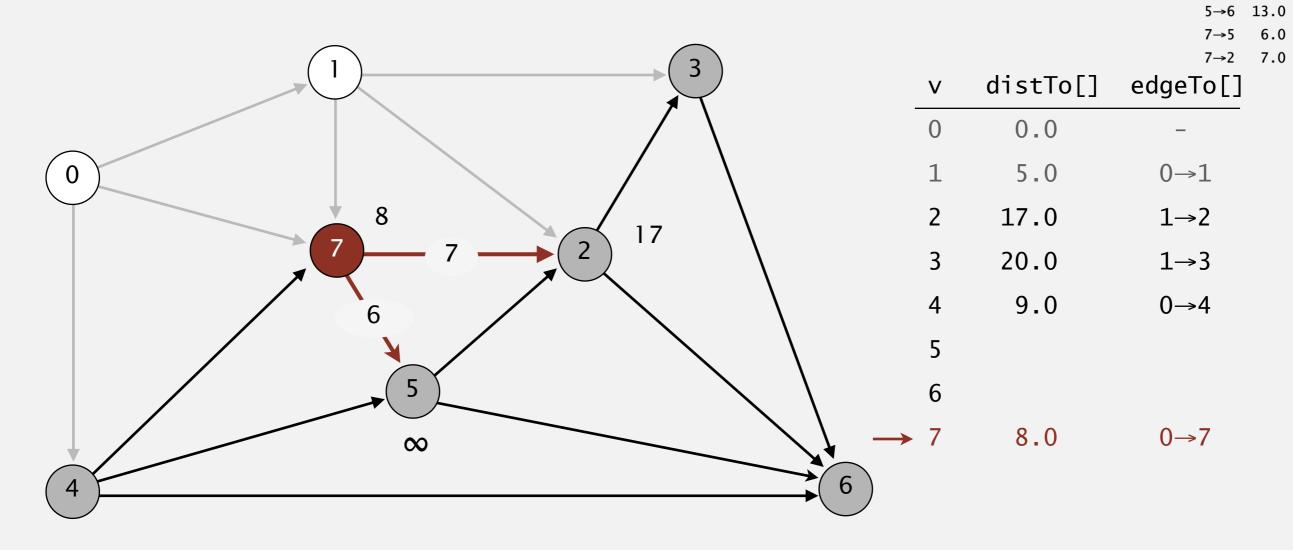
choose vertex 7

0→1

1→2 12.0

5.0

- Consider vertices in increasing order of distance from s
 (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 7

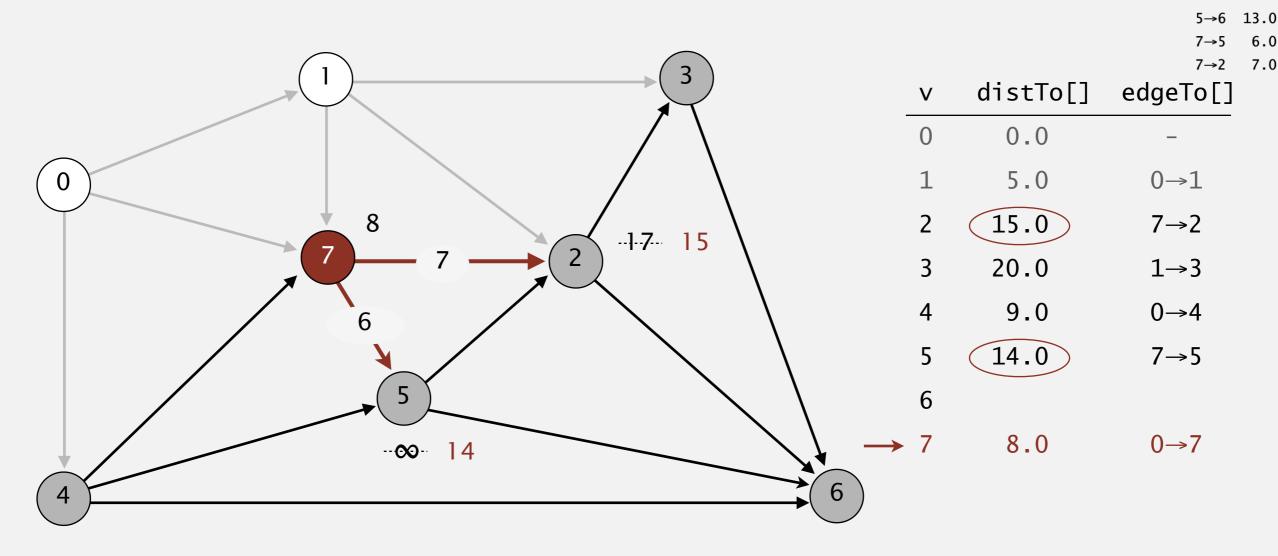
0→1

1→2 12.0

5.0

3.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 7

0→1

5.0

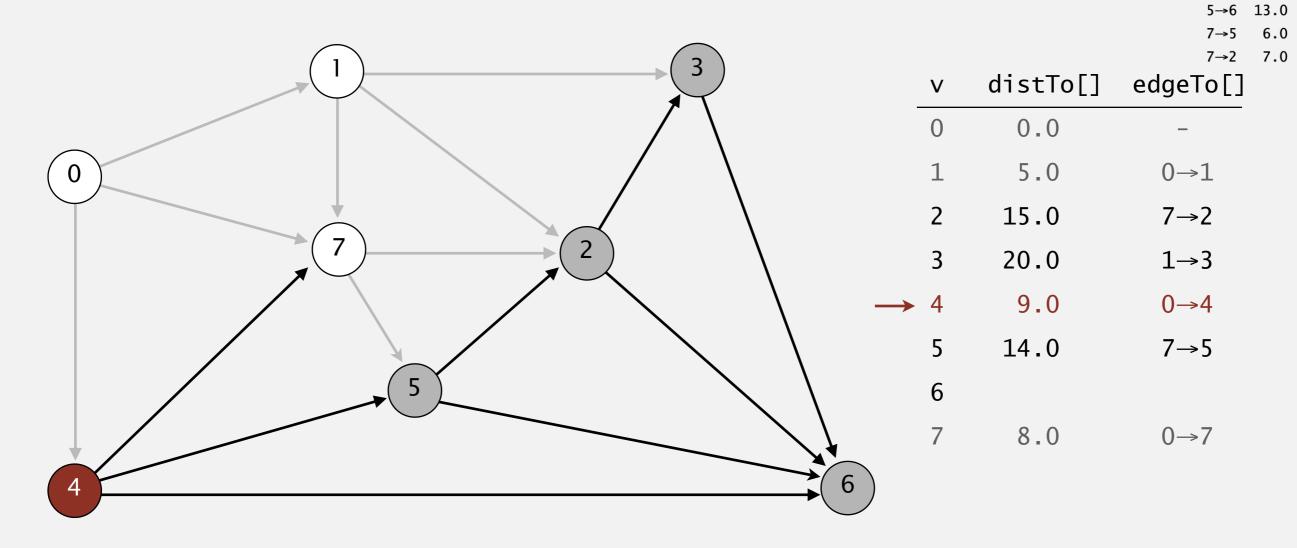
3.0

11.0

20.0

1→2 12.0

- Consider vertices in increasing order of distance from s
 (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



select vertex 4

0→1

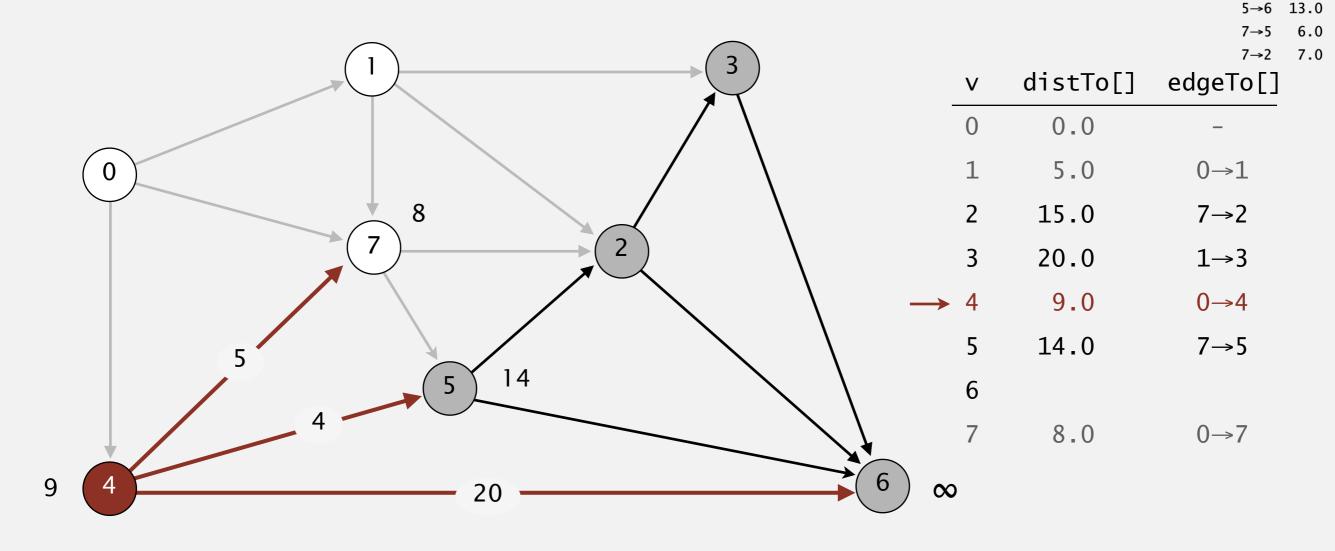
5.0

3.0

20.0

1→2 12.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 4

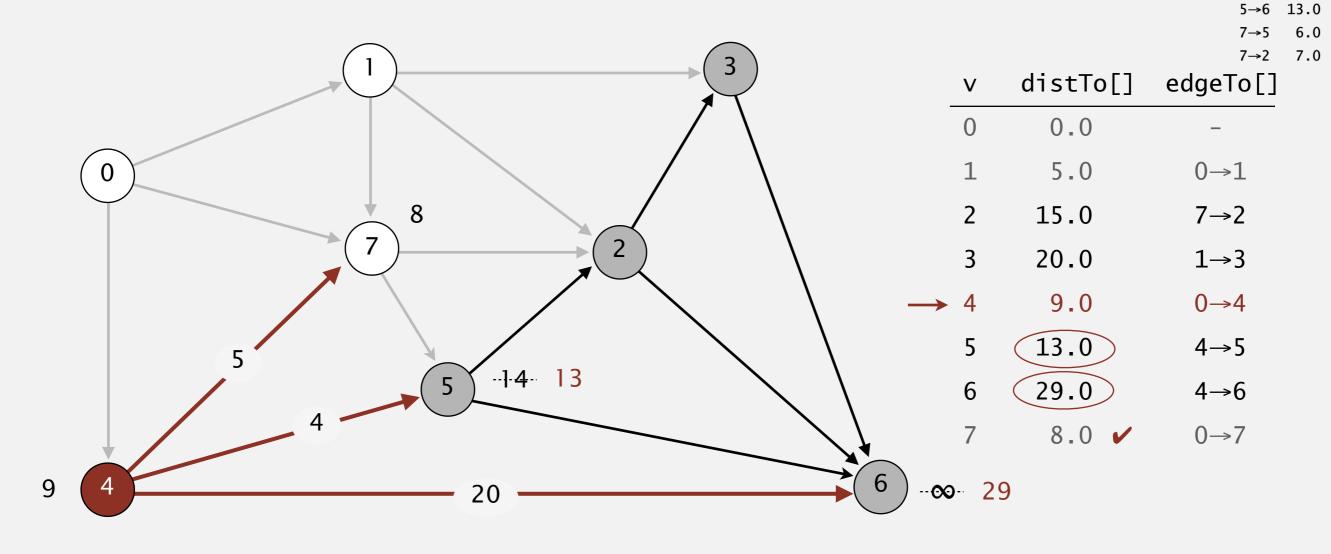
0→1

1→2 12.0

5.0

3.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 4

0→1

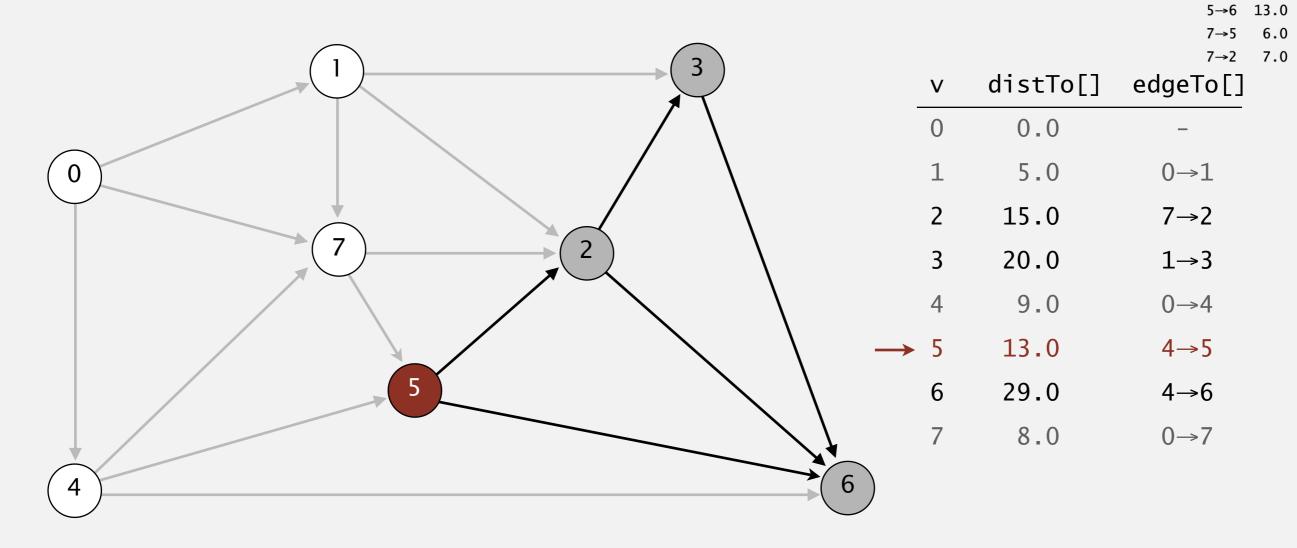
1→2 12.0

5.0

3.0

11.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



select vertex 5

0→1

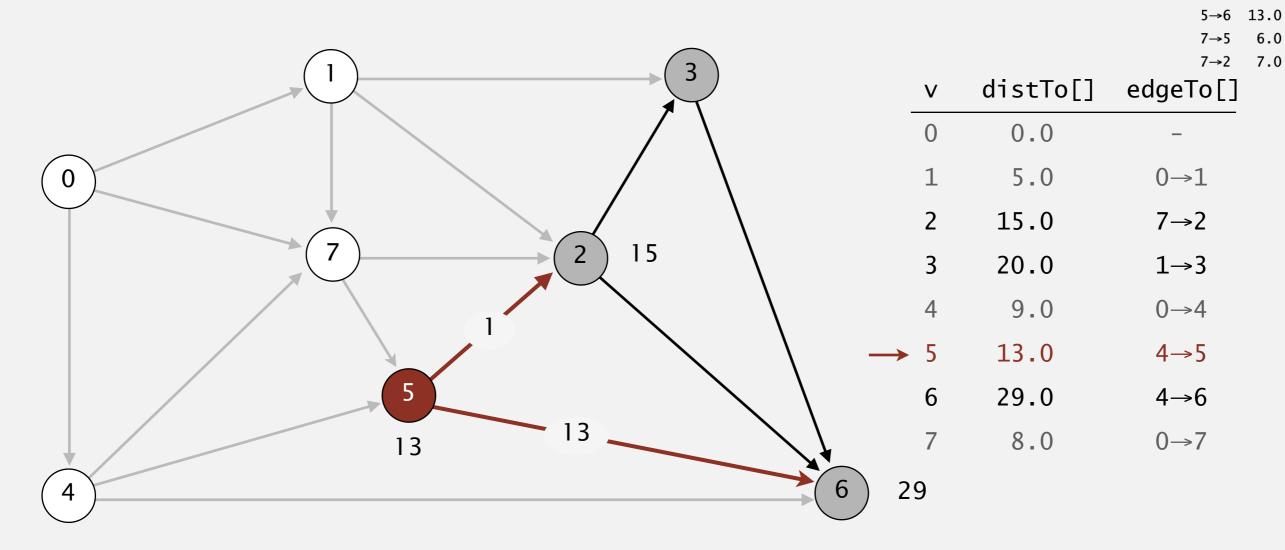
5.0

3.0

20.0

1→2 12.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 5

0→1

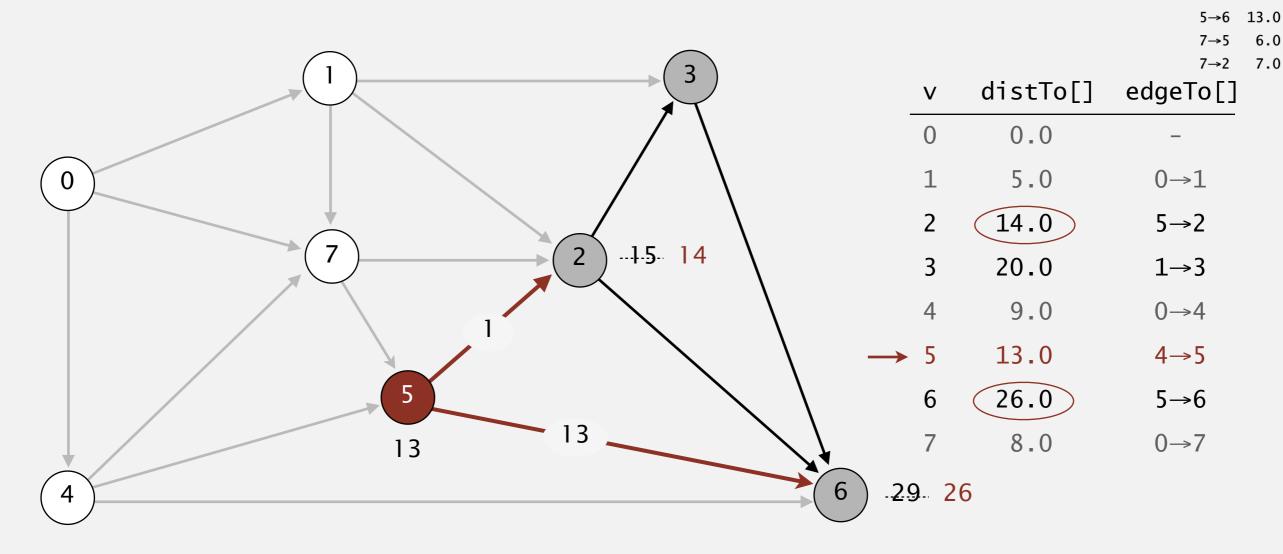
5.0

3.0

20.0

1→2 12.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 5

0→1

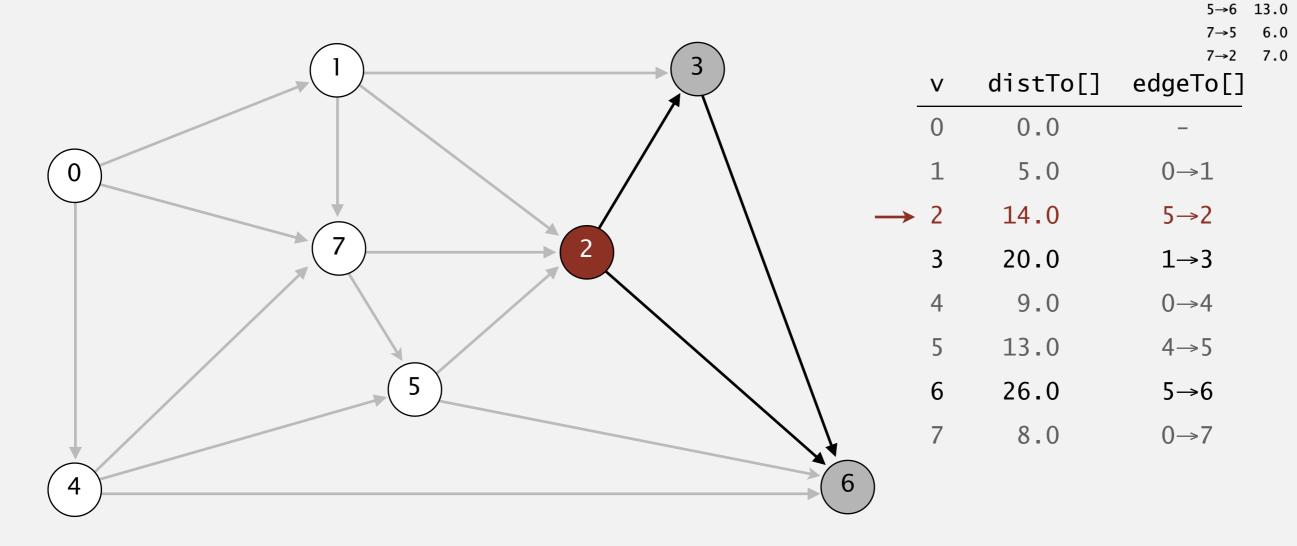
5.0

3.0

11.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



select vertex 2

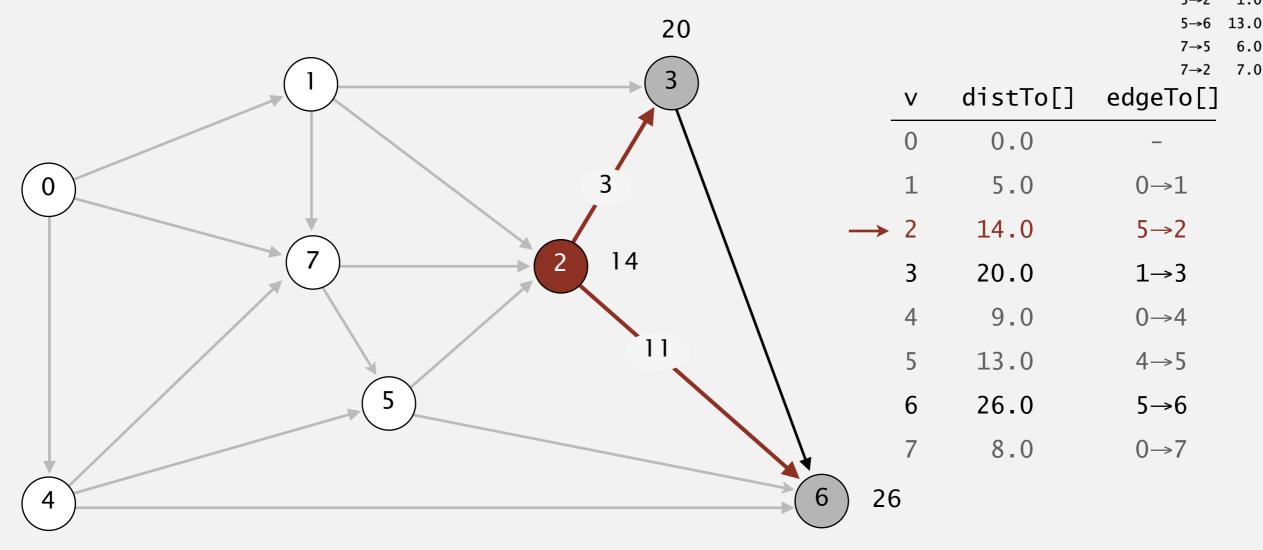
0→1

5.0

3.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 2

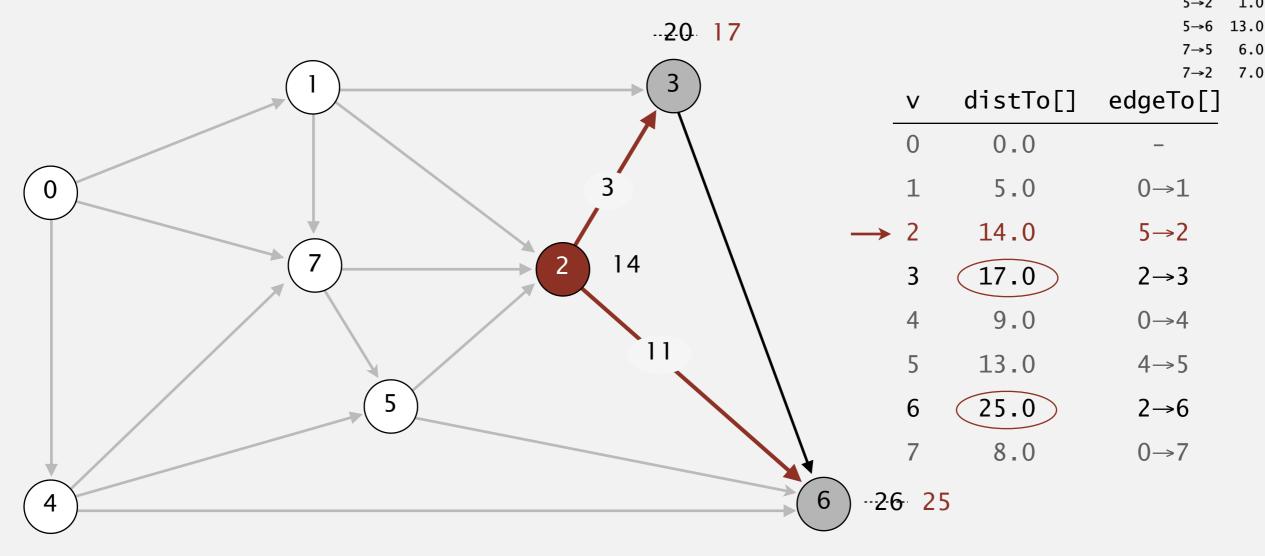
0→1

5.0

3.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 2

0→1

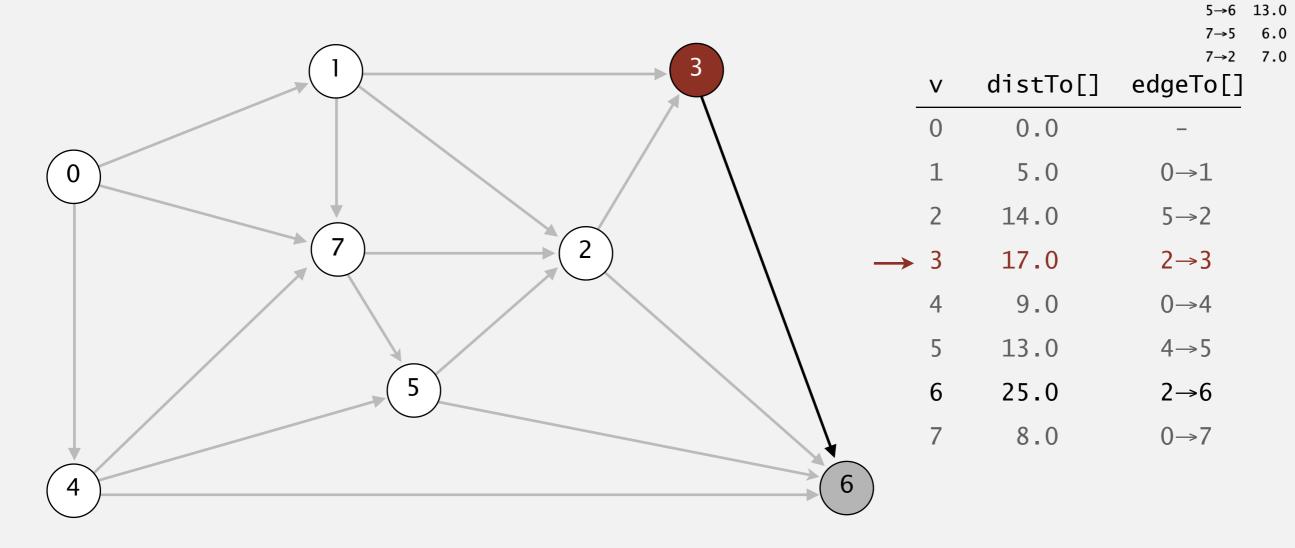
5.0

3.0

11.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



select vertex 3

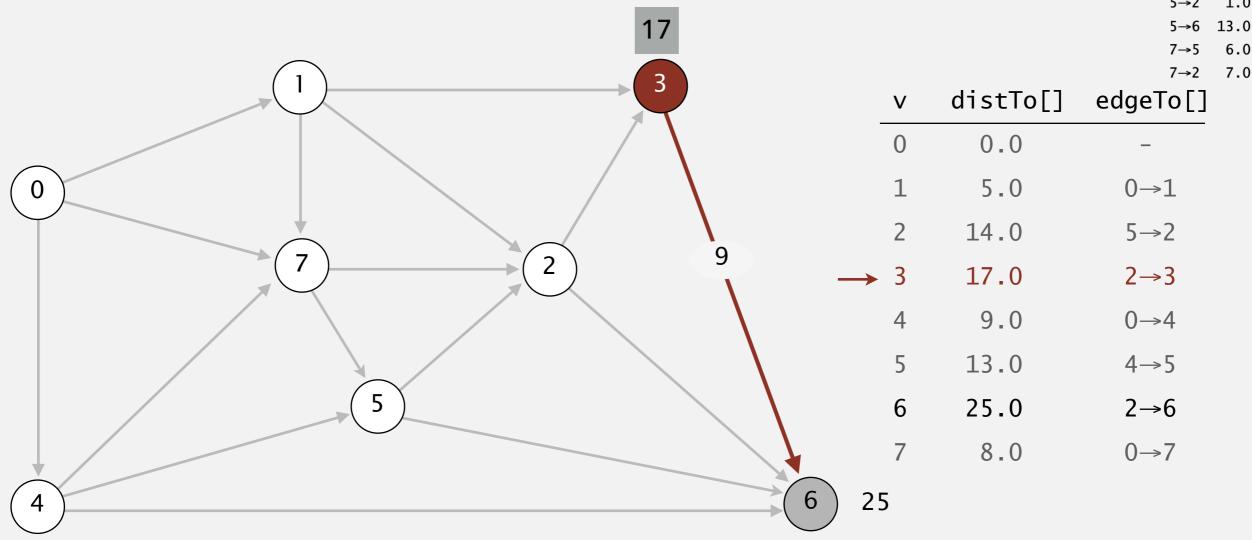
0→1

5.0

3.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 3

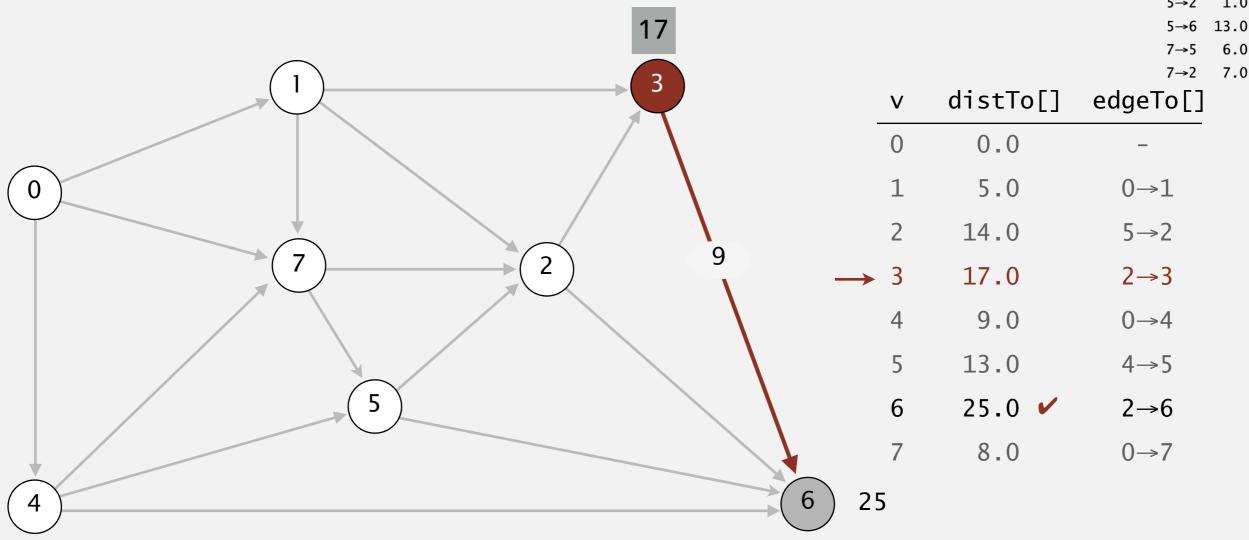
0→1

5.0

3.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 3

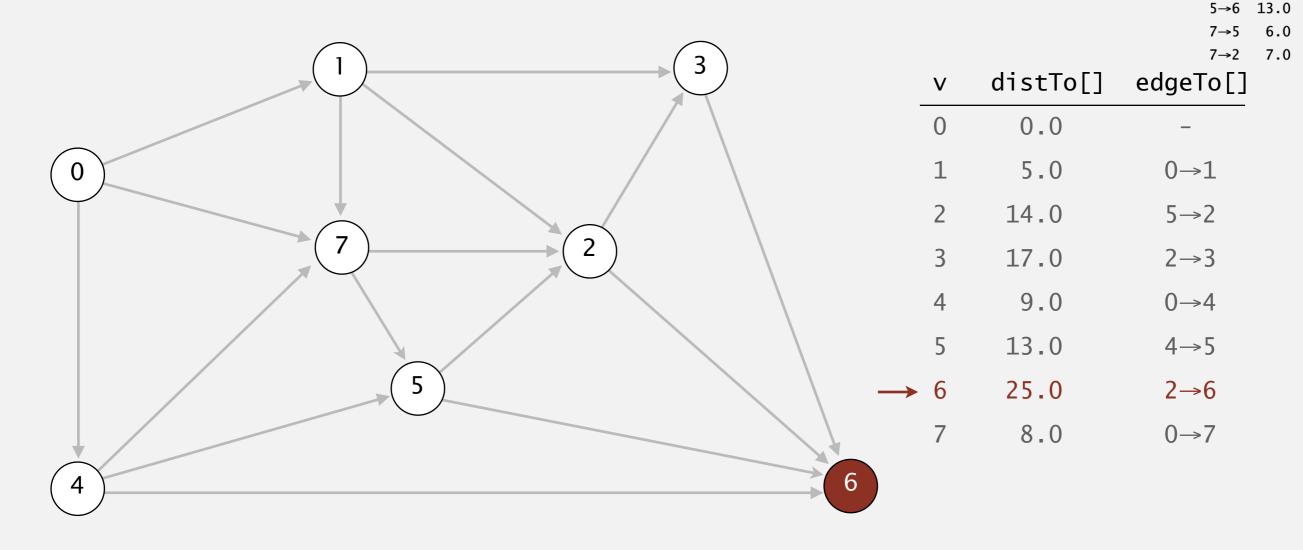
0→1

5.0

3.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



select vertex 6

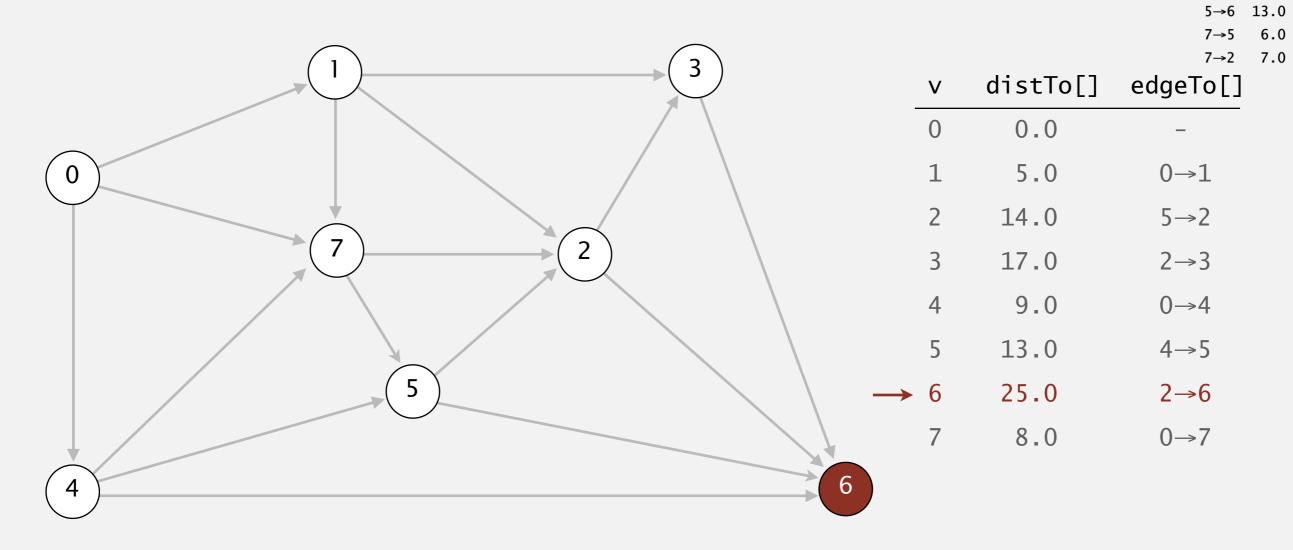
0→1

5.0

3.0

20.0

- Consider vertices in increasing order of distance from s (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



relax all edges adjacent from 6

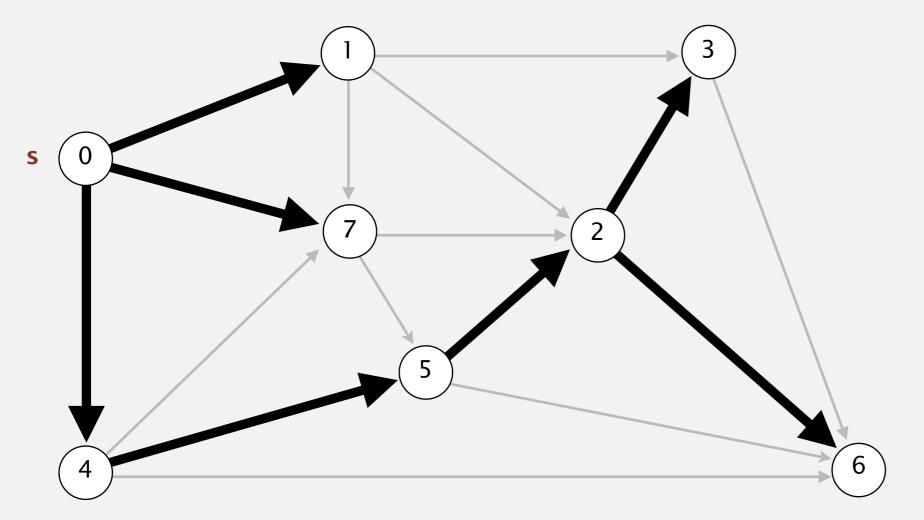
0→1

5.0

3.0

20.0

- Consider vertices in increasing order of distance from s
 (non-tree vertex with the lowest distTo[] value).
- Add vertex to tree and relax all edges adjacent from that vertex.



		7→2
V	distTo[]	edgeTo[]
0	0.0	_
1	5.0	0→1
2	14.0	5→2
3	17.0	2→3
4	9.0	0→4
5	13.0	4→5
6	25.0	2→6
7	8.0	0→7

shortest-paths tree from vertex s

5.0

7.0

Indexed min-priority queue (Section 2.4 in recommended textbook)

- Associate an index between 0 and n-1 with each key in a priority queue.
 - Insert a key associated with a given index.
 - Delete a minimum key and return associated index.
 - Decrease the key associated with a given index.
- public class IndexMinPQ<Key extends Comparable<Key>>
 - IndexMinPQ(int n)
 - ▶ Create indexed PQ with indices 0,1,...n-1
 - void insert(int i, Key key)
 - Associate key with index i.
 - int delMin()
 - Remove a minimal key and return its associated index.
 - void decreaseKey(int i, Key key)
 - Decrease the key with index i to the specified value.

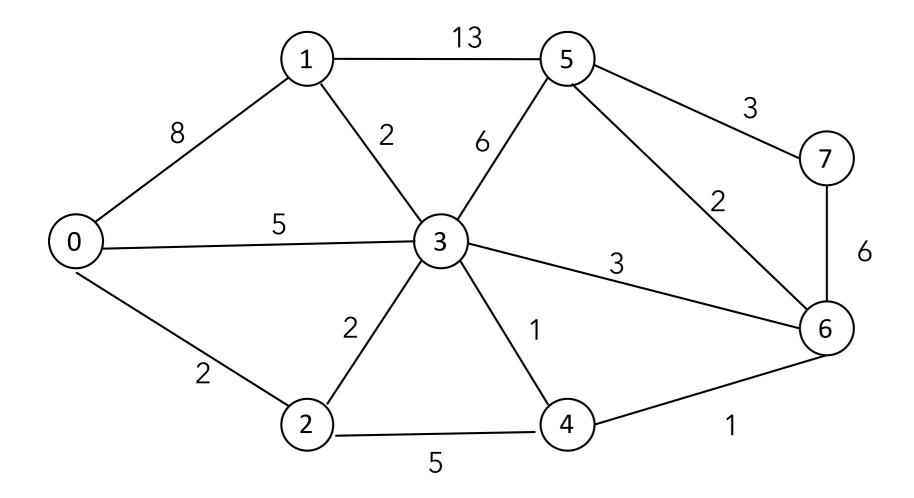
```
public class DijkstraSP {
                                    // distTo[v] = distance of shortest s->v path
   private double[] distTo;
   private DirectedEdge[] edgeTo; // edgeTo[v] = last edge on shortest s->v path
   private IndexMinPQ<Double> pq;
                                      // priority queue of vertices
   public DijkstraSP(EdgeWeightedDigraph G, int s) {
        distTo = new double[G.V()];
        edgeTo = new DirectedEdge[G.V()];
        for (int v = 0; v < G.V(); v++)
           distTo[v] = Double.POSITIVE INFINITY;
        distTo[s] = 0.0;
        // relax vertices in order of distance from s
        pq = new IndexMinPQ<Double>(G.V());
        pq.insert(s, distTo[s]);
       while (!pq.isEmpty()) {
            int v = pq.delMin();
           for (DirectedEdge e : G.adj(v))
                relax(e);
   }
   // relax edge e and update pq if changed
   private void relax(DirectedEdge e) {
        int v = e.from(), w = e.to();
        if (distTo[w] > distTo[v] + e.weight()) {
            distTo[w] = distTo[v] + e.weight();
            edgeTo[w] = e;
           if (pq.contains(w)) pq.decreaseKey(w, distTo[w]);
                                pq.insert(w, distTo[w]);
            else
   }
```

Running time depends on PQ implementation

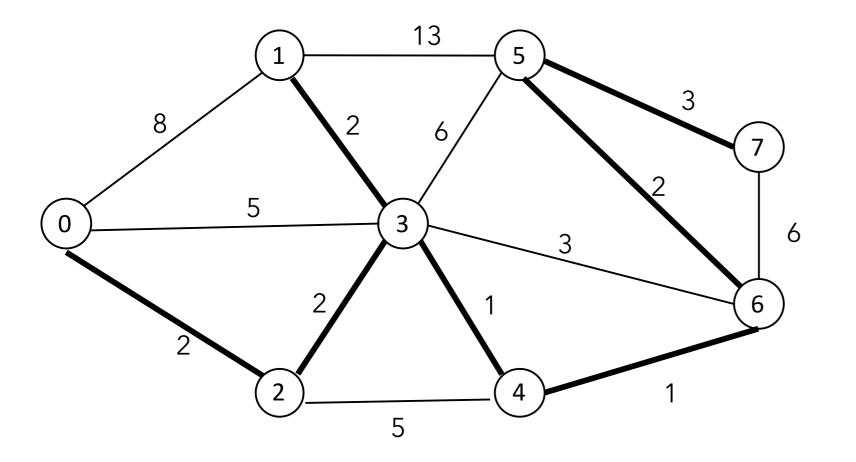
- Many variations. Assuming binary heap, running time is proportional to $|E|\log|V|$ and |V| extra space.
 - lacksquare Cost of insert, delete-min, decrease-key are all $\log |V|$.
- More complicated version with a Fibonacci heap takes $O(|E| + |V| \log |V|)$ time but in practice it's not worth implementing.

Practice Time

▶ Run Dijkstra's algorithm on the following graph with 0 being the starting vertex.



Answer



V	distTo[]	edgeTo[]
0	0	-
1	6	3->1
2	2	0->2
3	4	2->3
4	5	3->4
5	8	6->5
6	6	4->6
7	11	5->7

Lecture 23: Shortest Paths

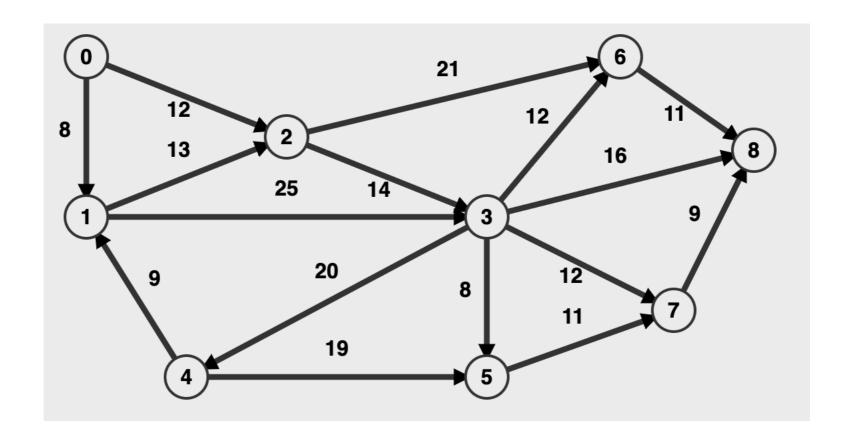
- Introduction to Shortest Paths
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- Properties
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Readings:

- Recommended Textbook: Chapter 4.4 (Pages 638-676)
- Website:
 - https://algs4.cs.princeton.edu/44sp/
- Visualization
 - https://visualgo.net/en/sssp

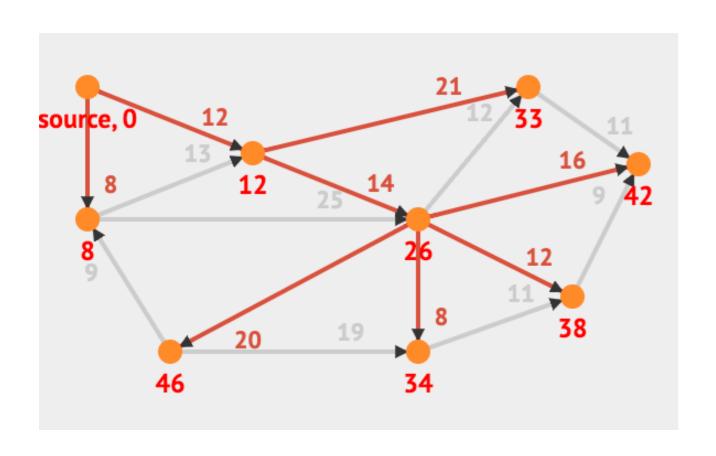
Problem

Run Dijkstra's algorithm on the following graph with 0 being the starting vertex.



Answer

Run Dijkstra's algorithm on the following graph with 0 being the starting vertex.



V	distTo[]	edgeTo[]
0	0	-
1	8	0->1
2	12	0->2
3	26	2->3
4	46	3->4
5	34	3->5
6	33	3->6
7	38	3->7
8	42	3->8