CS062 DATA STRUCTURES AND ADVANCED PROGRAMMING

10: Doubly Linked Lists



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Now that we've (hopefully) understood how singly linked lists work, let's see a very close linear data structure called doubly linked lists.



We will start by seeing how we would go about implementing doubly linked lists and we'll finish with the default Java implementation.



As with singly linked lists, we can recursively define doubly linked lists as either being empty (null) or a node having a reference to a doubly linked list. A node is a data type that holds any kind of data (think generics) but instead of one reference it now has two, one to the previous and one to the next node. We will use arrows to symbolize the links, rectangles for the nodes, and we will use slashes to indicate the first node called head/beginning/front/or first and the last node, called tail/end/back and last.



Nodes will be again represented through a private inner class that contains an element of type E and two reference to the previous and next node.

```
ZRAYLIST
public interface List <E> {
    void add(E element);
    void add(int index, E element);
    void clear();
    E get(int index);
    boolean isEmpty();
    E remove(int index);
    E set(int index, E element);
    int size();
}
```

Let's refresh our memory about the List interface. If we implement it, we promise to implement the methods:

void add(E element);

void add(int index, E element);

void clear();

E get(int index);

boolean isEmpty();

E remove();

E remove(int index);

E set(int index, E element);

Standard Operations

- DoublyLinkedList(): Constructs an empty doubly linked list.
- * isEmpty():Returns true if the doubly linked list does not contain any element.
- size(): Returns the number of elements in the doubly linked list.
- * E get(int index): Returns the element at the specified index.
- addFirst(E element): Inserts the specified element at the head of the doubly linked list.
- * addLast(E element): Inserts the specified element at the tail of the doubly linked list.
- * add(E element): Inserts the specified element at the tail of the doubly linked list.
- * add(int index, E element): Inserts the specified element at the specified index.
- E set(int index, E element): Replaces the specified element at the specified index and returns the old element
- * E removeFirst(): Removes and returns the head of the doubly linked list.
- * E removeLast(): Removes and returns the tail of the doubly linked list.
- * E remove(): Removes and returns the head of the doubly linked list.
- * E remove(int index): Removes and returns the element at the specified index.
- clear(): Removes all elements.

These are the standard operations we expect to have. We will have a constructor and usual methods for checking the size, whether it is empty, a getter, three adds, one set, three removes, and one clear.

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DoublyLin	<edlist(): an="" constructs="" dll<="" empty="" th=""><th></th></edlist():>	
	head	
	tail	
	size	
What sho	uld happen?	
DoublyLi	inkedList <string> dll = new DoublyLinkedList<string>();</string></string>	

Let's say someone creates a doubly linked lists of strings. what do you think should happen to the head, tail, and size?

DOUBLY LINKED LISTS		8
DoublyLinkedList((): Constructs ar	n empty DLL
DoublyLinkedList <st< td=""><td>ring> dll = new Doubl</td><td>yLinkedList<string>();</string></td></st<>	ring> dll = new Doubl	yLinkedList <string>();</string>
	head = null	
	tail = null	What should happen?
	size = 0	dll.add("CS062");

the head and tail will be null and the size zero. What would happen if we call dll.add("CS062");



The head and tail point to the node that contains CS062 and the size is 1. What if we call dll.addFirst("ROCKS");



The addition will happen at the head. The head now points to the node that contains ROCKS while the tail points to cs062. The size is 2. What should happen if we type dll.addLast("!");



The addition happens at the end. The head remains pointing to the node that contains ROCKS. The insertion creates a new node that contains ! and the tail is moved to it. The size is now 3. What should happen if we call dll.add(1,"?");



We will make room for a new node to be created at index 1 and increase the size by 1. What if we call remove?



remove removes and returns the old head of the doubly linked list while moving the head pointer to the next node. And of course reduces the size by 1. removeFirst works exactly the same way. How about removeLast?



The removal happens at the tail and the size reduces by 1. How about remove(1)?



The node at index 1 (second node) will be removed and the size will be reduced by 1.



Our own implementation of doubly linked lists will lead us to work with generics. we will use the list interface and an inner class for nodes.



That means that we will have three instance variables, head and tail of type Node, and size of type int along with our inner private class for Node.



isEmpty can either check whether the head and tail is null or the size 0. size is very simple.



Knowing what we know about pointers, let's try to implement get.



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The get method will check that the index is within bounds and if not will throw an exception. We will next use the usual trick: we will create a reference that points to where head points to (NOT A NEW NODE!) We will move index steps to the right by pointing finger to finger.next. Eventually, when finger points to the right node, we will return the element it holds.



Let's try addFirst to add an element to the head of the doubly linked list.



Did you get something similar?

PRACTICE TIME: Insert element at tail of doubly linked list



How about adding to the tail?



It should look familiar.



As a note, add is just a call to addLast.

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PRACTICE TIME: Insert element at a	specified index
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/** * Inserts the specified element at the specified ir *	dex.
<pre>* @param index * @param index * the index to insert the element * @param element * the element to insert * @pre 0<=index<=size */</pre>	
<pre>public void add(int index, E element) {</pre>	
<pre>// if index is 0, call addFirst</pre>	
<pre>// if index is size, call addLast</pre>	
<pre>// else // Make two new Node references, previ</pre>	ous and finger. Set previous to null and finger to head
// search for index- \underline{th} position. Set previo	us to finger and move finger to next position
<pre>// create new Node, update its element, and are</pre>	fix its pointers taking into account where finger and previous
// increase number of nodes	
}	

Let's try to add at a specific index.



This is more work. We will need to double down on our trick and have two pointers. Let's call them previous and finger. Finger will start at the head and previous will trail it one step behind. Eventually, finger will reach the index we want to insert the new node which we will reference with current. We will use these two pointers, to make the previous point to current (and vice versa), and current point to finger (and vice versa).



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Replacing an element a specified index will look exactly the same as with singly linked lists.



Let's try to remove the head.



Retrieve and remove head



Did you get something like this? Don't forget that removing the last node is an edge case we need to handle by fixing the tail.



How about removeLast?

DOUBLY LINKED LISTS

Retrieve and remove tail



Very similar idea.



remove is just a call to removeFirst.

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PRACTICE TIME: Retrieve and remove element from a specific index

/**	
* Rem	oves and returns the element at the specified index.
*	
* @pa	ram index
*	the index of the element to be removed
* @re	turn the element previously at the specified index
* @pr	e Ø<=index <size< th=""></size<>
*/	
public	E remove(int index) {
	// check whether index is valia
	// if index is 0
	// return removeFirst
	<pre>// else if index is size-1</pre>
	// return removeLast
	// else
	// Make two new Node references, previous and finger. Set previous to null and finger to head
	// search for index- \underline{th} position. Set previous to finger and move finger to next position
	// update pointers for previous and finger
	// decrease number of nodes
	// return the element that finger points to
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-	
}	

What about removing at a specific index?



We will use again the same trick to find the node at the index we want to remove it.



Clear is super simple. Just set the head and tail to null and the size to 0. The garbage collector will take care of the rest.



Let's look now into the running time complexity of addFirst. It will be O(1). It does not depend on how many elements already exist in the doubly linked list. The fact that we need to do a couple of operations doesn't matter. they don't scale linearly with the size of the doubly linked list.



Same idea for addLast (and as a consequence for add)



Get is another story. It can take O(n) for worst case if we need to hop n steps to find the desired index.



same idea for add, worst case is O(n).

set(int index, E element) in singly linked lists is O(n) for worst case



and for set.



removeFirst (and remove) from the head in contrast is O(1) like with addFirst.



Same idea for removeLast



But remove at a specific index can be O(n)





clear() in singly linked lists is O(1) for worst case

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Clear is O(1)!



That's all for our own implementation. Let's see Java's default implementation



LinkedList also implements the list interface like array list and vector.



If you want to use it, you will have to import the java.util.LinkedList;



And that's all for today



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Readings:



Feel free to download the code and play with our implementation.