

Lecture 15: Capabilities

CS 181S

November 5, 2018

Where we were...

- **Authentication:** mechanisms that bind principals to actions
- **Authorization:** mechanisms that govern whether actions are permitted
 - Discretionary Access Control
 - Mandatory Access Control



Access Control Policy

- An **access control policy** specifies which of the **operations** associated with any given **object** each **principal** is authorized to perform
- Expressed as a relation *Auth*:

<i>Auth</i>	Objects	
	dac.tex	dac.pptx
principals	eibirrell	r,w
	faculty	r
	student	r

Capability Lists

Access Control Lists

Capability Lists

- The capability list for a principal P is a list
$$\langle O_1, Privs_1 \rangle, \langle O_2, Privs_2 \rangle, \dots, \langle O_n, Privs_n \rangle$$
 - e.g., $\langle \text{dac.tex}, \{r,w\} \rangle \langle \text{dac.pptx}, \{r,w\} \rangle$
- **Capabilities** carry privileges.
 - 1) **Authorization:** Performing operation op on object O_i requires a principal P to hold a capability $C_i = \langle O_i, Privs_i \rangle$ such that $op \in Privs_i$
 - 2) **Unforgeability:** Capabilities cannot be counterfeited or corrupted.
- Note: Capabilities are (typically) transferable

Capabilities

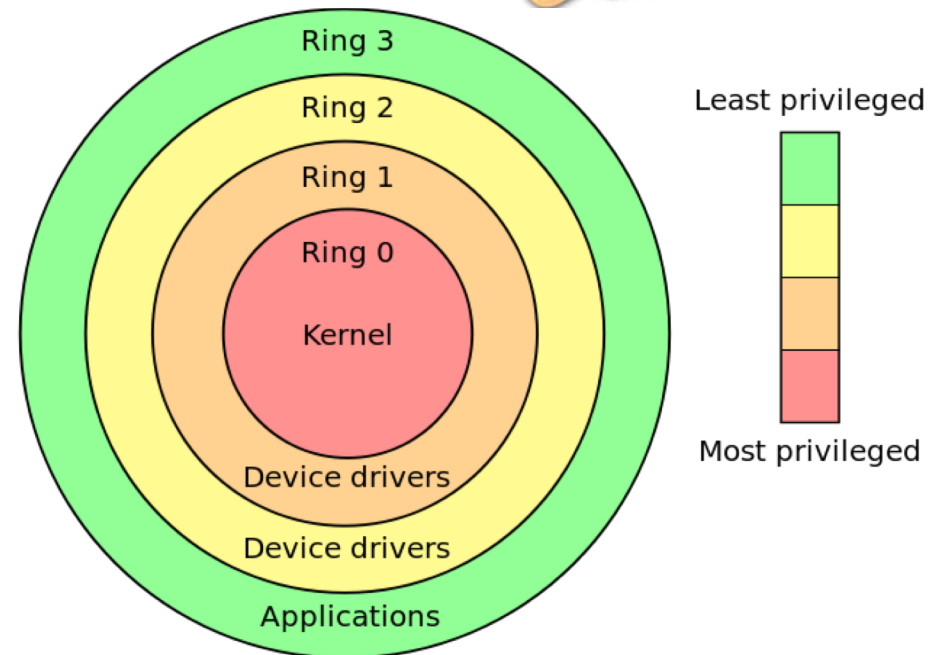
- Advantages:
 - Natural approach for user-defined objects
 - Eliminates confused deputy problems

- Disadvantages:
 - Review of permissions?
 - Revocation?
 - Delegation?
 - Privacy?

C-Lists

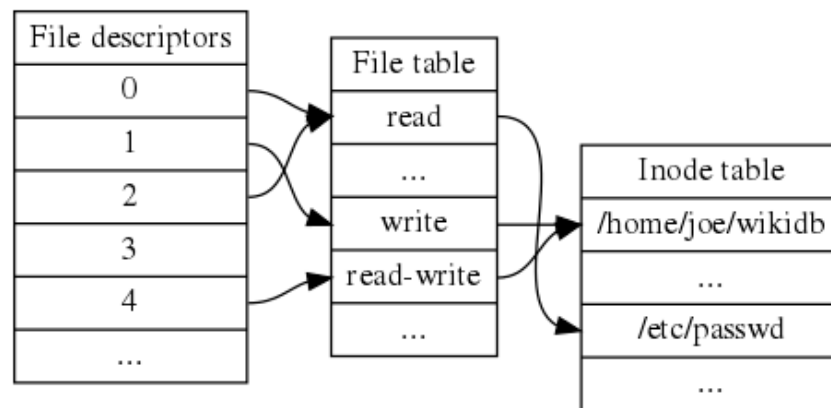
- OS maintains and stores stores list of capabilities $C_i = \langle O_i, Privs_i \rangle$ for each principal (process)

- 1) **Authorization:** OS mediates access to objects, checks process capabilities
- 2) **Unforgeability:** capabilities are stored in protected memory region (kernel memory)



Example: File Descriptor Table

- In Unix etc, a file descriptor is a handle used to reference files and I/O resources
- File descriptors have modes (read, write) and are stored in per-process file descriptor table
- File descriptors can be passed between processes using `sendmsg()`



Example: Google Fuchsia

- new OS in development by Google
- possibly intended as a universal across-platform OS for the IoT era (lots of speculation)
- capability-based microkernel embraces capabilities (handles) for all kernel objects
 - socket, port, virtual memory region, process, thread, etc.

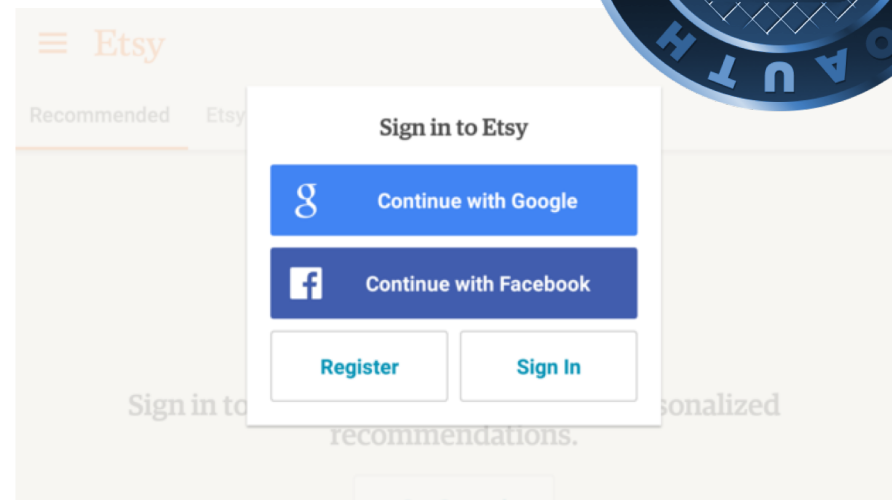


Fuchsia



Example: OAuth2

- Industry standard authorization protocol
- Used for single sign-on by major IDPs
 - Facebook, Google
- A **bearer token** contains a unique identifier



Cryptographically-protected capabilities

- Object owner creates capabilities using a digital signature scheme
- Capabilities are triples $C = \langle O, Privs, \text{Sig}(O, Privs; k_O) \rangle$
- **Authorization:** P is permitted to perform op on O if P produces a capability for O with $op \in Privs$ and a valid signature
- **Unforgeability:** digital signatures are unforgeable to adversaries who don't know private key k_O
- Note: assumes PKI

Restricted Delegation

- $C_0 = \langle O, Privs_0, k_1, s_0 \rangle$
 - where $s_0 = \text{Sig}(O, Privs_0, k_1; k_0)$
- $C_1 = \langle O, Privs_1, k_2, d_0, s_1 \rangle$
 - Where $d_0 = \text{Sig}(O, Privs_0, k_1; k_0)$ and $s_1 = \text{Sig}(O, Privs_1, k_2, d_0; k_1)$
- $C_2 = \langle O, Privs_2, k_3, d_0, d_1, s_2 \rangle$
 - Where $d_1 = \text{Sig}(O, Privs_1, k_2; k_1)$ and $s_2 = \text{Sig}(O, Privs_2, k_3, d_0, d_1; k_2)$

- $C_n = \langle O, Privs_n, k_{n+1}, d_0, \dots, d_{n-1}, s_n \rangle$
 - Where $d_i = \text{Sig}(O, Privs_i, k_{i+1}; k_i)$
 $s_i = \text{Sig}(O, Privs_i, k_{i+1}, s_0, \dots, s_{i-1}; k_i)$

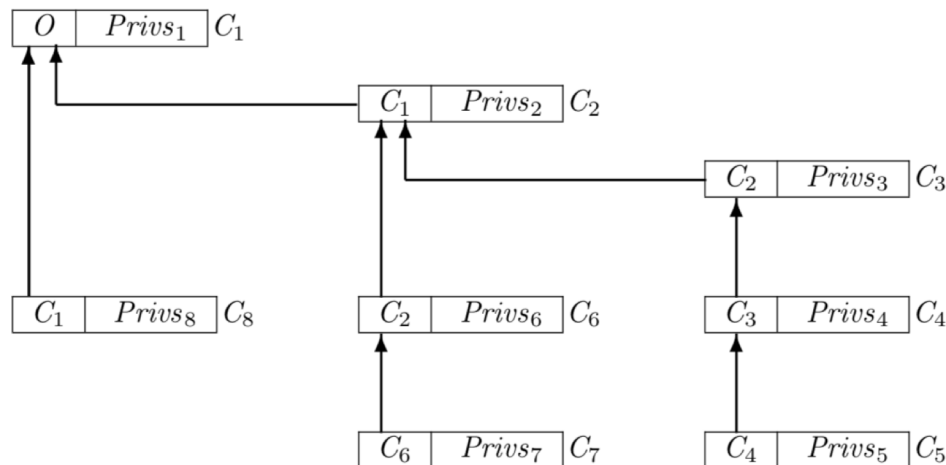
Revocation

- Revocation Tags

- Capabilities are tuples $C = \langle O, Privs, rt_c, \text{Sig}(O, Privs, rt_c; k) \rangle$
- Access to object O is guarded by a reference monitor; monitor maintains a list of revoked tags rt_c

- Capability Chains

- Objects can be other capabilities!
- P is authorized to perform op on O if P holds a capability C_i and $op \in Privs_k$ holds for every capability C_k in the chain from C_i to C_1



Keys as capabilities

- Encrypt object
- Decryption method functions as reference monitor:
 - **Authorization:** correct key will decrypt object -> allow access
 - **Unforgeability:** incorrect key will not decrypt
- Note: no notion of separate privileges

Example: Mac keychains

- OSX/iOS password manager
- uses password-based encryption (AES-256) to store username/password credentials
- supports multiple keychains



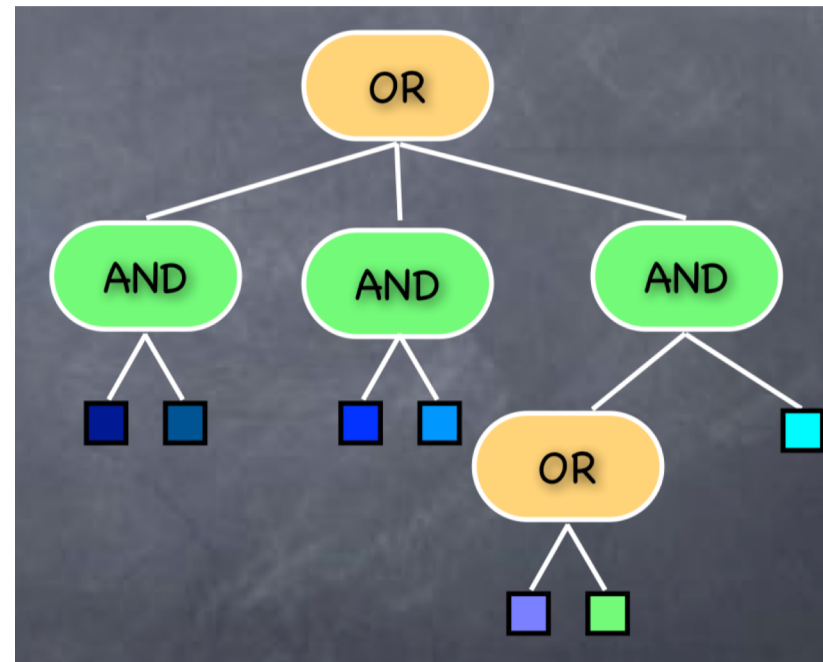
Example: CryptDB

- Encrypted database system. Inspiration for several application-grade encrypted database systems
- Processes database queries on encrypted data
- Uses chains of keys (starting with user password) to decrypt values/authorize users
 - onion encryption



Attribute-based encryption

- Type of public-key encryption in which secret keys depend on user attributes
- Users can only decrypt a ciphertext if they hold a key for appropriate attributes
- A KDC creates secret keys for users based on attributes



What about privacy?

