Lecture 32: Concurrency II

CS 62
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Some slides based on those from Dan Grossman, U. of Washington
Race Conditions

• A race condition occurs when the computation result depends on scheduling (how threads are interleaved)
  • If T1 and T2 happened to get scheduled in a certain way, things go wrong
  • Since we do not control scheduling, we need to write programs that work independent of scheduling

• Race conditions are bugs that exist only due to concurrency
  • No interleaved scheduling problems with only 1 thread.

• Typically, problem is that some intermediate state can be seen by another thread; screws up other thread.
Data Races vs Bad Interleavings

• We will make a big distinction between these terms
• Both are kinds of race-condition bugs
• Confusion often results from not distinguishing these or using the ambiguous “race condition” to mean only one
Data races (briefly)

- A **data race** is a specific type of *race condition* that can happen in 2 ways:
  - Two different threads *potentially* write a variable at the same time
  - One thread *potentially* writes a variable while another reads the variable
- Not a race: simultaneous reads provide no errors
- “Potentially” is important
  - We claim the code itself has a data race independent of any particular actual execution
- Data races are bad, but we can still have a race condition, and bad behavior, when no data races are present...through **bad interleavings** (what we will discuss now).
Stack Example

class Stack<E> {
    private E[] array;
    private int index = 0;
    Stack(int size) {
        array = (E[]) new Object[size];
    }
    synchronized boolean isEmpty() {
        return index==0;
    }
    synchronized void push(E val) {
        if(index==array.length)
            throw new StackFullException();
        array[index++] = val;
    }
    synchronized E pop() {
        if(index==0)
            throw new StackEmptyException();
        return array[--index];
    }
}
Let’s implement `peek()`

```java
synchronized E peek() {
    if (index == 0)
        throw new StackEmptyException();
    return array[index-1];
}
```

```
correct
```
Example of race condition, not data race

class C {
    static <E> E myPeekHelperWrong(Stack<E> s) {
        E ans = s.pop();
        s.push(ans);
        return ans;
    }
}

• *No overall* effect on the shared data. State should be the same at the end
• But the way it is implemented creates an inconsistent *intermediate state*
• There is still a *race condition* though. This intermediate state should not be exposed → *bad interleavings*
myPeekHelperWrong() and isEmpty()

Thread 1

boolean b = isEmpty();

Thread 2 (calls myPeekHelperWrong)

E ans = stk.pop();
stk.push(ans);
return ans;

array = 
0 1 2 3 4 5 6 7
12

index = 0

array = 
0 1 2 3 4 5 6 7

boolean b = isEmpty();

index = 1

array = 12
0 1 2 3 4 5 6 7
myPeekHelperWrong() and push()

Thread 1
-------
stk.push(x);
stk.push(y);
E z = stk.pop();

Thread 2 (calls myPeekHelperWrong)
-----------------------------------
E ans = stk.pop();
stk.push(ans);
return ans;
**myPeekHelperWrong()** and **pop()**

**Thread 1**

```
stk.push(x);
stk.push(y);
E z = stk.pop();
```

**Thread 2 (calls myPeekHelperWrong)**

```
E ans = stk.pop();
stk.push(ans);
return ans;
```
myPeekHelperWrong() and myPeekHelperWrong() on 1 element

Thread 1
--------
E ans = stk.pop();
stk.push(ans);
return ans;

Thread 2 (calls myPeekHelperWrong)
----------------------------------
E ans = stk.pop(); // exception!
myPeekHelperWrong() and myPeekHelperWrong() on > 1 element

---

Thread 1
---

E ans = stk.pop();
stk.push(ans);
return ans;

---

Thread 2 (calls myPeekHelperWrong)
---

E ans = stk.pop();
stk.push(ans);
return ans;
The fix

- **peek** needs synchronization to disallow interleavings
  - The key is to make a *larger critical section*
  - That intermediate state of **peek** needs to be protected
- Use re-entrant locks; will allow calls to **push** and **pop**
- Code on right is example of a peek external to the **Stack** class

```java
class Stack<E> {
    synchronized E peek() {
        E ans = pop();
        push(ans);
        return ans;
    }
}
class C {
    <E> E myPeek(Stack<E> s) {
        synchronized (s) {
            E ans = s.pop();
            s.push(ans);
            return ans;
        }
    }
}
```
The wrong fix

- **Focus so far**: problems from peek doing writes that lead to an incorrect intermediate state
- **Tempting but wrong**: If an implementation of peek (or isEmpty) does not write anything, then maybe we can skip the synchronization?
- Does not work due to *data races* with push and pop...
class Stack<E> {
    private E[] array = (E[])new Object[SIZE];
    int index = -1;
    boolean isEmpty() { // unsynchronized: wrong?!
        return index == -1;
    }
    synchronized void push(E val) {
        array[++index] = val;
    }
    synchronized E pop() {
        return array[index--];
    }
    E peek() { // unsynchronized: wrong!
        return array[index];
    }
}
Why wrong?

• It looks like `isEmpty` and `peek` can “get away with this” since `push` and `pop` adjust the state “in one tiny step”

• But this code is still wrong and depends on language-implementation details you cannot assume
  • Even “tiny steps” may require multiple steps in the implementation:
    • `array[++index] = val;` probably takes at least two steps
      • Code has a data race, allowing very strange behavior

• Moral: Do not introduce a data race, even if every interleaving you can think of is correct
Getting it right

• Avoiding race conditions on shared resources is difficult
  • What “seems fine” in a sequential world can get you into trouble when multiple threads are involved.
  • Decades of bugs have led to some conventional wisdom: general techniques that are known to work

• Next we discuss this conventional wisdom!
3 choices

- For every memory location (e.g., object field) in your program, you must obey at least one of the following:
  1. Thread-local: Do not use the location in > 1 thread
  2. Immutable: Do not write to the memory location
  3. Shared-and-mutable: Use synchronization to control access to the location
1. Thread-local

• Whenever possible, do not share resources
  • Easier to have each thread have its own thread-local copy of a resource than to have one with shared updates
  • This is correct only if threads do not need to communicate through the resource
    • That is, multiple copies are a correct approach
• Note: Because each call-stack is thread-local, never need to synchronize on local variables
• In typical concurrent programs, the vast majority of objects should be thread-local: shared-memory should be rare – minimize it!
2. Immutable

- Whenever possible, don’t update fields of objects
  - Make new objects instead
- One of key tenets of functional programming
  - You did study this in 52/54
  - Generally helpful to avoid side-effects
  - Much more helpful in a concurrent setting
- If a location is only read, never written, no synchronization is necessary!
  - Simultaneous reads are not races and not a problem
- *Programmers over-use mutation – minimize it!*
3. The rest: keep it synchronized

- After minimizing the amount of memory that is (1) thread-shared and (2) mutable, we need guidelines for how to use locks to keep other data consistent

- **Guideline**: No data races
  - *Never allow two threads to read/write or write/write the same location at the same time* (use locks!)
  - Even if it ‘seems safe’

- Necessary: A Java or C program with a data race is almost always wrong

- *But Not sufficient*: Our **peek** example had no data races, and it’s still wrong...
Worse than you think

Assertion always true w/ single threaded.
• Looks always true for multithreaded.
• OK if \( f \) not called at all
• OK after \( f \) completes
• Looks OK if in middle of \( f \)
• But has race condition

```java
class C {
  private int x = 0;
  private int y = 0;

  void f() {
    x = 1; // line A
    y = 1; // line B
  }

  void g() {
    int a = y; // line C
    int b = x; // line D
    assert(b >= a);
  }
}
```
Memory reordering

• For performance reasons, compiler and hardware reorder memory operations.

• But, but, ...
  • Compiler/hardware will never perform a memory reordering that affects the result of a single-threaded program
  • The compiler/hardware will never perform a memory reordering that affects the result of a data-race-free multi-threaded program

• So: If no interleaving of your program has a data race, then need not worry: result will be equivalent to some interleaving