

OO Keywords

- Object
- Message
- Class
- Instance
- Method
- Subtype
- Subclass

Objects

- Internal data abstractions
- Hide representation
- Have associated state
- Methods have access to its state
- Self

Object Types

- Allow objects to be first class
- Allow use in assignment, parameters, components of structures
- Allow objects to be classified via subtyping

Classes

- Templates for creation of objects
 - Initialization code
 - Contain definitions for all methods
- Can modify or extend by creating subclasses.
- Can be used as types

Dynamic Method Invocation

- Each object responsible for keeping track of implementation of its own operations
- When evaluate o.m(...) at run-time, code run depends on operations associated with o.
- Static overloading is different
 - resolved at compile time.

Multi-methods

- Code executed depends on more than one argument.
- Example m(a,b) -- choice of code to be executed depends on run-time types of a and b.
- CLOS is example of such language
- Behavior and implementation quite different from single-dispatch languages.

Subtyping

- Already discussed
- Relation between interfaces, independent of implementations

Subclasses

Support incremental modification of code.

```
class Point
  var
    x = 0: Int;
  methods
    fun getx():int {return x}
    proc move(dx: int)
        {x := x + dx}
end class;
```

Subclasses

Static Overloading vs Dynamic Dispatch

- Dynamic dispatch object receiving message determines which code will be executed.
 - Determined at run-time.
- Static overloading occurs when an object supports two or more implementations of a message name -- generally w/different types.
 - Determined at compile-time.
 - e.g., moveTo(x,y) & moveTo(locn)
- Confusion when coexist in same language.

Overloading vs Dynamic Dispatch

```
class C { ...
 fun eq(other: Ctype):boolean {...} (1)
class SC of C modifying eq { ...
 fun eq(other: Ctype):boolean {...} (2) override
 fun eq(other: SCtype):boolean {...} (3) overload
}
c,c': Ctype; sc: SCtype;
c = new C;
                  c' = new SC;
                                    sc = new SC;
           What code is executed?
c.eq(c);
                  c.eq(c');
                                    c.eq(sc);
c'.eq(c);
                  c'.eq(c');
                                    c'.eq(sc);
                  sc.eq(c');
sc.eq(c);
                                    sc.eq(sc);
```

Type Restrictions

Type systems limit changes in method types in subclasses.

```
class C
    ...
    methods
    clone(): CType {...};
    equals(other: CType) {...};
end class
How can we change these types in subclasses?
```

More flexible subclasses

Why restrict changing types in subclasses?

Methods can be mutually recursive!

```
class Example
   :
   methods
   proc m(s:S,...) {... self.n(s) ...}
   fun n(x: S): T {...}
end class;
subclass SubExample of Example modifying n
    :
   methods
   fun n(x: S'): T' {...}
   proc newMethod(...){...}
end class;
```

What must be relationship of new type of n to old to preserve type safety?

Changing Types in Subclasses

Subtype will always be fine!

I.e., S <: S' and T' <: T equivalently, S' → T' <: S → T

E.g., can change return type of clone in subclass to type of objects from subclass.
Cannot specialize parameter types in equals *Binary methods*If subclass updates method types so they are subtypes of original then type-safe.
If restricted in this way, subclasses will always generate subtypes.

What about instance variables?

Instance variables can be values and receivers

- No subtypes!

If Circle has instance variable center: Point, ColorCircle's center must have same type.

Hard to redefine getCenter in ColorCircle, even if legal!

Important problem in OO language design and type theory

OO Languages

- Simula 67
- Smalltalk-72, -74, ... -80
- C++, Object Pascal, Object Cobol, ...
- Eiffel, Sather
- Java
- Scala
- Dart, Grace?

Simula 67

Simula 67

- First OO language
- You read in text
- Also added coroutines
- Use of "inner" rather than "super" in constructors

Inner

Class A; Begin startA; Inner; endA End; A Class B; Begin startB Inner; endB End; B Class C; Begin startC Inner; endC End; Ref(C) X;

X :- New C;

• Results in execution of:

startA startB startC endC endB endA

• Beta supports similar in all methods & classes

Smalltalk

Smalltalk

- New features:
 - Everything is an object, including classes
 - No operations -- only message-sending
 - Used to build customizable environment
 - Abstraction -- private instance variables, public methods
- Dynamically typed

Dynabook

- Laptop computer -- Alan Kay 1970's
 - Turing award 2003
- Proposed in 1970's aimed at children & adults
 - Neal Stephenson's "The Diamond Age or, a young lady's illustrated primer" is the next step
- Programmable environment
- Smalltalk as OS and programming language

Syntax

- n <- 3+4
 - send "+" message to 3 w/param 4 and insert in n
- •n between: 10 and: 100
 - send "between: and:" message to n w/params 10, 100
- •[:params | <message-expressions>]
 - lexical closure equiv to lambda expression

Smalltalk class

class name Point super class Object class var instance var х У class messages and methods !...names and code for methods..." instance messages and methods moveDx: dx Dy: dy || $x \leq dx + x$ $y \leq dy + y$ х х . . .



Dynamic Method Invocation

- Start with object's class and search up superclasses.
- When call method inside, start search from self again.
- Most other OO languages do not implement dmi in this way -- too inefficient!

Key ideas of Smalltalk

- Everything is an object
- Information hiding instance variables protected.
- Dynamic typing, so subtyping determined by whether can masquerade -- "message not understood"
- Inheritance distinct from subtyping

