SEARCHING

CS062 DATA STRUCTURES AND ADVANCED PROGRAMMING

20: Left-Leaning Red-Black Trees



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Lecture 20: Left-leaning Red-Black Trees

Introduction

- Elementary red-black BST operations
- Insertion
- Mathematical analysis
- Historical context

Left-leaning red-black BSTs correspond 1-1 with 2-3 trees

- Start with standard BSTs which are made up of 2-nodes.
- Add extra information to encode 3-nodes. We will introduce two types of links.
- Red links: bind together two 2-nodes to represent a 3-node.
 - Specifically, 3-nodes are represented as two 2-nodes connected by a single red link that leans left (one of the 2-nodes is the left child of the other).
- Black links: bind together the 2-3 tree.
- Advantage: Can use BST code with minimal modification.

Left-leaning red-black BSTs correspond 1-1 with 2-3 trees



1-1 correspondence between red-black BSTs and 2-3 trees

Definition

- A left-leaning red-black tree is a BST such that:
 - No node has two red links connected to it.
 - Red links lean left.
 - Every path from root to leaves has the same number of black links (perfect black balance).



Search

- Exactly the same as for elementary BSTs (we ignore the color).
 - But runs faster because of better balance.

```
public Value get(Key key) {
    if (key == null) throw new IllegalArgumentException("argument to get() is null");
    return get(root, key);
}
// value associated with the given key in subtree rooted at x; null if no such key
private Value get(Node x, Key key) {
    while (x != null) {
        int cmp = key.compareTo(x.key);
        if (cmp < 0) x = x.left;
        else if (cmp > 0) x = x.right;
        else return x.val;
    }
    return null;
}
```

Operations such as floor, iteration, rank, selection are also identical.

Representation

- Each node is pointed to by one node, its parent. We can use this to encode the color of the links in nodes.
- True if the link from the parent is red and false if it is black. Null links are black.

```
private static final boolean RED = true;
private static final boolean BLACK = false;
private Node root;
                       // root of the BST
// BST helper node data type
private class Node {
    private Key key;
                               // key
    private Value val;
                               // associated data
    private Node left, right; // links to left and right subtrees
    private boolean color;
                               // color of parent link
    private int size;
                               // subtree count
private boolean isRed(Node x) {
    if (x == null) return false;
    return x.color == RED;
}
```



Practice Time

Which of the following are legal LLRB trees?



Answer

- Which of the following are legal LLRB trees?
- iii and iv
 - i is not balanced and ii is also not in symmetrical order



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Left rotation: Orient a (temporarily) right-leaning red link to lean left





Left rotate (right link of h)

Right rotation: Orient a left-leaning red link to a (temporarily) lean right



Right rotate (left link of h)

Color flip: Recolor to split a (temporary) 4-node



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Basic strategy: Maintain 1-1 correspondence with 2-3 trees

- During internal operations, maintain:
 - symmetric order
 - perfect black balance.
- But we might violate color invariants. For example:
 - Right-leaning red link.
 - Two red children (temporary 4-node).
 - Left-left red (temporary 4-node).
 - Left-right red (temporary 4-node).
- To restore color invariant we will be performing rotations and color flips.

Insertion into a LLRB

- Do standard BST insertion and color the new link red.
- Repeat until color invariants restored:
 - Both children red? Flip colors.
 - Right link red? Rotate left.
 - Two left reds in a row? Rotate right.

Red-black BST construction demo



Implementation

Only three cases:

- Right child red; left child black: rotate left.
- Left child red; left-left grandchild red: rotate right.
- Both children red: flip colors.

```
// insert the key-value pair in the subtree rooted at h
 private Node put(Node h, Key key, Value val) {
     if (h == null) return new Node(key, val, RED, 1);
     int cmp = key.compareTo(h.key);
     if
             (cmp < 0) h.left = put(h.left, key, val);</pre>
     else if (cmp > 0) h.right = put(h.right, key, val);
                       h.val
     else
                               = val;
     // fix-up any right-leaning links
     if (isRed(h.right) && !isRed(h.left))
                                                 h = rotateLeft(h);
     if (isRed(h.left) && isRed(h.left.left)) h = rotateRight(h);
     if (isRed(h.left) && isRed(h.right))
                                                 flipColors(h);
     h.size = size(h.left) + size(h.right) + 1;
     return h;
 }
```



Visualization of insertion into a LLRB tree

> 255 insertions in ascending order.

Visualization of insertion into a LLRB tree

> 255 insertions in descending order.

Visualization of insertion into a LLRB tree

> 255 insertions in random order.

Practice Time - Worksheet #20

Draw the LLRB tree that results when you insert the keys 10, 18, 7, 15, 16, 30, 25, 40, 60 in that order in an initially empty tree. ANSWER - Worksheet #20

Draw the LLRB tree that results when you insert the keys 10, 18, 7, 15, 16, 30, 25, 40, 60 in that order in an initially empty tree.



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Balance in LLRB trees

- Height of LLRB trees is $\leq 2 \log n$ in the worst case.
- Worst case is a 2-3 tree that is all 2-nodes except that the left-most path is made up of 3-nodes.
- All ordered operations (min, max, floor, ceiling) etc. are also O(log n).

Summary for dictionary operations

	Worst case			Average case		
	Search	Insert	Delete	Search	Insert	Delete
BST	п	п	п	log n	log n	\sqrt{n}
2-3 search tree	log n	log n	log n	log n	log n	log n
Red-black BSTs	log n	log n	log n	log n	log n	log n

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Red-black trees

- A dichromatic framework for balanced trees. [Guibas and Sedgewick, 1978].
- Why red-black? Xerox PARC had a laser printer and red and black had the best contrast...
- Left-leaning red-black trees [Sedgewick, 2008]
 - Inspired by difficulties in proper implementation of RB BSTs.
- RB BSTs have been involved in lawsuit because of improper implementation.

Balanced trees in the wild

- Red-black trees are widely used as system dictionaries.
 - e.g., Java: java.util.TreeMap and java.util.TreeSet.
- Other balanced BSTs: AVL, splay, randomized.
- 2-3 search trees are a subset of b-trees.
 - See recommended textbook for more.
 - B-trees are widely used for file systems and databases.

Readings:

- Recommended Textbook: Chapter 3.3 (Pages 424-447)
- Website:
 - https://algs4.cs.princeton.edu/33balanced/
- Visualization:
 - <u>https://algs4.cs.princeton.edu/GrowingTree/</u> (for LLRB trees)
- Worksheet:
- Lecture 20 worksheet

Problem 1

Draw the left-leaning red-black BST that results when you insert items with the keys E, A, S, Y, Q, U, T, I, O, N in that order into an initially empty tree.

ANSWER 1

Draw the left-leaning red-black BST that results when you insert items with the keys E, A, S, Y, Q, U, T, I, O, N in that order into an initially empty tree.

