

Lecture 2: Representing Integers

CS 105

Fall 2026

Review: Bits Require Interpretation

10001100 00001100 10101100 00000000

might be interpreted as

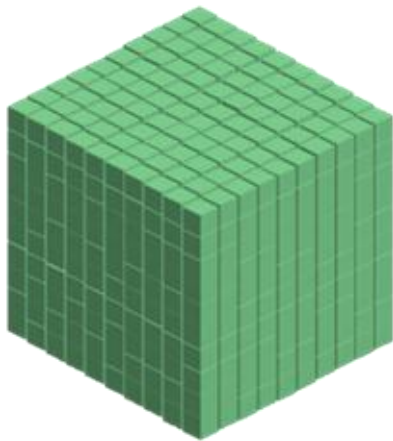
- The integer 3,485,745
- A floating point number close to 4.884569×10^{-39}
- The string "105"
- A portion of an image or video
- An address in memory

Representing Integers

- Arabic Numerals: 47
- Roman Numerals: XLVII
- Brahmi Numerals: 𑌔𑌖
- Tally Marks: 

Base-10 Integers

1000 (10^3)

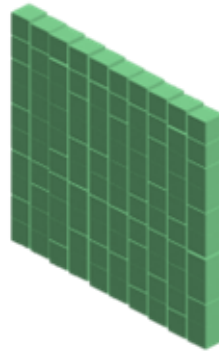


0

0

1

100 (10^2)



0

0

8

10 (10^1)



0

4

8

1 (10^0)

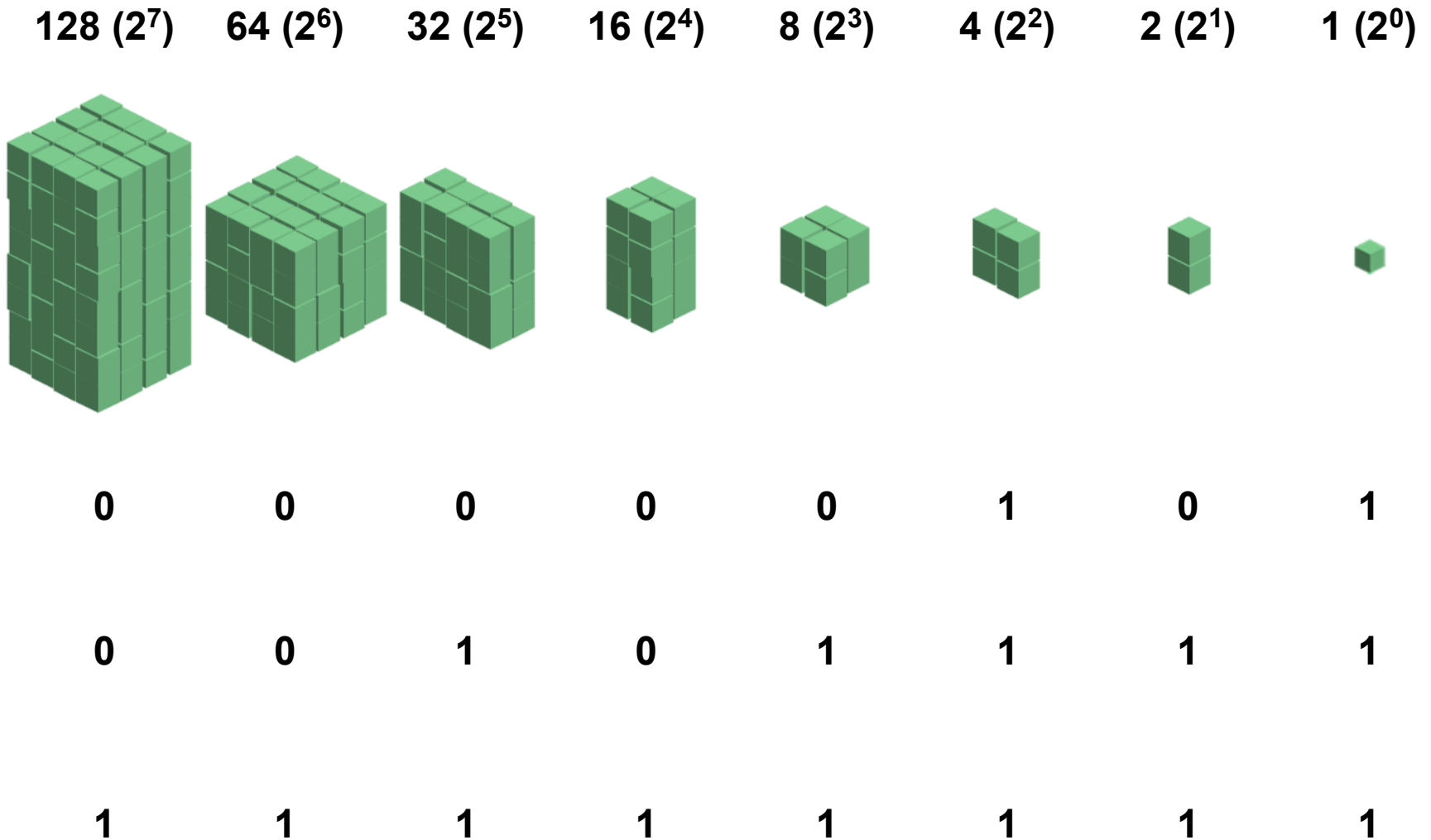


5

7

7

Base-2 Integers (aka Binary Numbers)



Binary Numbers

- Decimal (Base-10):

1011

$$\begin{aligned} &= 1 \cdot 10^3 + 0 \cdot 10^2 + 1 \cdot 10^1 + 1 \cdot 10^0 \\ &= 1011 \end{aligned}$$

- Binary (Base-2):

1011

$$\begin{aligned} &= 1 \cdot 2^3 + 0 \cdot 2^2 + 1 \cdot 2^1 + 1 \cdot 2^0 \\ &= 11 \end{aligned}$$

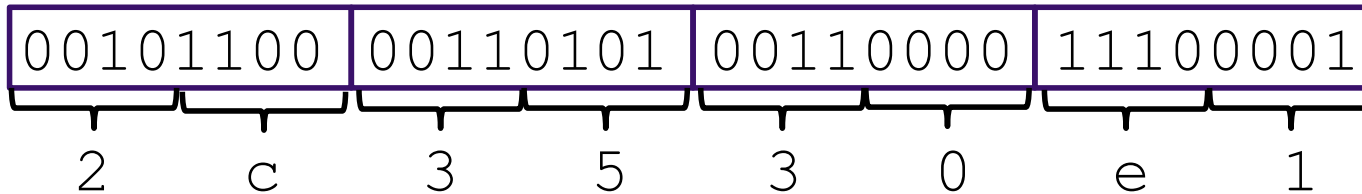
There are
10 types
of people
in the world:

Those who
understand binary,
and those
who don't.

Exercise 1: Binary Numbers

- Consider the following four-bit binary values. What is the (base-10) integer interpretation of these values?
 1. 0001
 2. 1010
 3. 0111
 4. 1111

Hexadecimal Numbers



0x2c3530e1

Dec	Hex
0	0
1	1
2	2
3	3
4	4
5	5
6	6
7	7
8	8
9	9
10	a
11	b
12	c
13	d
14	e
15	f

Exercise 2: Hexadecimal Numbers

- Consider the following hexadecimal values. What is the representation of each value in binary?
 1. 0x0a
 2. 0x11
 3. 0x2f

Addition Example

- Compute $5 + 6$ assuming all ints are stored as eight-bit (1 byte) unsigned values

$$\begin{array}{r} 1 \\ 00000101 \\ + 00000110 \\ \hline 00001011 \end{array} = 11 \text{ (Base-10)}$$

Like you learned in grade school, only binary!
... and with a finite number of digits

Addition Example with Overflow

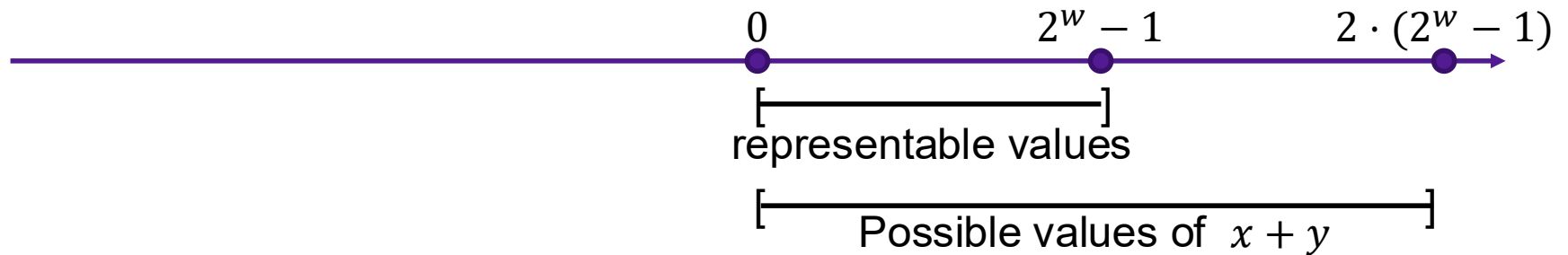
- Compute $200 + 100$ assuming all ints are stored as eight-bit (1 byte) unsigned values

$$\begin{array}{r} 11 \\ 11001000 \\ + 01100100 \\ \hline 00101100 \end{array} = 44 \text{ (Base-10)}$$

Like you learned in grade school, only binary!
... and with a finite number of digits

Error Cases (Unsigned Ints)

- Assume w -bit unsigned values



- $$x + y = \begin{cases} x + y & \text{(normal)} \\ x + y - 2^w & \text{(overflow)} \end{cases} \quad \left. \vphantom{\begin{cases} x + y \\ x + y - 2^w \end{cases}} \right\} \text{modular addition}$$

- overflow has occurred iff $x + y < x$

Exercise 3: Binary Addition

- Given the following 5-bit unsigned values, compute their sum and indicate whether or not an overflow occurred

x	y	x+y	overflow?
00010	00101		
01100	00100		
10100	10001		

Multiplication Example

- Compute 5 x 6 assuming all ints are stored as eight-bit (1 byte) unsigned values

$$\begin{array}{r} 00000101 \\ \times 00000110 \\ \hline 00000000 \\ 000001010 \\ + 0000010100 \\ \hline 00011110 \end{array} = 30 \text{ (Base-10)}$$

Like you learned in grade school, only binary!
... and with a finite number of digits

Multiplication Example

- Compute 200×3 assuming all ints are stored as eight-bit (1 byte) unsigned values

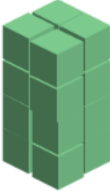
$$\begin{array}{r} 11001000 \\ \times 00000011 \\ \hline 11001000 \\ + 110010000 \\ \hline 1001011000 = 88 \text{ (Base-10)} \end{array}$$

Like you learned in grade school, only binary!
... and with a finite number of digits

What about
positive/negative
integers?

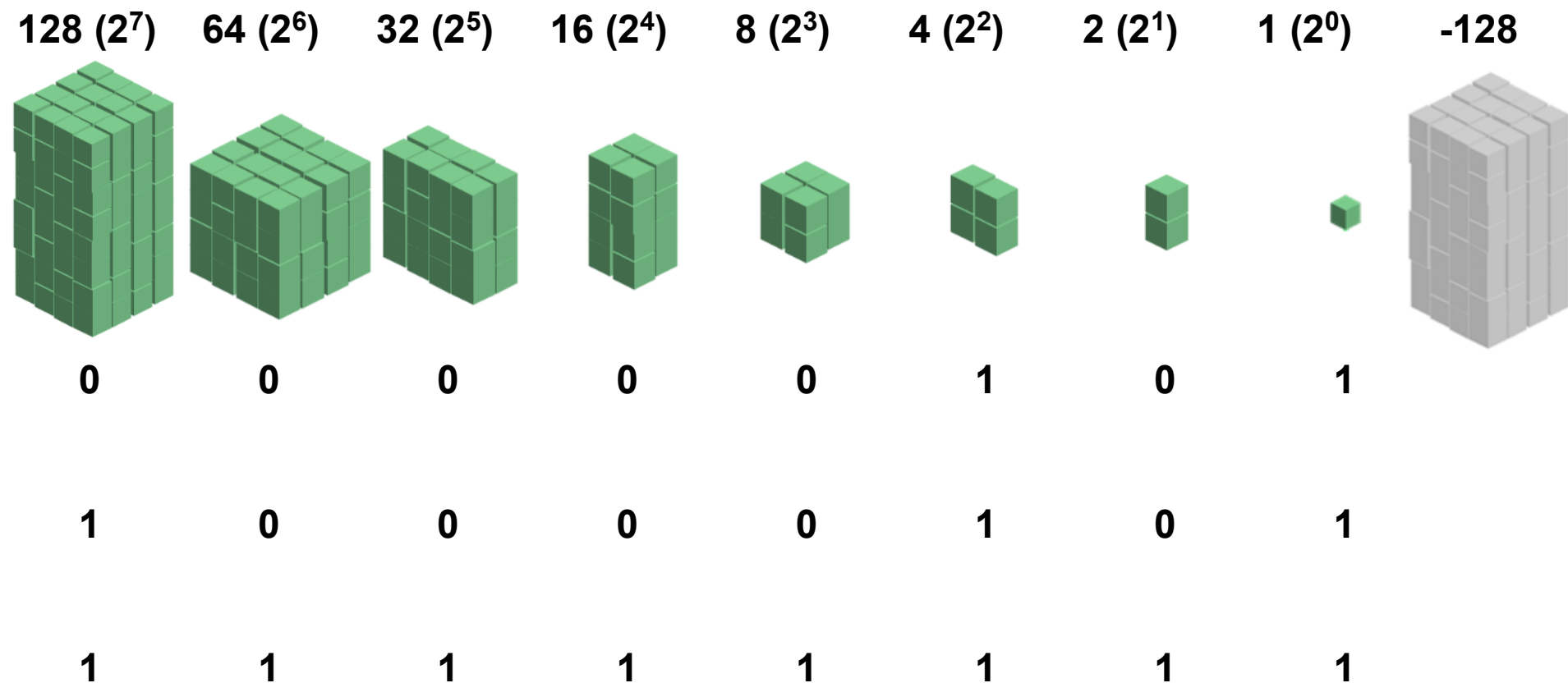
Representing Signed Integers

- Option 1: sign-magnitude
 - One bit for sign; interpret rest as magnitude
 - $Signed(x) = (-1)^{x_{w-1}} \cdot \sum_{i=0}^{w-2} x_i \cdot 2^i$

+/-	64 (2^6)	32 (2^5)	16 (2^4)	8 (2^3)	4 (2^2)	2 (2^1)	1 (2^0)
—							
0	0	0	0	0	1	0	1
1	0	0	0	0	1	0	1
1	1	1	1	1	1	1	1

Representing Signed Integers

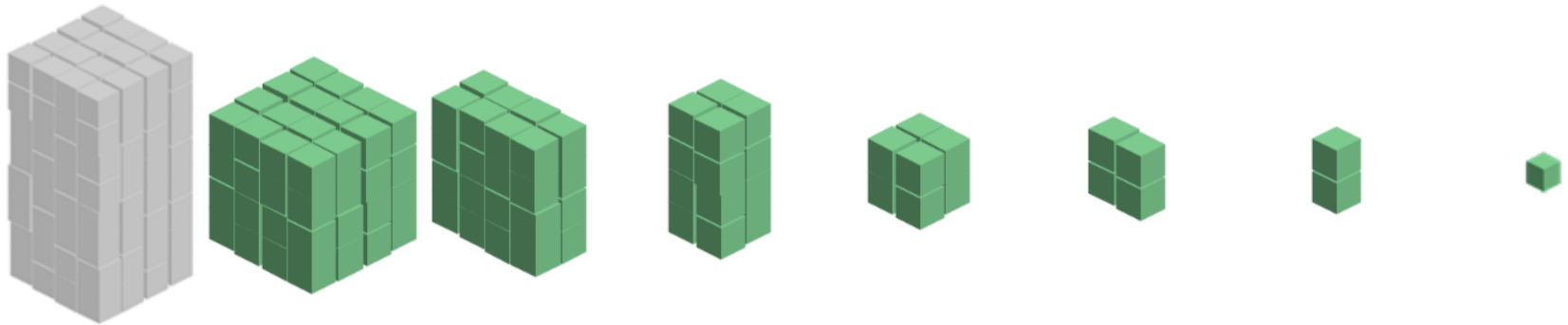
- Option 2: excess-K
 - Choose a positive K in the middle of the unsigned range
 - $Signed(x) = \sum_{i=0}^{w-1} x_i \cdot 2^i - 2^{w-1}$



Representing Signed Integers

- Option 3: two's complement
 - Like unsigned, except the high-order contribution is *negative*
 - $Signed(x) = -x_{w-1} \cdot 2^{w-1} + \sum_{i=0}^{w-2} x_i \cdot 2^i$

-128 (-2^7) **64 (2^6)** **32 (2^5)** **16 (2^4)** **8 (2^3)** **4 (2^2)** **2 (2^1)** **1 (2^0)**



0	0	0	0	0	1	0	1
1	0	0	0	0	1	0	1
1	1	1	1	1	1	1	1

Important Signed Integers

	8	16	32	64
TMax	0x7F	0x7FFF	0x7FFFFFFF	0x7FFFFFFFFFFFFFFF
TMin	0x80	0x8000	0x80000000	0x8000000000000000
0	0x00	0x0000	0x00000000	0x0000000000000000
-1	0xFF	0xFFFF	0xFFFFFFFF	0xFFFFFFFFFFFFFFFF

Exercise 4: (Signed) Binary Numbers

- Consider the following one-byte binary values. What is the (signed) integer interpretation of these values?
 1. 01101001
 2. 10010111
 3. 01000010
 4. 10111110
- What is the one-byte binary representation of the following integers?
 - 47
 - -47

$$-x == -x+1-1 == -1-x+1 == 111...1-x+1 = \sim x+1$$

Integers in C

C Data Type	Size (bytes)
unsigned char	1
unsigned short	2
unsigned int	4
unsigned long	8

C Data Type	Size (bytes)
char	1
short	2
int	4
long	8

Addition Example

- Compute $5 + -3$ assuming all ints are stored as four-bit signed values

$$\begin{array}{r} 1 1 \\ 0 1 0 1 \\ + 1 1 0 1 \\ \hline 0 0 1 0 \end{array} = 2 \text{ (Base-10)}$$

Exactly the same as unsigned numbers!

... but with different error cases

Addition/Subtraction with Overflow

- Compute $5 + 6$ assuming all ints are stored as four-bit signed values

$$\begin{array}{r} 1 \\ 0101 \\ + 0110 \\ \hline 1011 \end{array} = -5 \text{ (Base-10)}$$

Multiplication Example

- Compute 3×2 assuming all ints are stored as four-bit signed values

$$\begin{array}{r} 0011 \\ \times 0010 \\ \hline 0000 \\ + 00110 \\ \hline 0110 \end{array} = 6 \text{ (Base-10)}$$

Exactly like unsigned multiplication!
... except with different error cases

Multiplication Example

- Compute 5×2 assuming all ints are stored as four-bit signed values

$$\begin{array}{r} 0101 \\ \times 0010 \\ \hline 0000 \\ + 01010 \\ \hline 1010 \end{array} = -6 \text{ (Base-10)}$$

Exercise 5: (Signed) Binary Multiplication

- Given the following 3-bit signed values, compute their product and indicate whether or not an overflow occurred

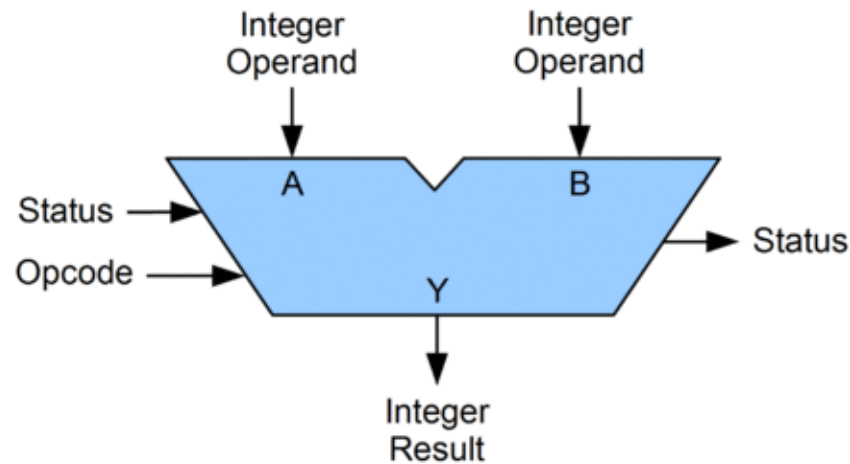
x	y	x*y	overflow?
100	101		
010	011		
111	010		

Casting between Numeric Types

- Casting from shorter to longer types preserves the value
- Casting from longer to shorter types drops the high-order bits
- Casting between signed/unsigned types preserves the bits (it just changes the interpretation)
- Implicit casting occurs in assignments and parameter lists. In mixed expressions, signed values are implicitly cast to unsigned
 - Source of many errors!

Arithmetic Logic Unit (ALU)

- circuit that performs bitwise operations and arithmetic on integer binary types



Review: Bitwise vs Logical Operations

- Bitwise Operators `&`, `|`, `~`, `^`
 - View arguments as bit vectors
 - operations applied bit-wise in parallel
- Logical Operators `&&`, `||`, `!`
 - View 0 as “False”
 - View anything nonzero as “True”
 - Always return 0 or 1
 - **Early termination**
- Shift operators `<<`, `>>`
 - Left shift fills with zeros
 - For unsigned integers, right shift is logical (fills with zeros)

Exercise 6: Bitwise vs Logical Operations

- What is the binary representation of each of the following expressions? Assume signed char data type (one byte).

1. $\sim(-30)$

2. $-30 \& 22$

3. $-30 \&\& 22$

4. $22 \ll 1$

5. $22 \gg 1$

6. $-30 \gg 1$

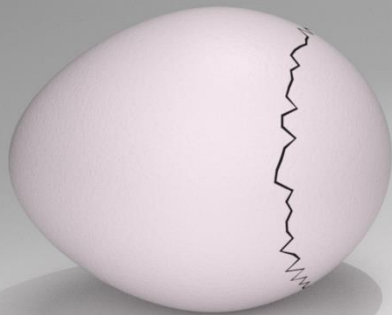
Multiplying with Shifts

- Multiplication is slow
- Bit shifting is kind of like multiplication, and is often faster
 - $x * 8 = x \ll 3$
 - $x * 10 = x \ll 3 + x \ll 1$
- Most compilers will automatically replace multiplications with shifts where possible

Endianness

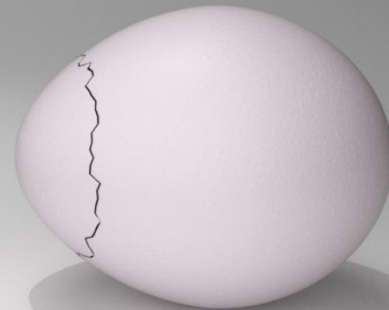
47 vs 74

The way traditionally
lilliputians broke their boiled eggs
on the larger end



BIG ENDIAN

The way the king then ordered
lilliputians to break their boiled eggs
on the smaller end



Little ENDIAN

Endianness

- **Big Endian:** low-order bits go on the right (47)
 - I tend to think in big endian numbers, so examples in class will generally use this representation
 - Networks generally use big endian (aka network byte order)
- **Little Endian:** low-order bits go on the left (74)
 - Most modern machines use this representation
- I will try to always be clear about whether I'm using a big endian or little endian representation
- When in doubt, ask!